

Ex. define $T(n)$ for $n \in \mathbb{N}^+$:

$$T(n) = \begin{cases} 1 & \text{if } 1 \leq n \leq 2 \\ 4T(\lfloor \frac{n}{2} \rfloor) + n & \text{if } n \geq 3 \end{cases}$$

show $\forall n \geq 1$: $\boxed{T(n) \leq n^2}$ hence $T(n) = O(n^2)$.
 $\nwarrow$ $P(n)$

Strong ind., 2 base cases: $n=1, n=2$.

I. $T(1) = 1 \leq 1 = 1^2$ ✓

$T(2) = 1 \leq 4 = 2^2$ ✓

II. $\forall n \geq 2 : P(1) \wedge \dots \wedge P(n-1) \rightarrow P(n)$

Let $n \geq 2$. Assume for all k
in the range $1 \leq k < n$ that

$$T(k) \leq k^2.$$

we must show

$$T(n) \leq n^2.$$

Then

$$T(n) = 4T\left(\left\lfloor \frac{n}{3} \right\rfloor\right) + n$$

$$\leq 4\left\lfloor \frac{n}{3} \right\rfloor^2 + n \quad \left\{ \begin{array}{l} \text{by the ind hyp.} \\ \text{with } k = \left\lfloor \frac{n}{3} \right\rfloor, \\ \text{since } 1 \leq \left\lfloor \frac{n}{3} \right\rfloor < n. \end{array} \right.$$

$$\leq 4\left(\frac{n}{3}\right)^2 + n \quad \left\{ \text{since } \lfloor x \rfloor \leq x \right.$$

$$= \frac{4}{9}n^2 + n \leq n^2$$

$\uparrow$
why?

$$n > 2 \Leftrightarrow \frac{9}{5} \leq n$$

$$\Leftrightarrow n \leq \frac{5}{9} n^2$$

$$\Leftrightarrow \frac{4}{9} n^2 + n \leq n^2 \quad \checkmark$$

$$\therefore T(n) \leq n^2.$$



n	1	2	3	4	5	6	7	8	9
$\lfloor n/3 \rfloor$	.	.	1	1	1	2	2	2	3


 Base
 cases.

Graph Theory

Defn a graph $G = (V, E)$ is a pair of sets

Vertices: $V \neq \emptyset$ vertices

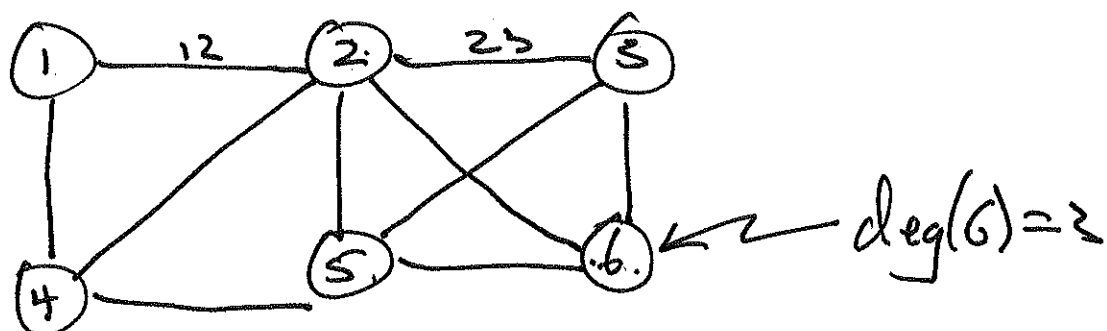
Edges: $E \subseteq V^{(2)} = \{2\text{-element sets of } V\}$



$$\{x, y\} = xy = yx$$

Ex.

5



$$V = \{1, 2, 3, 4, 5, 6\}$$

$$E = \{12, 14, 23, 24, 25, 26, 35, 36, 45, 56\}$$

We say: 2 is adjacent to 3

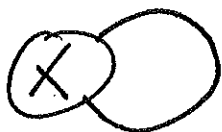
23 joins 2 to 3

23 has ends 2 and 3

12 is adjacent to 23

2 is incident with 23

Note: not allowed



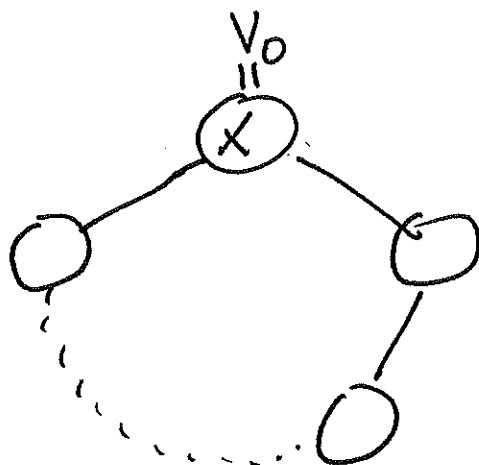
Let $x, y \in V$. An x - y Path in G is a sequence

$$X = v_0, v_1, v_2, \dots, v_k = y$$

of vertices in which $\{v_i, v_{i+1}\} \in E(G)$ for all $i = 0, 1, \dots, k-1$. Also there are no repeated vertices (except possibly when $x = y$.) If $x = y$ the Path is closed.

Trivial Path: $x = v_0 = y$

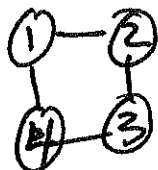
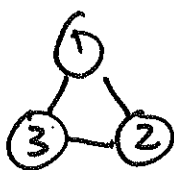
a non-trivial closed path



is called a cycle.

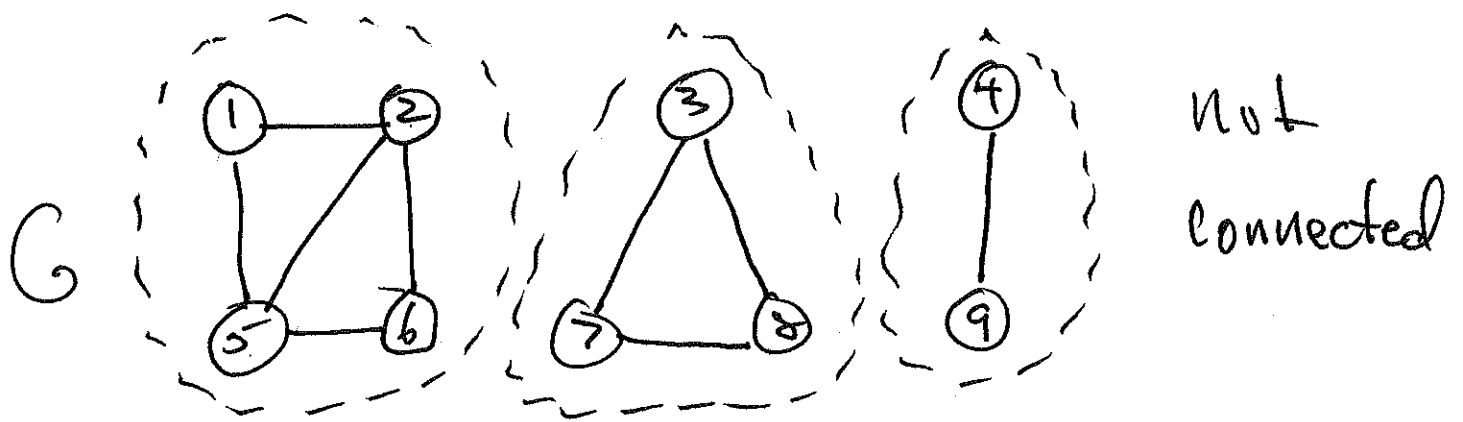
length of Path = # edges traversed.

Cycles:



len = 3, 4, 5, 6, ...

G is called connected if for all $x, y \in V$, G contains an x - y Path.

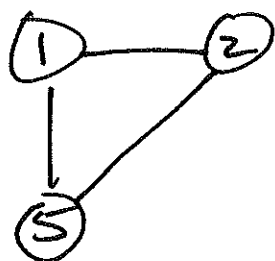


A subgraph of G is a graph H such that

$$V(H) \subseteq V(G)$$

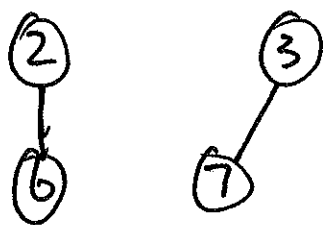
$$E(H) \subseteq E(G)$$

Note: $(\{1, 2, 5\}, \{12, 15, 25\})$



is a connected subgraph of G .

Also: $(\{2, 3, 6, 7\}, \{26, 37\})$



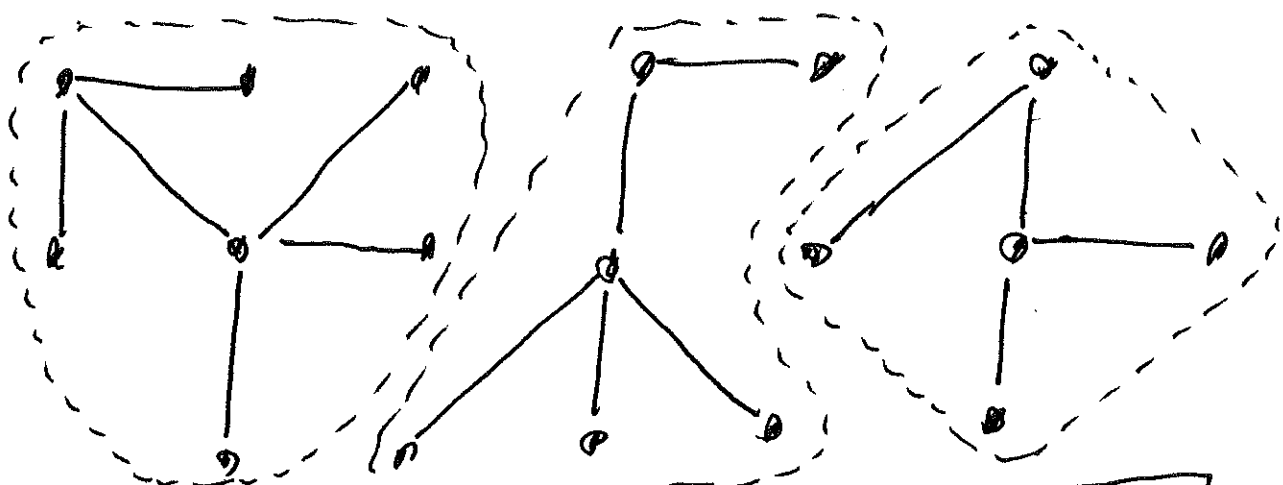
is a disconnected subgraph of G .

A subgraph H of G is a connected component iff

(1) H is connected

(2) H is maximal w.r.t. (1)

G is called acyclic iff it contains no cycles. (also called a forest.)



#vert:	7	6	5
#edge:	6	5	4

A graph is a tree iff it is both connected and acyclic.

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Ex. Prove $\forall n \geq 1$: if T is
a tree on n vertices,
then T has $(n-1)$ edges.

$$\forall n \geq 1: |V(T)| = n \Rightarrow |E(T)| = n-1$$

$\uparrow$
 $P(n)$

Proof:

I. let T be a tree with 1 vertex.
Then $T = (\{1\}, \emptyset)$.

①

In this case T has $1-1=0$
edges.

III d. $\forall n > 1 : P(1) \wedge \dots \wedge P(n-1) \rightarrow P(n)$

Let $n > 1$ be arbitrary. Assume for all k in the range $1 \leq k < n$ that: if T' is a tree on k vertices, then T' has $(k-1)$ edges.

We must show that: if T is a tree on n vertices, then T has $(n-1)$ edges.

Let T be a tree on n vertices.

Pick any edge $e \in E(T)$ and remove it from T . This

disconnects T into 2 subtrees, call them T_1 and T_2 .

Let $|V(T_1)| = k_1$, and $|V(T_2)| = k_2$

Note $1 \leq k_1 < n$ and $1 \leq k_2 < n$

By the ind. hyp. we have

$$|E(T_1)| = k_1 - 1 \text{ and } |E(T_2)| = k_2 - 1,$$

Note also that $k_1 + k_2 = n$, since no vertices were removed. Therefore

$$|E(T)| = |E(T_1)| + |E(T_2)| + 1$$

$$= (k_1 - 1) + (k_2 - 1) + 1$$

$$= (k_1 + k_2) - 1$$

$$= n - 1.$$



Induction Principle

Theorem 1

for any Prop. fcn $P(n)$, the following sentence is true.

$$\left[P(1) \wedge (\forall n \geq 1: P(n-1) \Rightarrow P(n)) \right] \rightarrow \left[\forall n \geq 1: P(n) \right]$$

will use: well ordering property $\mathbb{Z}^+$.

any non-empty subset of $\mathbb{Z}^+$
contains a least element.

Proof of Thm 1

Assume $P(1)$ and $\forall n \geq 1: P(n-1) \rightarrow P(n)$ are both true. Define

$$S = \{n \in \mathbb{Z}^+ \mid P(n) \text{ is false}\}$$

it's sufficient to show $S = \emptyset$.

Assume, to get a X , that $S \neq \emptyset$.

By the W.O.P, S contains a least element, call it m .

Since $P(1)$ is true, $1 \notin S$, so $m \neq 1$. $\therefore m > 1$, and $m-1$ is a positive int.

Note $m-1 \notin S$, so $P(m-1)$ is

true. we have $\forall n > 1: P(n-1) \Rightarrow P(n)$

$\therefore$ letting $n=m$ we have

$P(m-1) \Rightarrow P(m)$ is true

so $P(m)$ is true.

$\therefore m \notin S$.

But this contradicts defn of m as least element in S . $\therefore X$

conclude $S = \emptyset$.

