

handout on common funcs. (sec. 3.2)

Floor & ceiling funcs.

let $x \in \mathbb{R}$

$$(1) \lfloor x \rfloor = \text{greatest int. } \leq x$$

$$(2) \lceil x \rceil = \text{least int } \geq x$$

alternative definition

let $N \in \mathbb{Z}, x \in \mathbb{R}$.

$$(1) N = \lfloor x \rfloor \text{ iff } N \leq x < N+1$$

$$(2) N = \lceil x \rceil \text{ iff } N-1 < x \leq N$$

lemma 1: let $x \in \mathbb{R}$, $a, b \in \mathbb{Z}$. then

$$(1) \quad a \leq x < b \iff a \leq \lfloor x \rfloor < b$$

$$(2) \quad a < x \leq b \iff a < \lceil x \rceil \leq b$$

lemma 2: if $x \in \mathbb{R}$, $m \in \mathbb{Z}^+$, then

$$(1) \quad \left\lfloor \frac{\lfloor x \rfloor}{m} \right\rfloor = \left\lfloor \frac{x}{m} \right\rfloor$$

$$(2) \quad \left\lceil \frac{\lceil x \rceil}{m} \right\rceil = \left\lceil \frac{x}{m} \right\rceil$$

lemma 3: if $a, b \in \mathbb{Z}^+$, $n \in \mathbb{Z}^+$, then

$$(1) \quad \left\lfloor \frac{\lfloor n/a \rfloor}{b} \right\rfloor = \left\lfloor \frac{n}{ab} \right\rfloor$$

$$(2) \quad \left\lceil \frac{\lceil n/a \rceil}{b} \right\rceil = \left\lceil \frac{n}{ab} \right\rceil$$

see handout for proofs.

log

Let $x, a, b \in \mathbb{R}$ where $x > 0, a > 1, b > 1$.

$\log_a(x)$ is inverse of $\exp_a(x) = a^x$.

This means:

$$a^{\log_a x} = x$$

and

$$\log_a(a^x) = x$$

Therefore

$$(*) \quad x = a^{\log_a x} = \left(b^{\log_b(a)} \right)^{\log_a x} = b^{\log_b(a) \cdot \log_a(x)}$$

Take $\log_b()$ of both sides

$$(**) \quad \log_b(x) = \underbrace{\log_b(a)}_{\text{const.}} \cdot \log_a(x)$$

hence

$$\log_b(n) = \text{const} \cdot \log_a(n)$$

$$\therefore \boxed{\log_b(n) = \Theta(\log_a(n))}$$

conversion formula :

$$\log_a(x) = \frac{\log_b(x)}{\log_b(a)} \quad (\text{by } **)$$

also by (**) we have

$$\begin{aligned} a^{\log_b(x)} &= a^{\log_b(a) \cdot \log_a(x)} \\ &= \left(a^{\log_a(x)}\right)^{\log_b(a)} \\ &= x^{\log_b(a)} \end{aligned}$$

$$\therefore \boxed{a^{\log_b(x)} = x^{\log_b(a)}}$$

Stirling's formula

Let $n \in \mathbb{Z}^+$. Then

$$n! = \sqrt{2\pi n} \cdot \left(\frac{n}{e}\right)^n \cdot \left(1 + \mathcal{O}\left(\frac{1}{n}\right)\right)$$

weaker version:

$$n! \sim \sqrt{2\pi n} \cdot \left(\frac{n}{e}\right)^n$$

Corollary

(1) $n! = o(n^n)$ ✓

(2) $n! = \omega(b^n)$ for any $b > 0$

(3) $\log(n!) = \Theta(n \log n)$

Proof of (1)

$$\frac{n!}{n^n} = \frac{\sqrt{2\pi n} \cdot \cancel{n^n} \cdot (1 + \Theta(\frac{1}{n}))}{\cancel{n^n}}$$

$$= \sqrt{2\pi} \cdot \frac{n^{1/2}}{e^n} \cdot (1 + \Theta(\frac{1}{n})) \rightarrow 0 \text{ as } n \rightarrow \infty$$

Exercise: Prove (2)

Proof of (3)

$$\log(n!) = \log\left(\sqrt{2\pi} \cdot n^{1/2} \cdot \frac{n^n}{e^n} \cdot (1 + \Theta(\frac{1}{n}))\right)$$

$$= \log(\sqrt{2\pi}) + \frac{1}{2} \log n + n \log n - n \log e + \log(1 + \Theta(\frac{1}{n}))$$

$$\therefore \frac{\log(n!)}{n \log n} = \frac{\log(\sqrt{2\pi})}{n \log n} + \frac{1}{2n} + \frac{1}{n} - \frac{\log e}{\log n} + \frac{\log(\quad)}{n \log n}$$

$\downarrow \quad \quad \downarrow \quad \downarrow \quad \downarrow \quad \downarrow$
 $0 \quad \quad 0 \quad 1 \quad 0 \quad 0$

$$\therefore \lim_{n \rightarrow \infty} \left(\frac{\log n!}{n \log n} \right) = 1$$

$$\therefore \log n! = \Theta(n \log n).$$

$$\underline{\text{Ex.}} \quad \binom{2n}{n} = \Theta\left(\frac{4^n}{\sqrt{n}}\right)$$

$$\text{Recall: } \binom{m}{k} = \frac{m!}{k! (m-k)!}$$

Proof:

$$\binom{2n}{n} = \frac{(2n)!}{n! (2n-n)!}$$

$$= \frac{(2n)!}{(n!)^2}$$

$$= \frac{\sqrt{2\pi \cdot 2n} \cdot \left(\frac{2^n}{e}\right)^{2n} \cdot \left(1 + O\left(\frac{1}{2n}\right)\right)}{2\pi n \cdot \left(\frac{n}{e}\right)^{2n} \cdot \left(1 + O\left(\frac{1}{n}\right)\right)^2}$$

$$= \frac{1}{\sqrt{\pi}} \cdot \frac{1}{\sqrt{n}} \cdot \frac{2^{2n} \cdot \cancel{n^{2n}}}{\cancel{e^{2n}}} \cdot \frac{\cancel{e^{2n}}}{\cancel{n^{2n}}} \cdot \left[\frac{1 + O\left(\frac{1}{2n}\right)}{\left(1 + O\left(\frac{1}{n}\right)\right)^2} \right]$$

$$= \frac{1}{\sqrt{\pi}} \cdot \frac{4^n}{\sqrt{n}} \cdot \left[\frac{\dots}{\dots} \right]$$

$$\therefore \frac{\binom{2n}{n}}{\frac{4^n}{\sqrt{n}}} = \frac{1}{\sqrt{\pi}} \cdot \left[\frac{\dots}{\dots} \right] \rightarrow \frac{1}{\sqrt{\pi}} \text{ as } n \rightarrow \infty$$

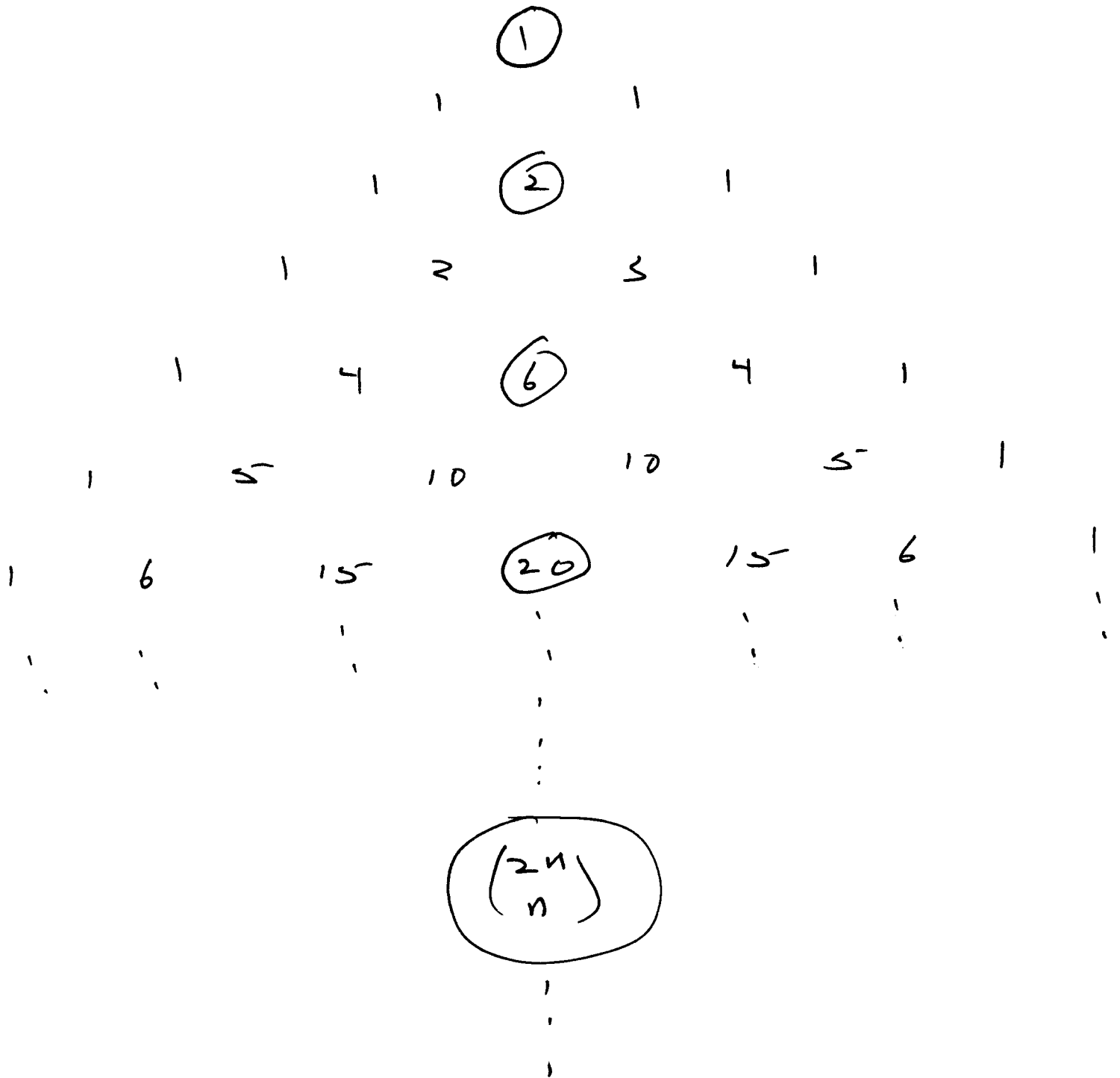
Since $0 < \frac{1}{\sqrt{\pi}} < \infty$, we have

$$\binom{2n}{n} = O\left(\frac{4^n}{\sqrt{n}}\right).$$

■

Pascal's Δ .

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Exercise

determine $a > 0$ such that

$$\binom{3^n}{n} = \Theta\left(\frac{a^n}{\sqrt{n}}\right)$$

Exercise

for each $k \geq 2$, determine $a(k)$ such that

$$\binom{k^n}{n} = \Theta\left(\frac{a(k)^n}{\sqrt{n}}\right)$$

Know : $a(2) = 4$

$a(3) =$ Prev. exercise .

$\vdots$

Induction handout

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Let $P(n)$ be a Propositional Function

$P(n) = \left(\begin{array}{l} \text{a sentence whose truth or falsity} \\ \text{depends on } n \in \mathbb{N} = \{0, 1, 2, \dots\} \end{array} \right)$

in a way

$$P: \mathbb{N} \rightarrow \{\text{false}, \text{true}\}$$

Goal: Prove statements of the form

$$\forall n \geq n_0 : P(n)$$

Domino analogy:



$P(n)$ = "the n^{th} domino falls".

Goal: $\forall n \geq n_0 : P(n)$

= "all dominoes (after n_0) fall".

Induction Proof.

analogy

I. Base step: Prove $P(n_0)$ } "domino n_0 falls"

IIa Induction step:

$\forall n \geq n_0 : P(n) \rightarrow P(n+1)$ } "if any domino falls, then the next also falls"

once I and II_a are Proven,
we may conclude

$$\forall n \geq n_0: P(n)$$

} "all dominos
fall"

continue