

Ex. Prove $\sqrt{n+10} = \Theta(\sqrt{n})$

Proof

we must show there exist positive c_1, c_2, n_0 s.t. $\sqrt{n} \geq n_0$:

$$0 \leq c_1 \sqrt{n} \leq \sqrt{n+10} \leq c_2 \sqrt{n}$$

Let $c_1 = 1$, $c_2 = \sqrt{2}$, $n_0 = 10$.

Then if $n \geq n_0 = 10$, we have

$$-10 \leq 0 \quad \text{and} \quad 10 \leq n$$

$$\circ \circ \quad -10 \leq (1-1)n \quad \text{and} \quad 10 \leq (2-1)n$$

$$\circ \circ \quad -10 \leq (1-c_1^2)n \quad \text{and} \quad 10 \leq (c_2^2-1)n$$

∴

$$c_1^2 n \leq n+10 \quad \text{and} \quad n+10 \leq c_2^2 n$$

$$\therefore 0 \leq c_1^2 n \leq n+10 \leq c_2^2 n$$

$$\therefore 0 \leq c_1 \sqrt{n} \leq \sqrt{n+10} \leq c_2 \sqrt{n} \quad \blacksquare$$

Exercise

find other constants c_1, c_2, n_0 that work.

Exercise

show that $c_1 = \sqrt{\frac{1}{2}}, c_2 = \sqrt{\frac{3}{2}}, n_0 = 20$

Also work.

Exercise

Let $a, b \in \mathbb{R}$ where $b > 0$.

show that

$$(n+a)^b = \Theta(n^b)$$

Theorem

If $h(n) = O(g(n))$ and $f(n) \leq h(n)$

for all suff. large n , then

$$f(n) = O(g(n))$$

Proof.

Since $h(n) = O(g(n))$, we have
positive c_1, n_1 st.

$$(1) \quad \forall n \geq n_1 : 0 \leq h(n) \leq c_1 g(n)$$

since $f(n) \leq h(n)$ for suff. large
 n , we have pos. n_2 st.

$$(2) \quad \forall n \geq n_2 : 0 \leq f(n) \leq h(n)$$

we must find positive c_3, n_3 st.

$$(3) \quad \forall n \geq n_3 : 0 \leq f(n) \leq c_3 g(n).$$

Let $c_3 = c_1$, and $n_3 = \max(n_1, n_2)$.

Thus c_3, n_3 are positive and

(1) and (2) \Rightarrow (3). \square

Exercise

If $h(n) = \Omega(g(n))$ and

$f(n) \geq h(n)$ for all suff. large n ,

then

$$f(n) = \Omega(g(n))$$

Exercise

If $h_2(n) = O(g(n))$, $h_1(n) = \Omega(g(n))$

and if

$$h_1(n) \leq f(n) \leq h_2(n)$$

for all suff. large n , then

$$f(n) = \Theta(g(n)).$$

Ex.

Let $k > 0$. Prove that

$$\sum_{i=1}^n i^k = \Theta(n^{k+1})$$

Remark:

$$\sum_{i=1}^n i = \frac{n(n+1)}{2}$$

$$\sum_{i=1}^n i^2 = \frac{n(n+1)(2n+1)}{6}$$

$$\sum_{i=1}^n i^3 = \left(\frac{n(n+1)}{2} \right)^2$$

⋮

Proof

$$\sum_{i=1}^n i^k \leq \sum_{i=1}^n n^k = n \cdot n^k = n^{k+1} = O(n^{k+1})$$

Also

$$\sum_{i=1}^n i^k \approx \sum_{i=\lceil \frac{n}{2} \rceil}^n i^k$$

$$\approx \sum_{i=\lceil \frac{n}{2} \rceil}^n \lceil \frac{n}{2} \rceil^k$$

$$= (n - \lceil \frac{n}{2} \rceil + 1) \lceil \frac{n}{2} \rceil^k$$

$$= (\lfloor \frac{n}{2} \rfloor + 1) \lceil \frac{n}{2} \rceil^k$$

$$\approx \left(\frac{n}{2} - 1 + 1\right) \left(\frac{n}{2}\right)^k$$

$\left\{ \begin{array}{l} \lfloor x \rfloor > x - 1 \\ \lceil x \rceil \geq x \end{array} \right.$

$$\approx \left(\frac{n}{2}\right)^{k+1}$$

$$= \left(\frac{1}{2}\right)^{k+1} \cdot n^{k+1} = \Omega\left(n^{k+1}\right)$$

hence $\sum_{i=1}^n i^k = \Theta\left(n^{k+1}\right)$



Defn

$o(g(n))$

$$= \{ f(n) \mid \forall c > 0, \exists n_0 > 0, \forall n \geq n_0 : 0 \leq f(n) < c g(n) \}$$

write: $f(n) = o(g(n))$

say: " $g(n)$ is a strict asyml. U. b.
for $f(n)$ "

Recall:

$O(g(n)) =$

$$\{ f(n) \mid \exists c > 0, \exists n_0 > 0, \forall n \geq n_0 : 0 \leq f(n) \leq c g(n) \}$$

Obviously:

$$o(g(n)) \subseteq O(g(n))$$

Lemma 1

$$f(n) = o(g(n)) \text{ iff } \lim_{n \rightarrow \infty} \left(\frac{f(n)}{g(n)} \right) = 0$$

Prove:

$$f(n) = o(g(n)) \text{ iff}$$

$$\forall c > 0, \exists n_0 > 0, \forall n \geq n_0: 0 \leq f(n) < c g(n)$$

$$\therefore \forall c > 0, \exists n_0 > 0, \forall n \geq n_0: 0 \leq \frac{f(n)}{g(n)} < c$$

This is the definition of

$$\lim_{n \rightarrow \infty} \left(\frac{f(n)}{g(n)} \right) = 0$$



Ex. $\ln(n) = o(n)$

$$\lim_{n \rightarrow \infty} \left(\frac{\ln(n)}{n} \right) = \lim_{n \rightarrow \infty} \frac{1/n}{1} = 0$$

↑
L'Hop.

Exercise show $\log_b(n) = o(n)$ for any $b > 1$.

Ex. let $k > 0$, show $n^k = o(e^n)$

$$\begin{aligned} \lim_{n \rightarrow \infty} \left(\frac{n^k}{e^n} \right) &= \lim_{n \rightarrow \infty} \left(\frac{k \cdot n^{k-1}}{e^n} \right) \\ &= \lim_{n \rightarrow \infty} \left(\frac{k \cdot (k-1) n^{k-2}}{e^n} \right) \\ &\vdots \end{aligned}$$

= 0

after Γ KT app.
of L'Hop.

Exercise

show $n^k = o(b^n)$ for any

$b > 1$.

Exercise show

$$o(g(n)) \cap \Omega(g(n)) = \emptyset$$

Defn

$$w(g(n)) =$$

$$\{f(n) \mid \forall c > 0, \exists n_0 > 0, \forall n \geq n_0: 0 \leq c g(n) \leq f(n)\}$$

check: $w(g(n)) \subseteq \Omega(g(n))$

Exercise:

- $f(n) = o(g(n))$ iff $\lim_{n \rightarrow \infty} \left(\frac{f(n)}{g(n)} \right) = 0$
- $\omega(g(n)) \cap O(g(n)) = \emptyset$

Picture:

