

CS2 16 7-6-22

Ex. list all $P(4, 3) = 4 \cdot 3 \cdot 2 = 24$
3-permutations of $\{1, 2, 3, 4\}$.

1 2 3	2 1 3	3 1 2	4 1 2
1 3 2	2 3 1	3 2 1	4 2 1
1 2 4	2 1 4	3 1 4	4 1 3
1 4 2	2 4 1	3 4 1	4 3 1
1 3 4	2 3 4	3 2 4	4 2 3
1 4 3	2 4 3	3 4 2	4 3 2

Ex. (3.4 #7)

find # of 9-digit numbers satisfying

- no zeros : $\{1, 2, 3, 4, 5, 6, 7, 8, 9\}$
- no repetition
- all odd digits are to the left of all even digits.

odd digits = $\{1, 3, 5, 7, 9\}$

even digits = $\{2, 4, 6, 8\}$

(1) form an odd digit string. of
len. 5 : $5!$ ways

(2) form an even digit string of
len. 4 : $4!$

Then concatenate strings

$$\# \text{ strings} = 5! \cdot 4! = (120)(24)$$

$$= \boxed{2,880}$$

Exercise:

Same Problem, but 8-digit, 7-digit
and 6-digit strings.

(3.5) Combinations (Count subsets)

Defn

A k -combination of a finite set S is a k -element subset of S .
i.e. an unordered arrangement of k of its elements.

Notation

if $|S| = n$, the # of k -comb. of S is denoted

$$C(n, k) \quad \text{or} \quad \binom{n}{k}$$

read : "n-choose-k"

Ex $S = \{1, 2, 3\}$

	$C(3, k)$
$k=3 : \{1, 2, 3\}$	1
$k=2 : \{1, 2\}, \{1, 3\}, \{2, 3\}$	3
$k=1 : \{1\}, \{2\}, \{3\}$	3
$k=0 : \emptyset$	1

Theorem

If $0 \leq k \leq n$ then

$$\binom{n}{k} = C(n, k) = \frac{n!}{k! (n-k)!}$$

Proof.

compute $P(n, k)$ in a different way.

recall

$$P(n, k) = n(n-1) \cdots (n-k+1) = \frac{n!}{(n-k)!}$$

Let S be an n -element set.

new Process

(1) choose a k -element subset of S

$$\# \text{ways} = C(n, k)$$

(2) choose a permutation of these k elements : $\# \text{ways} = k!$

By the mult. Principle

$$P(n, k) = C(n, k) \cdot k!$$

$$\therefore \frac{n!}{(n-k)!} = C(n, k) \cdot k!$$

$$\therefore C(n, k) = \frac{n!}{k! (n-k)!}$$



Ex.

how many bit strings of len. 10
contain exactly 6 ones.

Each bit string of len. 10 encodes
with 6 1's encodes a 6-element
subset of $\{1, 2, 3, 4, 5, 6, 7, 8, 9, 10\}$.

for instance

$$\begin{array}{cccccccccc} 1 & 0 & 1 & 0 & 1 & 1 & 0 & 1 & 1 & 0 \\ \hline 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 10 \end{array}$$

corresponds to

$$\{1, 3, 5, 6, 8, 9\}$$

$\therefore (\# \text{ such strings}) = (\# \text{ 6-element sets of a } 10\text{-element set})$

$$= \binom{10}{6} = \frac{10!}{6! 4!} = \frac{10 \cdot \overset{3}{\cancel{9}} \cdot \cancel{8} \cdot 7}{\cancel{4} \cdot \cancel{3} \cdot 2 \cdot 1} = \boxed{210}$$

7

note: a bit str. of len. with
6 1's must 4 0's. Thus

(# bit str. of len 10 with 4 0's)

$$= 210$$

Also, choosing a 6-set of a 10-set
is equivalent to choosing a 4-set
of a 10-set (its complement)

Thus

$$\binom{10}{6} = \binom{10}{4}$$

↑
choose which
6 positions will
be occupied by 1's

↑
choose which
4 positions will
be occupied
by 0's.

in fact both are: $\frac{10!}{6!4!}$

Theorem

If $0 \leq k \leq n$, then

$$\binom{n}{k} = \binom{n}{n-k}$$

Algebraic Proof:

$$\begin{aligned} \binom{n}{k} &= \frac{n!}{k! (n-k)!} \\ &= \frac{n!}{(n-k)! \cdot (n-(n-k))!} \\ &= \binom{n}{n-k} \quad \square \end{aligned}$$

Combinatorial Proof:

The LHS counts the number of length n bit strings with exactly k 1's. The RHS counts the number of length n bit strings with exactly $(n-k)$ 0's.

But these 2 sets of bit strings are one and the same.

$$\therefore LHS = RHS, \text{ i.e.}$$

$$\binom{n}{k} = \binom{n}{n-k}.$$

Correspondence between sets & bitstrings

$$S = \{1, 2, 3\}$$

k	subset	bit string	#1's
0	$[\emptyset]$	000	0
1	$\{1\}$	100	1
	$\{2\}$	010	
	$\{3\}$	001	
2	$\{1, 2\}$	110	2
	$\{1, 3\}$	101	
	$\{2, 3\}$	011	
3	$\{1, 2, 3\}$	111	3

Exercise: write correspondence for $n=4$.

(2.6) Binomial Theorem

using : $\binom{n}{k} = \frac{n!}{k!(n-k)!} \quad (0 \leq k \leq n)$

we compute

n

0 $\binom{0}{0} = 1$

1 $\binom{1}{0} = 1, \binom{1}{1} = 1$

2 $\binom{2}{0} = 1, \binom{2}{1} = 2, \binom{2}{2} = 1$

3 $\binom{3}{0} = 1, \binom{3}{1} = 3, \binom{3}{2} = 3, \binom{3}{3} = 1$

4 $\binom{4}{0} = 1, \binom{4}{1} = 4, \binom{4}{2} = 6, \binom{4}{3} = 4, \binom{4}{4} = 1$

5 $\binom{5}{0} = 1, \binom{5}{1} = 5, \binom{5}{2} = 10, \binom{5}{3} = 10, \binom{5}{4} = 5, \binom{5}{5} = 1$

⋮

Theorem (Pascal's identity)

for $1 \leq k \leq n$:

$$(*) \quad \boxed{\binom{n+1}{k} = \binom{n}{k-1} + \binom{n}{k}}$$

equivalently, for $0 < k < n$:

$$(**) \quad \boxed{\binom{n}{k} = \binom{n-1}{k-1} + \binom{n-1}{k}}$$

Algebraic Proof of (*)

$$\text{RHS} = \binom{n}{k-1} + \binom{n}{k}$$

$$= \frac{n!}{(k-1)! (n-(k-1))!} + \frac{n!}{k! (n-k)!}$$

$$= \frac{k \cdot n!}{k(k-1)! (n-k+1)!} + \frac{(n-k+1) \cdot n!}{k! (n-k+1) \cdot (n-k)!}$$

$$= \frac{k \cdot n!}{k! (n-k+1)!} + \frac{(n-k+1) \cdot n!}{k! (n-k+1)!}$$

$$= \frac{(\cancel{k} + n - \cancel{k} + 1) n!}{k! (n-k+1)!}$$

$$= \frac{(n+1)!}{k! ((n+1)-k)!}$$

$$= \binom{n+1}{k}$$

$$= \text{LHS}$$



Combinatorial Proof (**)

Let $S = \{ \text{len. } n \text{ bit strings with } k \text{ 1's} \}$

Let $A = \{ x \in S \mid x \text{ ends with } 1 \}$

$B = \{ x \in S \mid x \text{ ends with } 0 \}$

clearly $S = A \cup B$ and $A \cap B = \emptyset$,

so by addition principle:

$$|S| = |A| + |B|$$

obviously $|S| = \binom{n}{k}$.

• each $x \in A$ can be obtained by choosing a bit str. of len. $(n-1)$ with $(k-1)$ 1's, and appending 1

$$\# \text{ ways} = \binom{n-1}{k-1}$$

$$\therefore |A| = \binom{n-1}{k-1}$$

- each $x \in R$ can be obtained by choosing a bit str. of len. $(n-1)$ containing k 1's, and appending 0:

$$\# \text{ways} = \binom{n-1}{k}$$

$$\therefore |R| = \binom{n-1}{k}$$

Therefore

$$\binom{n}{k} = |S| = |A| + |R|$$

$$= \binom{n-1}{k-1} + \binom{n-1}{k}$$

