

Quiz 4: today

lab 3: Posted, due Friday

New Problem in graph Theory:

Single Source Shortest Paths (SSSP)

Defn

Given $x, y \in V(G)$, the distance from x to y is

$$d(x, y) = \begin{cases} \text{min length of an } x-y \\ \text{Path in } G, \text{ if } y \text{ is reachable} \\ \text{from } x. \\ \infty \text{ if } y \text{ not reachable} \\ \text{from } x \end{cases}$$

SSSP Problem:

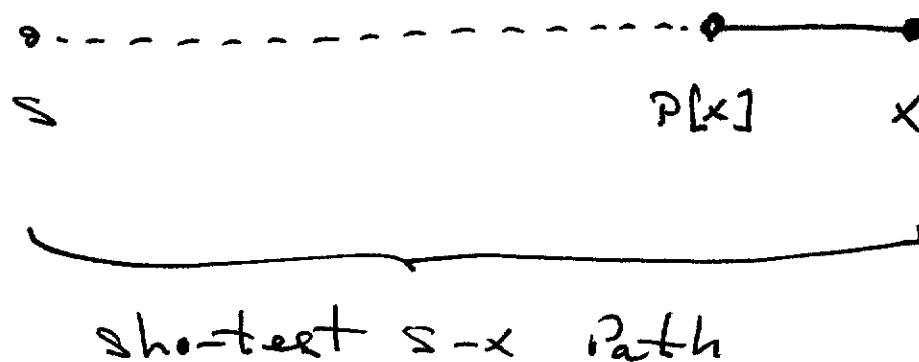
Given $s \in V(G)$, determine $d(s, x)$ for all $x \in V(G)$. Further, if $d(s, x) < \infty$ (i.e. x is reachable from s) determine a shortest s - x path in G .

Solution: Breadth First Search (BFS).

BFS is basically reachable() all over again.

Add 2 new attributes.

- A dictionary : distance
distance[x] gives $\delta(s, x)$, the length of a shortest $s-x$ Path.
- A dictionary : Predecessor
predecessor[x] gives the predecessor vertex of x along a shortest $s-x$ Path.

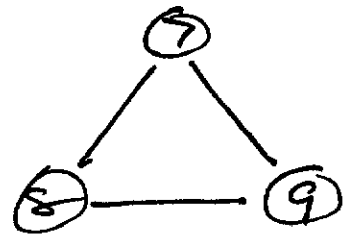
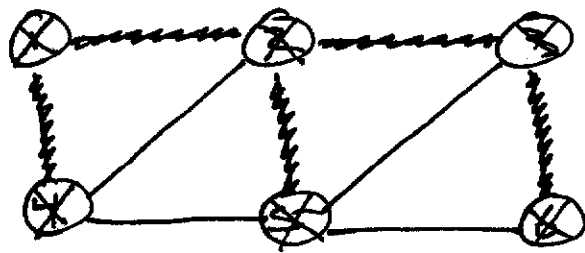


When a new vertex y is discovered,
we set $\text{Predecessor}[y]$ to be
the vertex x from which y was
discovered, i.e. as a member of
 $\text{Adj}[x]$.

we also set $\text{distance}[y]$ to be
one more than $\text{distance}[x]$.

$$\text{distance}[y] = \text{distance}[x] + 1$$

Ex. (same)



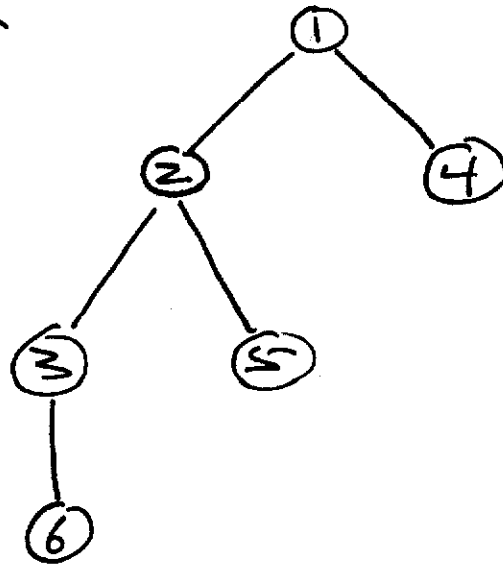
Source = 1

discovered queue:

$\neq \neq \neq \neq \neq \neq \neq \neq$

Predecessor
subgraph:

also called a
shortest
Path tree



dist:

0

1

2

3

Shortest 1-6 Path is

1, 2, 3, 6

See Recursive function

getPath()

Exercises

- run on same graph, different source.
- run on different graphs.

Next time!

• Graph coloring

• PaG