

CSE 201 1-7-20

Asymptotic growth of functions:

Defn. let $f(n)$ be a function

we say $f(n)$ is asymptotically non-negative

(a.n.n.) iff

$$\exists n_0 > 0, \forall n \geq n_0 : f(n) \geq 0$$

we say $f(n)$ is asymptotically Positive

(a.p.) iff

$$\exists n_0 > 0, \forall n \geq n_0 : f(n) > 0$$

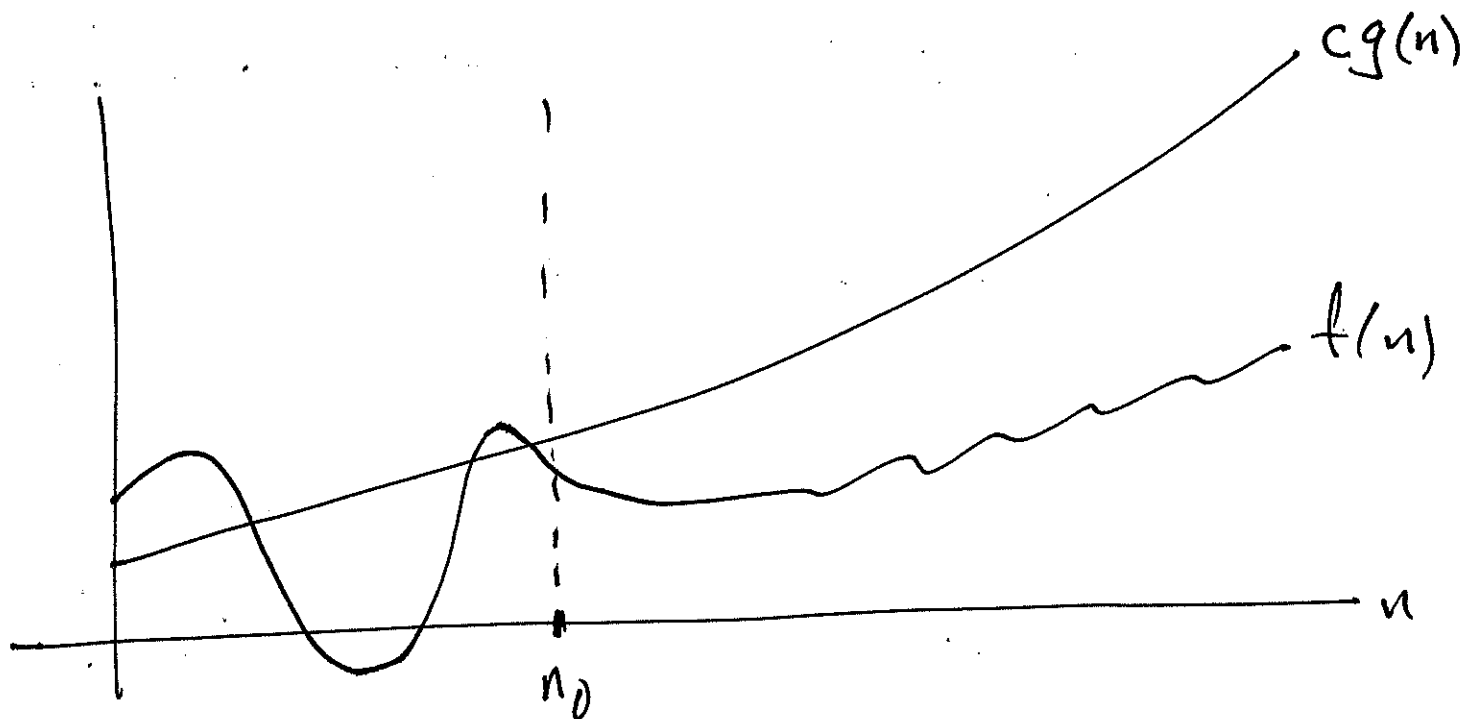
Defn

$$O(g(n)) =$$

$$\left\{ f(n) \mid \exists c > 0, \exists n_0 > 0, \forall n \geq n_0: 0 \leq f(n) \leq c g(n) \right\}$$

write $f(n) = O(g(n))$.

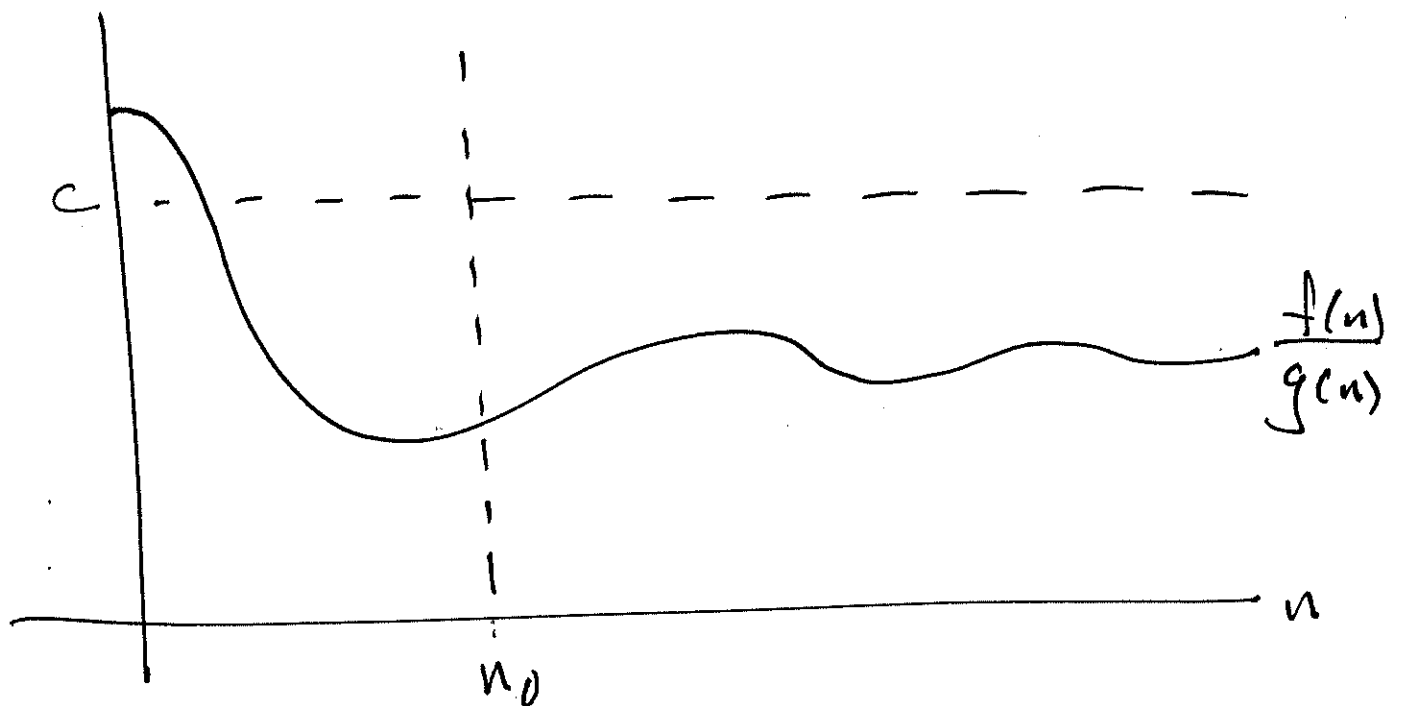
We say " $g(n)$ is an asymptotic upper bound for $f(n)$ ".



Blanket assumption: all fens.
under discussion are a.n.u.

Note if $g(n)$ is a.p., this is
equivalent to

$$\exists c > 0, \exists n_0 > 0, \forall n \geq n_0: 0 \leq \frac{f(n)}{g(n)} \leq c$$



Defn

$$\Omega(g(n)) =$$

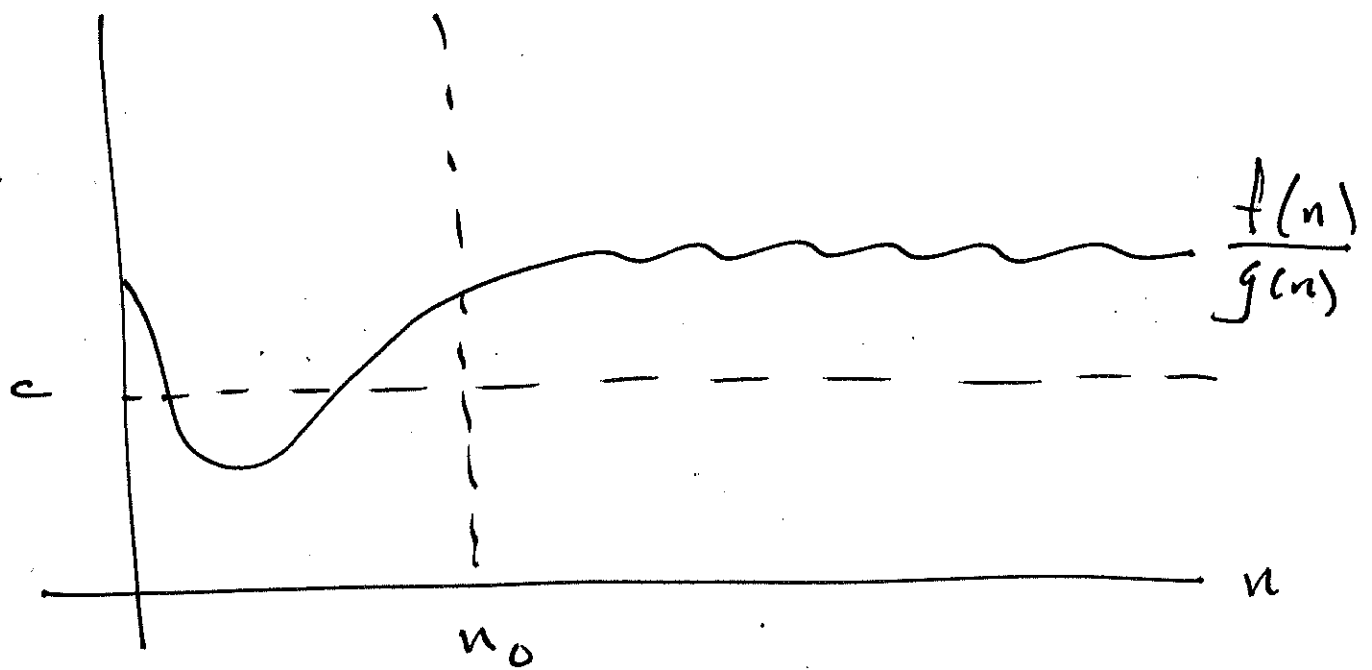
$$\{f(n) \mid \exists c > 0, \exists n_0 > 0, \forall n \geq n_0: 0 \leq c g(n) \leq f(n)\}$$

write $f(n) = \Omega(g(n))$

"asymptotic lower bound
for $f(n)$ "

note: if $g(n)$ is a.p., this is
equiv. to

$$\exists c > 0, \exists n_0 > 0, \forall n \geq n_0: 0 \leq c \leq \frac{f(n)}{g(n)}$$



Ex. If $P(n), q(n)$ are polynomials with

$$\deg(P(n)) \leq \deg(q(n))$$

then

$$P(n) = O(q(n))$$

and

$$q(n) = \Omega(P(n))$$

Theorem

$$f(n) = O(g(n)) \quad \begin{array}{c} \xrightarrow{\text{iff}} \\ \xleftarrow{\text{iff}} \end{array} \quad g(n) = \Omega(f(n))$$

Proof ($\Rightarrow$)

Assume $f(n) = O(g(n))$. Then

$$(1) \exists c_1 > 0, \exists n_1 > 0, \forall n \geq n_1: 0 \leq f(n) \leq c_1 g(n).$$

We must show $g(n) = \Omega(f(n))$, i.e.

$$(2) \exists c_2 > 0, \exists n_2 > 0, \forall n \geq n_2: 0 \leq c_2 f(n) \leq g(n).$$

$$\text{Let } c_2 = \frac{1}{c_1}, \text{ and } n_2 = n_1.$$

Then $c_2 > 0$, $n_2 > 0$ and $(1) \Rightarrow (2)$.



Defn

$$\Theta(g(n)) = O(g(n)) \cap \Omega(g(n))$$

write $f(n) = \Theta(g(n))$

say " $g(n)$ is a tight asymptotic bound on $f(n)$ "

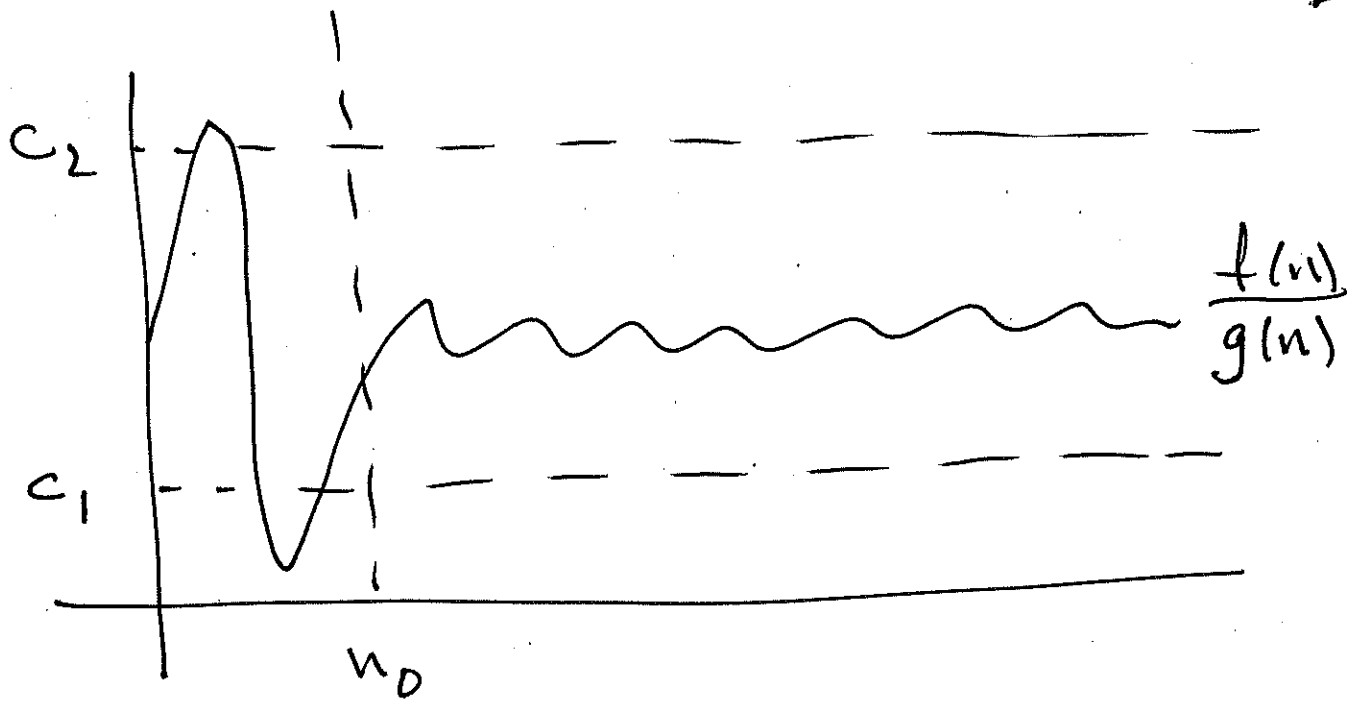
Equivalently: $f(n) = \Theta(g(n))$ iff

$$\exists c_1 > 0, \exists c_2 > 0, \exists n_0 > 0, \forall n \geq n_0:$$

$$0 \leq c_1 g(n) \leq f(n) \leq c_2 g(n)$$

if $g(n)$ is a.p. this is equiv. to

$$0 \leq c_1 \leq \frac{f(n)}{g(n)} \leq c_2$$



Exercise Prove

$$f(n) = \Theta(g(n)) \xrightarrow{\quad} \text{iff } g(n) = \Theta(f(n)) \xleftarrow{\quad}$$

Exercise let $g(n)$ be any function,

let $c > 0$. Then

$$c g(n) = O(g(n))$$

$$c g(n) = \Omega(g(n))$$

$$c g(n) = \Theta(g(n))$$

Analogy:

$$f(n) = O(g(n)) \sim x \leq y$$

$$f(n) = \Omega(g(n)) \sim x \geq y$$

$$f(n) = \Theta(g(n)) \sim x = y$$

$$f(n) = o(g(n)) \sim x < y$$

$$f(n) = \omega(g(n)) \sim y < x$$

Ex Prove $\sqrt{n+10} = \Theta(\sqrt{n})$

Proof

we must find positive c_1, c_2, n_0 st.

for all $n \geq n_0$:

$$0 \leq c_1 \sqrt{n} \leq \sqrt{n+10} \leq c_2 \sqrt{n}$$

Let $c_1 = 1$, $c_2 = \sqrt{2}$, $n_0 = 10$.

Then if $n \geq n_0 = 10$, we have

$$-10 \leq 0 \quad \text{and} \quad 10 \leq n$$

$$\therefore -10 \leq (1-1)n \quad \text{and} \quad 10 \leq (2-1)n$$

$$\therefore -10 \leq (1-c_1^2)n \quad \text{and} \quad 10 \leq (c_2^2-1)n$$

$$\therefore c_1^2 n \leq n+10 \quad \text{and} \quad n+10 \leq c_2^2 n$$

$$\therefore 0 \leq c_1^2 n \leq n+10 \leq c_2^2 n$$

$$\therefore 0 \leq c_1 \sqrt{n} \leq \sqrt{n+10} \leq c_2 \sqrt{n},$$

as required. 

Exercise

show $c_1 = \sqrt{\frac{1}{2}}$, $c_2 = \sqrt{\frac{3}{2}}$, $n_0 = 20$

Also work.

Exercise

let $a, b \in \mathbb{R}$ with $b > 0$. show
that

$$(n+a)^b = \Theta(n^b)$$

Theorem

If $h(n) = O(g(n))$ and $f(n) \leq h(n)$
 for all sufficiently large n ,
 then $f(n) = O(g(n))$.

Proof

Since $h(n) = O(g(n))$ we have

$$(1) \exists c_1 > 0, \exists n_1 > 0, \forall n \geq n_1 : 0 \leq h(n) \leq c_1 g(n).$$

Since $f(n) \leq h(n)$ for suff. large n
 we have

$$(2) \exists n_2 > 0, \forall n \geq n_2 : 0 \leq f(n) \leq h(n)$$

we must show: $f(n) = O(g(n))$, i.e.

(3) $\exists c_3 > 0, \exists n_3 > 0, \forall n \geq n_3$:

$$0 \leq f(n) \leq c_3 g(n)$$

Let $c_3 = c_1$, $n_3 = \max(n_1, n_2)$.

Then $c_3 > 0, n_3 > 0$ and (1) and (2) imply (3). □

Exercise Show

if $h(n) = \Omega(g(n))$ and $f(n) \geq h(n)$ for all suff. large n , then

$$f(n) = \Omega(g(n)).$$