

Asymptotic growth rates of functions

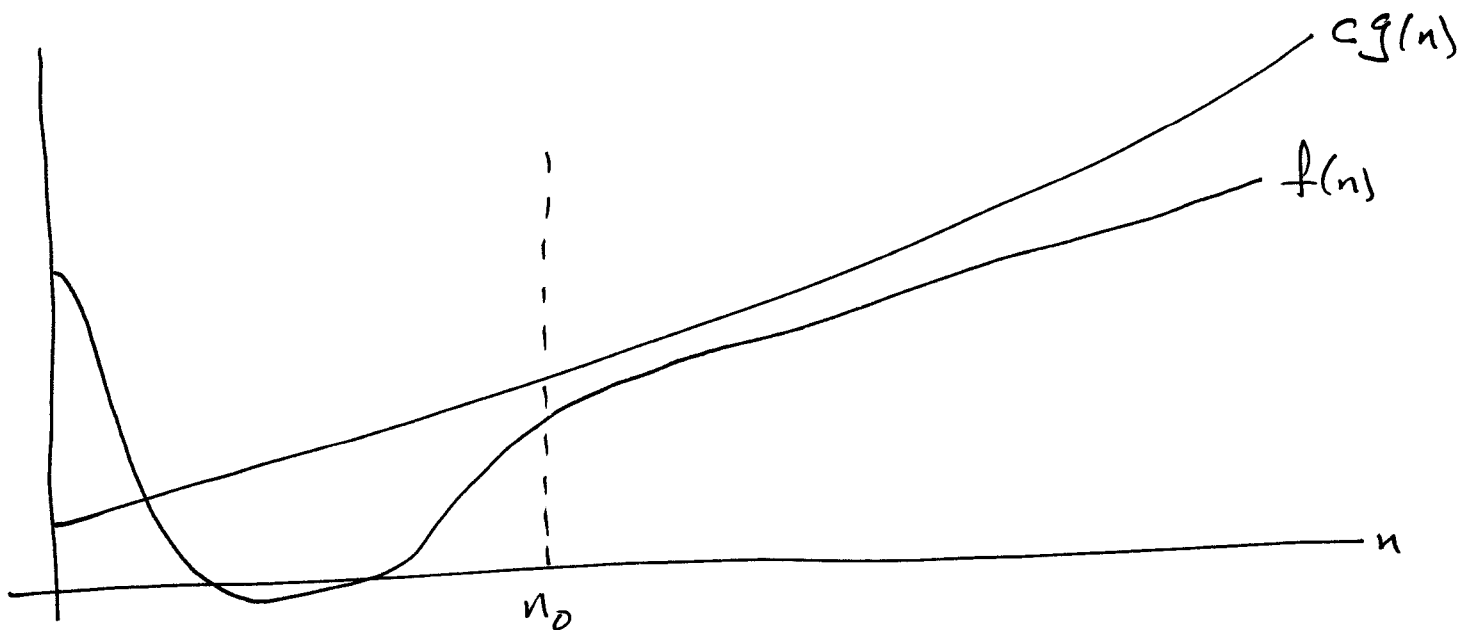
Defn

let $g(n)$ be a function.

$$O(g(n)) = \{f(n) \mid \exists c > 0, \exists n_0 > 0, \forall n \geq n_0 : 0 \leq f(n) \leq cg(n)\}$$

we say: $f(n)$ is Big O of $g(n)$

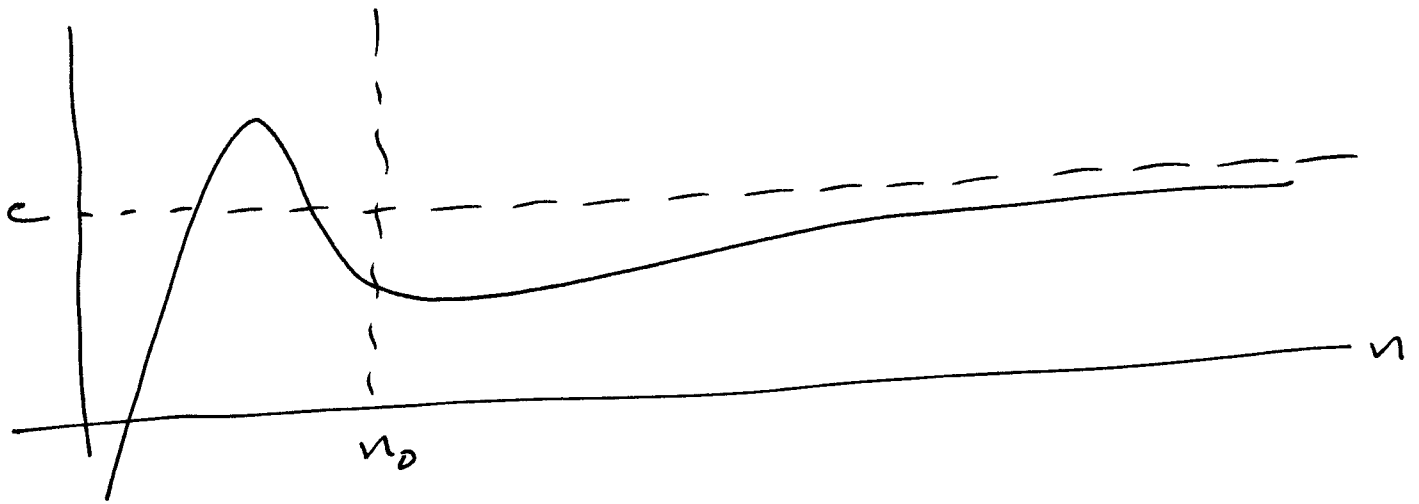
.. : $g(n)$ is an asymptotic upper bound for $f(n)$



equivalently:

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$$\exists c > 0, \exists n_0 > 0, \forall n \geq n_0 : 0 \leq \frac{f(n)}{g(n)} \leq c$$



Defn $f(n)$ is asymptotically non-negative (a.n.n) iff

$$\exists n_0 > 0, \forall n \geq n_0 : f(n) \geq 0.$$

Defn $f(n)$ is asymptotically positive (a.p) iff

$$\exists n_0 > 0, \forall n \geq n_0 : f(n) > 0.$$

Blanket assumption: all fns. under discussion are a.n.n.

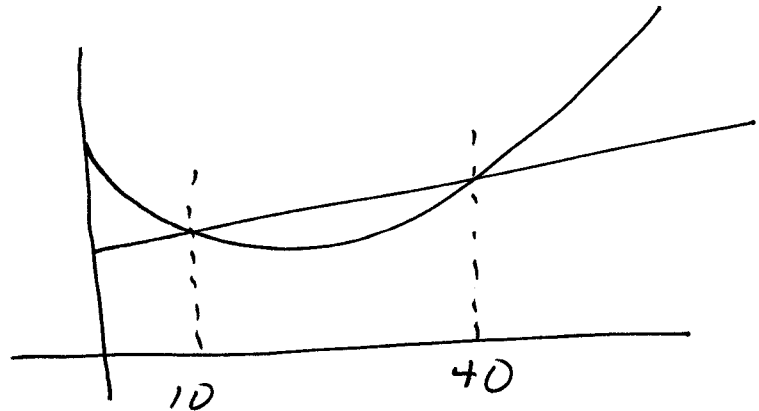
(occasionally we will need a.p.)

Ex $10n + 100 = O(n^2 - 40n + 500)$
 $\uparrow$
 means $\in$

why? observe

$$0 \leq (10n + 100) \leq (n^2 - 40n + 500)$$

for all $n \geq 40$.



In this case $n_0 = 40$, $c = 1$.

In fact $an+b = O(cn^2+dn+e)$ for any a, b, c, d, e with $a > 0, c > 0$.

Also $P(n) = O(Q(n))$ whenever $P(n), Q(n)$ are Polynomials with

$$\deg(P(n)) \leq \deg(Q(n)).$$

Defn

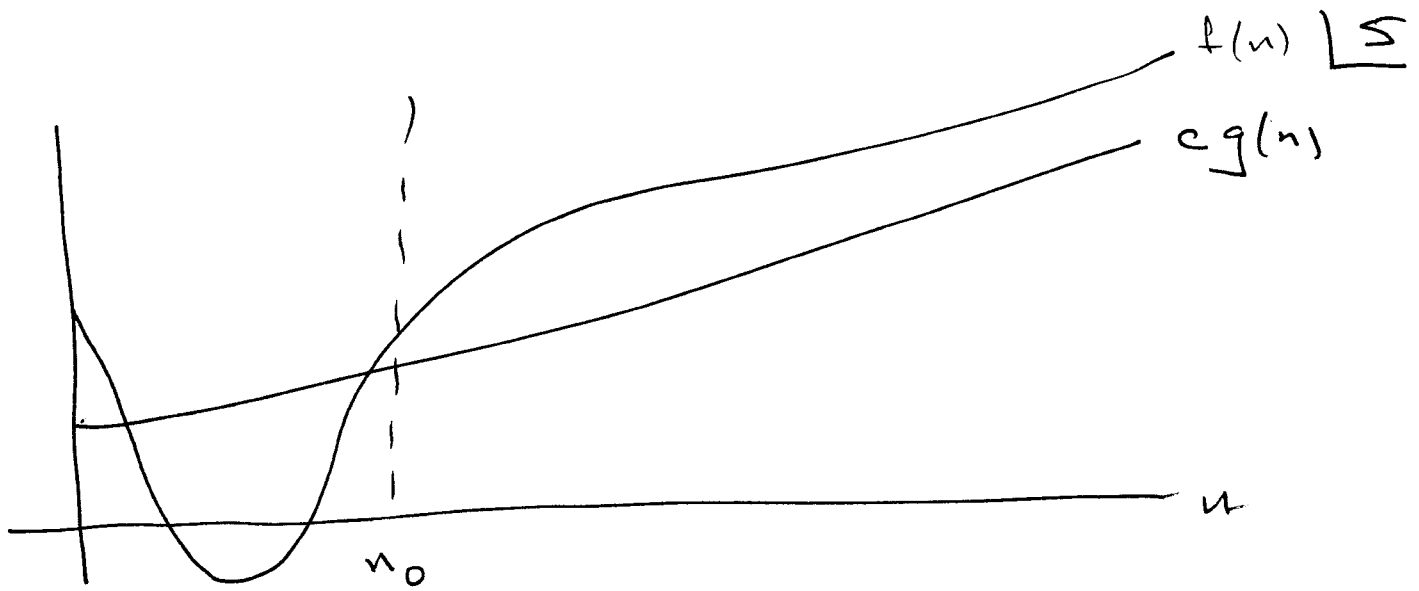
$$\Omega(g(n)) = \{f(n) \mid \exists c > 0, \exists n_0 > 0, \forall n \geq n_0 : 0 \leq c g(n) \leq f(n)\}$$

we say $f(n)$ is big omega of $g(n)$

" $g(n)$ is an asym. lower bound.

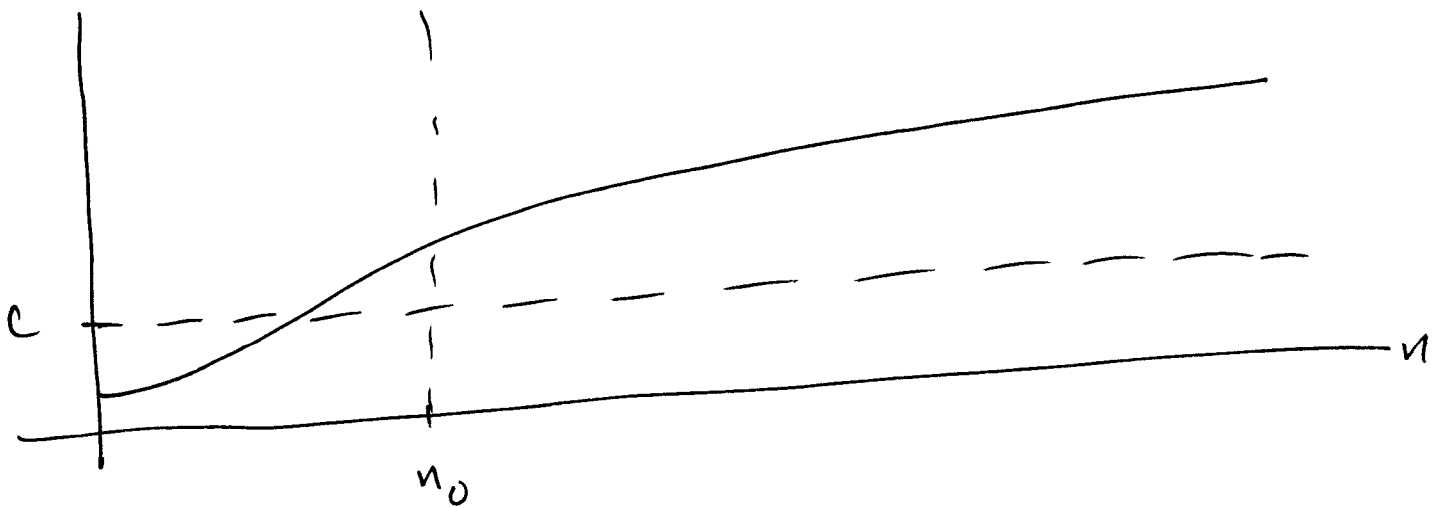
We write: $f(n) = \Omega(g(n))$

↑
means $f(n) \in \Omega(g(n))$.



equivalently:

$$\forall n \geq n_0 : 0 \leq c \leq \frac{f(n)}{g(n)}$$



Theorem

$$f(n) = O(g(n)) \iff g(n) = \Omega(f(n))$$

Proof of $(\Rightarrow)$:

Suppose $f(n) = O(g(n))$. Then there exist positive c, n_1 s.t.

$$(1) \quad \forall n \geq n_1 : 0 \leq f(n) \leq c_1 g(n)$$

we must show there exist positive c_2, n_2 s.t.

$$(2) \quad \forall n \geq n_2 : 0 \leq c_2 f(n) \leq g(n)$$

$$\text{Let } c_2 = \frac{1}{c_1} \text{ and } n_2 = n_1$$

Then $(1) \Rightarrow (2)$, and proof is complete. ~~■~~

Defn

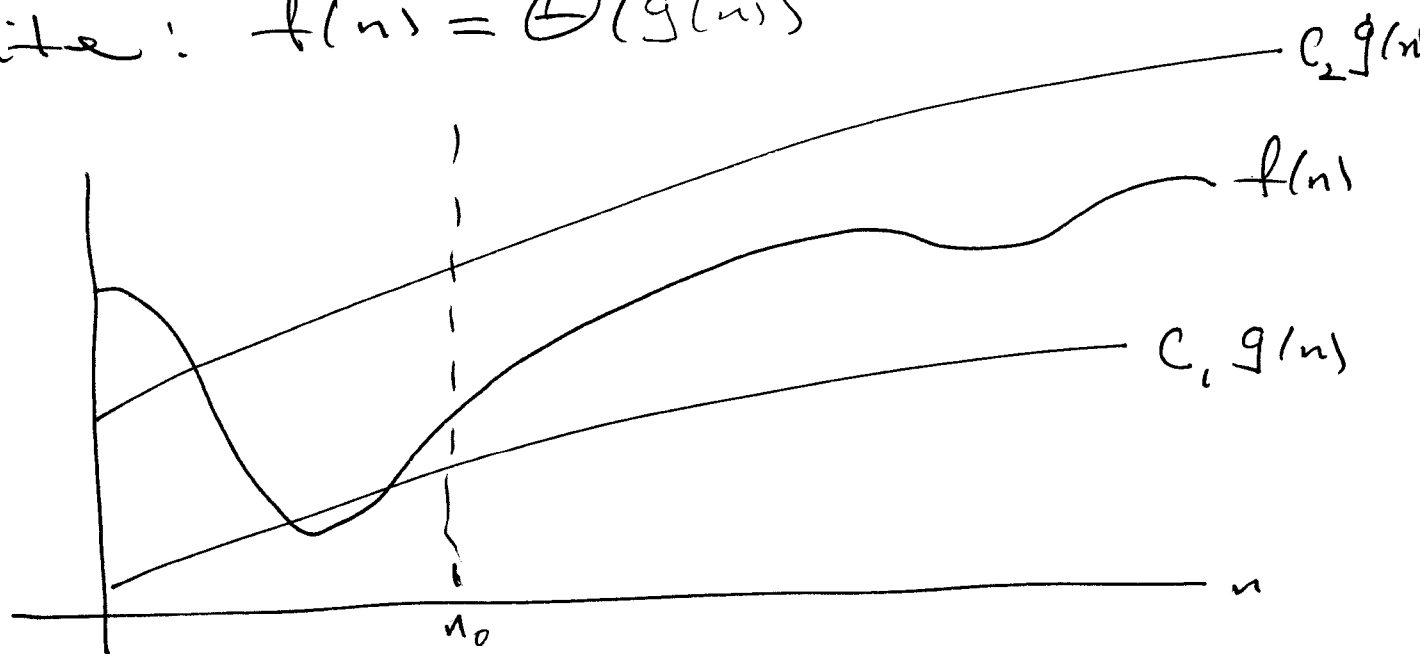
$$\Theta(g(n)) = \{f(n) \mid \exists c_1 > 0, c_2 > 0, n_0 > 0, \forall n \geq n_0:$$

$$0 \leq c_1 g(n) \leq f(n) \leq c_2 g(n)\}$$

we say: $f(n)$ is big theta of $g(n)$

... $g(n)$ is a tight asym. bound
for $f(n)$

write: $f(n) = \Theta(g(n))$



note: $\Theta(g(n)) = O(g(n)) \cap \Omega(g(n))$

equivalently (if f, g are a.p.)

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$$\forall n \geq n_0 : 0 \leq c_1 \leq \frac{f(n)}{g(n)} \leq c_2$$

