

cse

102

1-16-20

1

Recall Stirling's formula

$$n! = \sqrt{2\pi n} \cdot \left(\frac{n}{e}\right)^n \cdot \left(1 + \Theta\left(\frac{1}{n}\right)\right)$$

Corollary

(1)  $n! = o(n^n)$  ✓

(2)  $n! = \omega(b^n)$  any  $b > 0$  (Hint:  $b^n = e^{n \cdot \ln b}$ )

(3)  $\log(n!) = \Theta(n \log n)$  ✓

Proof of (1)

$$\frac{n!}{n^n} = \frac{\sqrt{2\pi n} \cdot \frac{n^4}{e^n} \cdot (1 + \Theta(\frac{1}{n}))}{n^4}$$

$$= \frac{\sqrt{2\pi} \cdot n^{1/2} \cdot (1 + \Theta(\frac{1}{n}))}{e^n} \rightarrow 0 \text{ as } n \rightarrow \infty$$

Exercise: Prove (2)

Proof of (2)

$$\log(n!) = \log \sqrt{2\pi} + \log \sqrt{n} + \log n^n - \log e^n + \log(\quad)$$

$$= \log \sqrt{2\pi} + \frac{1}{2} \log n + n \log n - n \log e + \log(\quad)$$

> 0

$$\frac{\log(n!)}{n \log n} = \frac{\log \sqrt{2\pi}}{n \log n} + \frac{1}{2n} + 1 - \frac{\log e}{\log n} + \frac{\log(\dots)}{n \log n}$$

$\downarrow$        $\downarrow$        $\downarrow$        $\downarrow$        $\downarrow$   
 0      0      1      0      0

1 as  $n \rightarrow \infty$

Ex  $\binom{2n}{n} = \Theta\left(\frac{4^n}{\sqrt{n}}\right)$  ✓

Recall  $\binom{m}{k} = \frac{m!}{k!(m-k)!}$

$\binom{m}{k}$  defn. # of  $k$ -subsets of an  $m$ -set,  $(0 \leq k \leq m)$

# Pascal's $\Delta$

0				①			
1			1		1		
2			1	②	1		
3		1	3	3	1		
4		1	4	⑥	4	1	
5	1	5	10	10	5	1	
6	1	6	15	②0	15	6	1
⋮				⋮			

## Proof

$$\begin{aligned}
 \binom{2n}{n} &= \frac{(2n)!}{n!(2n-n)!} \\
 &= \frac{(2n)!}{(n!)^2}
 \end{aligned}$$

$$\sqrt{2\pi \cdot 2n} \cdot \left(\frac{2n}{e}\right)^{2n} \cdot \left(1 + O\left(\frac{1}{2n}\right)\right)$$

$$\frac{\sqrt{2\pi \cdot 2n} \cdot \left(\frac{2n}{e}\right)^{2n} \cdot \left(1 + O\left(\frac{1}{2n}\right)\right)}{\left(\sqrt{2\pi n} \cdot \left(\frac{n}{e}\right)^n \cdot \left(1 + O\left(\frac{1}{n}\right)\right)\right)^2}$$

$$= \frac{1}{\sqrt{2\pi}} \cdot \frac{1}{\sqrt{n}} \cdot \frac{\frac{2^{2n} \cdot n^{2n}}{e^{2n}}}{\frac{n^{2n}}{e^{2n}}} \cdot \left[ \frac{\left(1 + O\left(\frac{1}{2n}\right)\right)}{\left(1 + O\left(\frac{1}{n}\right)\right)^2} \right]$$

$$= \frac{1}{\sqrt{2\pi}} \cdot \frac{4^n}{\sqrt{n}} \cdot \left[ \right]$$

so

$$\frac{\binom{2n}{n}}{4^n} = \frac{1}{\sqrt{2\pi}} \cdot \left[ \right] \xrightarrow[n \rightarrow \infty]{as} \frac{1}{\sqrt{2\pi}} \in (0, \infty)$$



Exercise

find const  $a > 0$  s.t.

$$\binom{Z_n}{n} = \mathcal{O}\left(\frac{a^n}{\sqrt{n}}\right)$$

In general find  $a_k > 0$  s.t.

$$\binom{K_n}{n} = \mathcal{O}\left(\frac{a_k^n}{\sqrt{n}}\right)$$

# Induction Proofs

Let  $P(n)$  be a propositional function,

$$P: \mathbb{Z}^+ \rightarrow \{ \text{Propositions depending on } n \}$$

↑  
or another  
subset of  $\mathbb{Z}$ .

Effectively  $P: \mathbb{Z}^+ \rightarrow \{ \text{true, false} \}$

Goal: Prove  $\forall n \geq 1: P(n)$

I. base case: show  $P(1)$

III.a. induction step

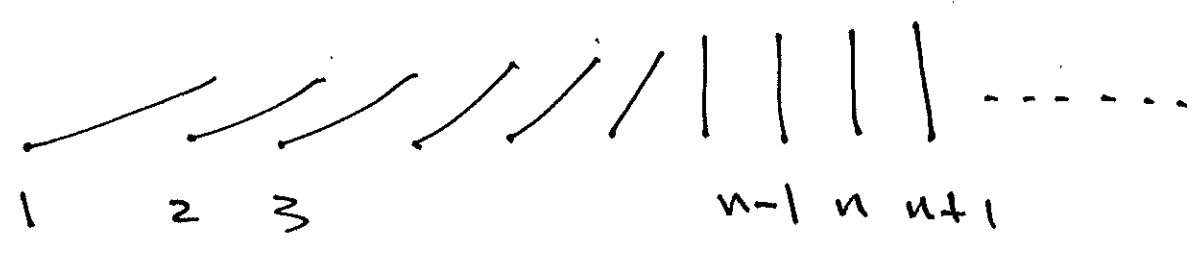
show  $\forall n \geq 1: \underbrace{P(n)}_{\text{induction hypothesis}} \rightarrow P(n+1)$

conclusion:  $\forall n \geq 1 : P(n)$

like an infinite "proof" :

- 1.  $P(1)$  I
- 2.  $P(1) \rightarrow P(2)$  II
- 3.  $P(2)$  modus ponens
- 4.  $P(2) \rightarrow P(3)$  II
- 5.  $P(3)$  M.P.
- ⋮

Domino analogy



$P(n)$  = "the  $n^{\text{th}}$  domino falls"

I:  $P(1)$  = "the 1<sup>st</sup> domino falls"

IIa.  $\forall n \geq 1: P(n) \rightarrow P(n+1)$

= "if any domino falls, then so does the next"

Conclusion:  $\forall n \geq 1: P(n)$

= "all dominoes fall"

Variation

IIb.  $\forall n > 1: P(n-1) \rightarrow P(n)$

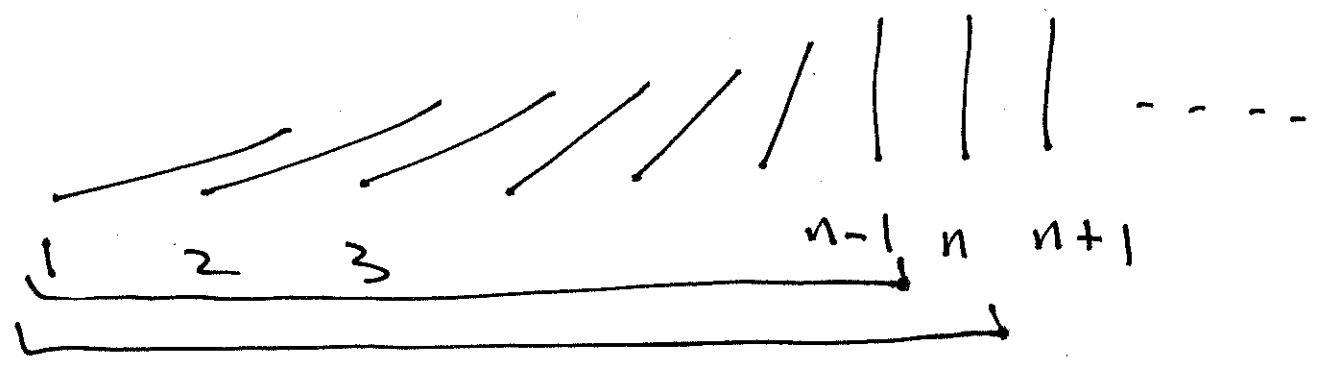
ind. hyp.

II a & II b are called

- weak induction
- 1<sup>st</sup> Principle of Mathematical Induction (PMI).

II c.  $\forall n \geq 1 : \underbrace{P(1) \wedge P(2) \wedge \dots \wedge P(n)} \rightarrow P(n+1)$

$\underbrace{\forall k : 1 \leq k \leq n : P(k)}$   
 induction hyp.



$$\text{III d. } \forall n > 1 : P(1) \wedge \dots \wedge P(n-1) \rightarrow P(n)$$

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$$\forall k : 1 \leq k < n : P(k)$$

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ind. hyp.

III e & III d are called

• strong induction

• 2<sup>nd</sup> PMI

Ex. define

$$T(n) = \begin{cases} 0 & \text{if } n=1 \\ T(\lfloor \frac{n}{2} \rfloor) + 1 & \text{if } n \geq 2 \end{cases}$$

show:  $\forall n \geq 1 : \boxed{T(n) \leq \lg(n)}$   $\leftarrow P(n)$

(hence  $T(n) = O(\lg n)$ )

Proof.

I.  $P(1)$   $\leftarrow$

show  $T(1) \leq \lg(1)$   $\leftarrow$

i.e.  $0 \leq 0$   $\leftarrow$

$$\text{III d. } \forall n > 1 : P(1) \wedge \dots \wedge P(n-1) \rightarrow P(n)$$

$$\forall k : 1 \leq k < n : P(k)$$

- Let  $n > 1$  be chosen arbitrarily.
- Assume for all  $k$  in the range  $1 \leq k < n$  that  $T(k) \leq \lg(k)$ .
- we must show that  $T(n) \leq \lg(n)$ .

Then

$$T(n) = T(\lfloor \frac{n}{2} \rfloor) + 1 \quad \text{by rec. for } T(n)$$

$$\leq \lg(\lfloor \frac{n}{2} \rfloor) + 1 \quad \left\{ \begin{array}{l} \text{by ind. hyp.} \\ \text{with } k = \lfloor \frac{n}{2} \rfloor \end{array} \right.$$

$$\leq \lg(\frac{n}{2}) + 1 \quad \text{since } \lfloor x \rfloor \leq x$$

$$= \lg n - x + x$$

$$= \lg n$$

Hence  $T(n) \leq \lg(n)$  for all  $n \geq 1$  by the 2<sup>nd</sup> PMI. ~~PMI~~

Multiple base case. ∴

I. Prove  $P(1), P(2), \dots, P(n_0)$

II.  $\forall n > n_0: P(1) \wedge \dots \wedge P(n-1) \rightarrow P(n)$

Conclusion:  $\forall n \geq 1: P(n)$

Ex. Define

$$T(n) = \begin{cases} 1 & \text{if } n=1 \\ & \text{or } n=2 \\ 4T(\lfloor \frac{n}{3} \rfloor) + n & \text{if } n \geq 3 \end{cases}$$

show  $\forall n \geq 1 : T(n) \leq n^2$   $\leftarrow P(n)$

(hence  $T(n) = O(n^2)$ )

Proof.

$P(1)$  says  $T(1) \leq 1^2$ , i.e.  
 $1 \leq 1^2$ , i.e.  $1 \leq 1$ . ✓

$P(2)$  says  $T(2) \leq 2^2$ , i.e.  
 $1 \leq 4$ , which is true. ✓

IId.  $\forall n > 2 : P(1) \wedge \dots \wedge P(n-1) \rightarrow P(n)$

$\forall k : 1 \leq k < n : P(k)$

Let  $n > 2$  be arbitrary. Then

$$T(n) = 4T\left(\left\lfloor \frac{n}{3} \right\rfloor\right) + n$$

$$\leq 4 \cdot \left\lfloor \frac{n}{3} \right\rfloor^2 + n$$

by ind. hyp.  
with  $k = \left\lfloor \frac{n}{3} \right\rfloor$   
since  $1 \leq \left\lfloor \frac{n}{3} \right\rfloor < n$

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