

Lemma 1

let $p_1, p_2, \dots, p_n$ be primes, and let

$$a = p_1^{x_1} p_2^{x_2} \dots p_n^{x_n}$$

and

$$b = p_1^{y_1} p_2^{y_2} \dots p_n^{y_n}$$

for some integers $x_i \geq 0, y_i \geq 0$ ($1 \leq i \leq n$).

Then $a \mid b \iff x_i \leq y_i$ ($1 \leq i \leq n$)

Proof.

observe

$$b = (p_1^{x_1} p_2^{x_2} \dots p_n^{x_n}) \cdot \overbrace{(p_1^{y_1-x_1} p_2^{y_2-x_2} \dots p_n^{y_n-x_n})}^K$$

$$= a \cdot K$$

($\Leftarrow$) If $x_i \leq y_i$ then $y_i - x_i \geq 0$ ($1 \leq i \leq n$),
so is an integer. $\therefore a \mid b$.

($\Rightarrow$) If $a|b$, then $b = ak'$ for some

$k' \in \mathbb{Z}$. Then $ak = ak'$, $\therefore k = k'$

$\therefore k$ is an integer

$\therefore x_i \leq y_i$ for $1 \leq i \leq n$. ■

Lemma 2

Let $a, b \in \mathbb{Z}^+$ and let

$$S = \{p_1, p_2, \dots, p_n\}$$

be the set of primes that divide a , or b , or both. i.e.

$$a = p_1^{x_1} p_2^{x_2} \dots p_n^{x_n}$$

and

$$b = p_1^{y_1} p_2^{y_2} \dots p_n^{y_n}$$

for some $x_i \geq 0, y_i \geq 0$.

Then

$$(1) \gcd(a, b) = p_1^{\min(x_1, y_1)} p_2^{\min(x_2, y_2)} \dots p_n^{\min(x_n, y_n)}$$

and

$$(2) \text{lcm}(a, b) = p_1^{\max(x_1, y_1)} p_2^{\max(x_2, y_2)} \dots p_n^{\max(x_n, y_n)}$$

Illustrate with : $a=30, b=75$

$$30 = 2 \cdot 3 \cdot 5 = 2^1 \cdot 3^1 \cdot 5^1$$

$$75 = 3 \cdot 5 \cdot 5 = 2^0 \cdot 3^1 \cdot 5^2$$

so

$$\gcd(30, 75) = 2^0 \cdot 3^1 \cdot 5^1 = 15$$

$$\text{lcm}(30, 75) = 2^1 \cdot 3^1 \cdot 5^2 = 150$$

Proof of lemma 2

let

$$d = \text{RHS of (1)}$$

and

$$m = \text{RHS of (2)}$$

we must show d, m satisfy definitions of gcd & lcm resp. i.e.

gcd (i) $d|a$ and $d|b$ ✓

(ii) if $e|a$ and $e|b$, then $e \leq d$

lcm (i) $a|m$ and $b|m$

(ii) if $a|k$ and $b|k$, then $m \leq k$

we have $\min(x_i, y_i) \leq x_i$ for $1 \leq i \leq n$. So by lemma 1, we have $d|a$. also $\min(x_i, y_i) \leq y_i$ for $1 \leq i \leq n$, so by lemma 1, we have $d|b$.

Now suppose $e|a$ and $e|b$. by unique factorization:

$$e = p_1^{z_1} p_2^{z_2} \cdots p_n^{z_n}$$

where $z_i \geq 0$ ($1 \leq i \leq n$). by lemma 1

$$\begin{cases} e|a \Rightarrow z_i \leq x_i \\ e|b \Rightarrow z_i \leq y_i \end{cases} \quad \left\{ \begin{array}{l} 1 \leq i \leq n \end{array} \right.$$

$\therefore z_i \leq \min(x_i, y_i)$ for $1 \leq i \leq n$.

Again by lemma 1, we have

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$$e \mid d.$$

It follows that $d = \gcd(a, b)$,

Proving (1). □

Exercise.

Prove part (2), i.e. m satisfies
defn. of $\text{lcm}(a, b)$.

Recall:

Thm

$$(*) \quad a \cdot b = \gcd(a, b) \cdot \text{lcm}(a, b)$$

P-root

□

let $p_1, \dots, p_n$ be as in previous

P-root, so

$$a = p_1^{x_1} p_2^{x_2} \dots p_n^{x_n}$$

$$b = p_1^{y_1} p_2^{y_2} \dots p_n^{y_n}$$

for $x_i \geq 0, y_i \geq 0$ and $1 \leq i \leq n$

Referring to (*)

$$\text{LHS} = a \cdot b = p_1^{x_1+y_1} p_2^{x_2+y_2} \dots p_n^{x_n+y_n}$$

$$\text{RHS} = p_1^{\min(x_1, y_1) + \max(x_1, y_1)} \dots p_n^{\min(x_n, y_n) + \max(x_n, y_n)}$$

using lemma 2. But

$$x_i + y_i = \min(x_i, y_i) + \max(x_i, y_i)$$

for $1 \leq i \leq n$. $\therefore$ LHS = RHS,

as required. ■

Ex. $a = 16,200$

$b = 756$

$\gcd(a, b) = 108$

$\text{lcm}(a, b) = \frac{(16200)(756)}{108} = \boxed{113,400}$

Euclid's Algorithm for $\gcd(a, b)$

Ex. $\gcd(16\,200, 756) = 108$

$16\,200 = 21 \cdot 756 + 324$

$756 = 2 \cdot 324 + \boxed{108}$

$324 = 3 \cdot \boxed{108} + 0$

↑
stop

Pseudo-code:

```

r = remainder of a on div. by b
while r > 0
    a = b
    b = r
    r = rem. of a on div. by b
end while
return b
    
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Ex. $\gcd(136, 235) = 1$

$$\begin{array}{rcl}
 \underline{a} & & \underline{b} & & \underline{r} \\
 136 & = & 0 \cdot 235 & + & 136 \\
 \swarrow & & \swarrow & & \swarrow \\
 235 & = & 1 \cdot 136 & + & 99 \\
 \swarrow & & \swarrow & & \swarrow \\
 136 & = & 1 \cdot 99 & + & 37 \\
 \swarrow & & \swarrow & & \swarrow \\
 99 & = & 2 \cdot 37 & + & 25 \\
 \swarrow & & \swarrow & & \swarrow \\
 & & & &
 \end{array}$$

$$\begin{array}{rcl}
 37 & = & 1 \cdot 25 + 12 \\
 25 & = & 2 \cdot 12 + 1 \\
 12 & = & 12 \cdot \boxed{1} + 0 \leftarrow \text{stop}
 \end{array}$$

$\uparrow$
gcd

Defn

we say $a, b \in \mathbb{Z}$ are relatively prime iff $\gcd(a, b) = 1$.

Equivalently, a and b have no prime factors in common.

Goal : Prove correctness of Euclid's Algorithm.

□

Exercise (hw ?)

Prove that: if $d|a$ and $d|b$,
then $d|(a+b)$.

Lemma 3

Let $a, d, n \in \mathbb{Z}$. Then

$$d|a \Rightarrow d|na.$$

Proof.

$$d|a \Rightarrow a = kd \text{ for some } k \in \mathbb{Z}.$$

$$\Rightarrow na = n(kd) = (nk)d$$

$$\Rightarrow d|na \text{ since } nk \in \mathbb{Z}.$$

■

Lemma 4

Let $a, b, n, m, d \in \mathbb{Z}$. Then

$$d|a \wedge d|b \Rightarrow d|(na+mb)$$

Proof:

$$\left. \begin{array}{l} d|a \Rightarrow d|na \\ d|b \Rightarrow d|mb \end{array} \right\} \Rightarrow d|(na+mb)$$

$\uparrow$ by lemma 3 $\uparrow$ exercise

Exercise

- Prove lemma 4 directly from defn. of divisibility, i.e. don't use exercise and lemma 3.
- Prove lemma 3 & exercise from lemma 4

lemma 5

Suppose $a = qb + r$ (for some $q \in \mathbb{Z}$).

Then

$$\gcd(a, b) = \gcd(b, r)$$

Proof: let $d \in \mathbb{Z}$. Then

$$d|a \wedge d|b \Rightarrow d|(a - qb) = r \Rightarrow d|b \wedge d|r$$

$\uparrow$
 lemma 4
 $n=1, m=-q$

Also

$$d|b \wedge d|r \Rightarrow d|(qb + r) = a \Rightarrow d|a \wedge d|b$$

$\uparrow$
 lemma 4
 $a=b$
 $b=r$
 $n=q$
 $m=1$

Thus

$$\{ \text{com. div of } a, b \} = \{ \text{com. div. of } b, r \}$$

$$\therefore \gcd(a, b) = \gcd(b, r)$$



Proof of correctness of Euclid's Alg.

The 1st division gives

$$a = q_1 b + r \quad (0 \leq r < b)$$

$$\text{Let } r_0 = a, q_1 = q, r_1 = b, r_2 = r \quad . \geq 0$$

$$r_0 = q_1 r_1 + r_2 \quad (0 \leq r_2 < r_1)$$

$$r_1 = q_2 r_2 + r_3 \quad (0 \leq r_3 < r_2)$$

⋮

The seq. of remainders is strictly decreasing

$$r_1 > r_2 > r_3 > r_4 > \dots \geq 0,$$

and therefore reach 0. let r_n be the last non-zero remainder.

⋮

$$r_{n-2} = q_{n-1} r_{n-1} + \boxed{r_n} \quad (0 \leq r_n < r_{n-1})$$

$$r_{n-1} = q_n r_n + 0$$

By lemma 5:

$$\begin{aligned} \gcd(a, b) &= \gcd(r_0, r_1) \\ &= \gcd(r_1, r_2) \end{aligned}$$

$$\begin{aligned}
 & \vdots \\
 &= \gcd(r_{n-1}, r_n) \\
 &= \gcd(r_n, 0) \\
 &= r_n
 \end{aligned}$$

Defn

Let $a, b, m \in \mathbb{Z}$ and $m > 0$. We say that a and b are congruent modulo m iff

$$m \mid (a - b)$$

notation : $a \equiv b \pmod{m}$

Alternate notation : $a \equiv_m b$.