

11

CSE 201 1-16-20

Ex. define

$$T(n) = \begin{cases} 0 & \text{if } n = 1 \\ T(\lfloor \frac{n}{2} \rfloor) + 1 & \text{if } n \geq 2 \end{cases}$$

show: $\boxed{T(n) \leq \lg(n)}$ for all $n \geq 1$. $\nwarrow P(n)$

(hence $T(n) = O(\log n)$.)

Proof.

I. $P(1)$ says

$$T(1) \leq \lg(1) \quad \checkmark$$

$$\text{i.e. } 0 \leq 0 \quad \checkmark$$

$$\text{Ind. } \forall n > 1: \underbrace{P(1) \wedge \dots \wedge P(n-1)}_{\forall k: 1 \leq k < n: P(k)} \rightarrow P(n)$$

- Let $n > 1$ be chosen arbitrarily.
- Assume for any k in the range $1 \leq k < n$ that $T(k) \leq \lg(k)$.
- We must show $T(n) \leq \lg(n)$.

Then

$$T(n) = T(\lfloor \frac{n}{2} \rfloor) + 1$$

$$\leq \lg(\lfloor \frac{n}{2} \rfloor) + 1$$

$$\leq \lg(\frac{n}{2}) + 1$$

by the recurrence
for $T(n)$

• By the ind. hyp.
with $k = \lfloor \frac{n}{2} \rfloor$

Since $\lfloor x \rfloor \leq x$
and since $\lg(\cdot)$
is increasing

$$= \lg(n) - \cancel{1} + \cancel{1} \dots$$

$$= \lg(n)$$

Hence $T(n) \leq \lg(n)$ for all $n \geq 1$

by 2nd P.M.I. ■

Multiple Base Cases:

I. show $P(1), P(2), \dots, P(n_0)$

IId. $\forall n > n_0: P(1) \wedge \dots \wedge P(n-1) \rightarrow P(n)$

Conclusion: $\forall n \geq 1: P(n)$

Ex. define

$$T(n) = \begin{cases} 1 & \text{if } n=1, n=2 \\ 4T(\lfloor \frac{n}{3} \rfloor) + n & \text{if } 3 \leq n \end{cases}$$

show $\forall n \geq 1 : T(n) \leq n^2$ $\leftarrow P(n)$
 (hence $T(n) = O(n^2)$.)

Proof..

I. $P(1)$ says: $T(1) \leq 1^2$

$$1 \leq 1 \quad \checkmark$$

$P(2)$ says: $T(2) \leq 2^2$

$$1 \leq 4 \quad \checkmark$$

Ind. $\forall n \geq 2: T(1) \dots T(n-1) \rightarrow T(n)$

- Let $n \geq 2$ be arbitrary.
- Assume for all k in range $1 \leq k < n$ that $T(k) \leq k^2$.

• we must show that: $T(n) \leq n^2$

Then

$$T(n) = 4T\left(\left\lfloor \frac{n}{3} \right\rfloor\right) + n \quad \text{by definition of } T(n).$$

$$\leq 4 \cdot \left(\left\lfloor \frac{n}{3} \right\rfloor\right)^2 + n$$

$$\leq 4 \left(\frac{n}{3}\right)^2 + n$$

$$\leq n^2$$

• by ind. hyp.
using $k = \left\lfloor \frac{n}{3} \right\rfloor$
note: $1 \leq \left\lfloor \frac{n}{3} \right\rfloor < n$

Since $\lfloor x \rfloor \leq x$
and $(\)^2$ is
increasing.

note $4\left(\frac{n}{3}\right)^2 + n \leq n^2$

pt. $n > 2 \Rightarrow \frac{9}{5} \leq n$

$$\Rightarrow n \leq \frac{5}{9} n^2$$

$$\Rightarrow \frac{4}{9} n^2 + n \leq n^2$$

$$\Rightarrow 4\left(\frac{n}{3}\right)^2 + n \leq n^2$$

Hence $T(n) \leq n^2$ for all $n \geq 1$

By 2nd PMI.

~~QED~~

why 2 base cases?

n	1	2	3	4	5	6	7	8	9	10
$\lfloor \frac{n}{2} \rfloor$	.	.	1	1	1	2	2	2		

note: $T(n)$ depends on

$$T\left(\left\lfloor \frac{n}{2^{\lfloor \log_2 n \rfloor}} \right\rfloor\right)$$

Induction Fallacies

Example C:

All Horses are of the same color.

Prove

$P(n)$

$\forall n \geq 1$: If S is a set of n horses, then all horses in S are of same color.

I. $P(1)$ says: "if S is a set of 1 horse, then all horses in S have same color".

IIa. $\forall n > 1 : P(n) \rightarrow P(n+1)$

- let $n > 1$ be arbitrary, hence $n \geq 2$.
- Assume that if S is a set of n horses, all of them have same color.

• We must show: If S is a set of $n+1$ horses, then all of them have same color.

Let S be a set of $n+1$ horses, say

$$S = \{h_1, h_2, h_3, \dots, h_{n+1}\}.$$

consider the sets

$$S' = S - \{h_1\}$$

$$S'' = S - \{h_2\}$$

~~the~~ each contain n horses, so

• by the incl. hyp., all in S' have same color, and all in S'' have same color.

note $h_3 \in S' \cap S''$

observe $n \geq 2 \Rightarrow n+1 \geq 3$ so

h_3 exists. Since h_3 has only one color, the horses in S' have same color as horses in S'' . Therefore, all horses in $S' \cup S'' = S$ have the same color. 

Justification of Induction

!!!

Theorem 1: 1st PMI.

For any Propositional function $P(n)$,
the following sentence is true

$$\left[P(1) \wedge (\forall n > 1: P(n-1) \rightarrow P(n)) \right] \rightarrow (\forall n: P(n))$$

Theorem 2: 2nd PMI

for any Prop.fcn $P(n)$, the following
sentence is true

$$\left[P(1) \wedge (\forall n > 1: (\forall k < n: P(k)) \rightarrow P(n)) \right] \rightarrow (\forall n: P(n))$$

Theorem 3 (well ordering prop. of $\mathbb{Z}^+$: WOP)

Any non-empty subset of $\mathbb{Z}^+$ contains a least element.

Proof that $\text{WOP} \Rightarrow \text{IPNI}$

Assume both

$$(1) \ P(1)$$

and

$$(2) \ \forall n > 1 : P(n-1) \rightarrow P(n)$$

are true. Let

$$S = \{n \in \mathbb{Z}^+ \mid P(n) \text{ is false}\}$$

It is sufficient to show
 $S = \emptyset$, since then we have

$$\forall n : P(n)$$

is true.

Assume, to get a contradiction,
 that $S \neq \emptyset$. By the W.O.P.

S contains a least element,
 call it m . so

- $m \in S$

and

- if $k < m$, then $k \notin S$

since $m \in S$, $P(m)$ is false.

But $P(1)$ is true, so $m \neq 1$.

Hence $m > 1$, and $m-1 \geq 1$.

so $P(m-1)$ is defined.

Now $m-1 < m$, so $m-1 \notin S$,

therefore $\boxed{P(m-1)}$ is true.

By (2) and Univ. Inst. we have that $\boxed{P(m-1) \rightarrow P(m)}$ is true.

so By Modus Ponens $P(m)$ is

true. It follows that

$m \notin S$. This contradicts

That $m \in S$. This $\times$ shows
our assumption was false.

Hence $S = \emptyset$ and

$\forall n: P(n)$

is true.

