

Induction Proofs

Goal: Prove stunts of the form

$$\forall n \geq n_0 : P(n)$$

where $P(n)$ is a Propositional function

$$P: \mathbb{Z} \rightarrow \{\text{true}, \text{false}\}$$

↑ or a subset of $\mathbb{Z}$,
usually $\{0, 1, 2, \dots\}$ or
 $\{1, 2, 3, \dots\}$

Have 2 steps

I. Base: Prove $P(n_0)$ is true

II a. Induction step:

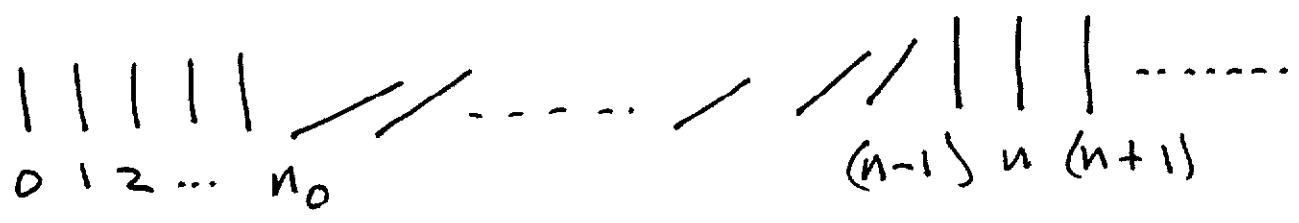
Prove $\forall n \geq n_0: \underbrace{P(n)}_{\text{induction hyp.}} \rightarrow P(n+1)$

when I & II a are done, we conclude

$$\forall n \geq n_0: P(n)$$

is true

Domino analogy:



$P(n)$ = "the n^{th} domino falls"

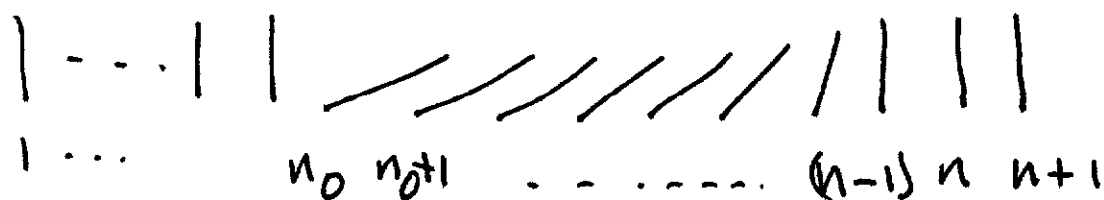
Variations on II :

II b. Prove $\forall n \geq n_0 : \underbrace{P(n-1)}_{\text{ind. hyp.}} \rightarrow P(n)$

II a & II b are called weak induction
or the 1st Principle of mathematical induction. (P.M.I.)

II c & II d below are strong induction
or the 2nd P.M.I.

II c : Show $\forall n \geq n_0 : \underbrace{P(n_0) \wedge P(n_0+1) \wedge \dots \wedge P(n)}_{\substack{(\forall k \leq n : P(k)) \\ \text{induction hypothesis}}} \rightarrow \underbrace{P(n+1)}_{\substack{\text{induction} \\ \text{conclusion}}}$



II d. Show $\forall n \geq n_0 : \underbrace{P(n_0) \wedge \dots \wedge P(n-1)}_{\substack{\forall k < n : P(k) \\ \text{ind. hyp.}}} \rightarrow P(n)$

same conclusion :

$$\forall n \geq n_0 : P(n)$$

is true.

Ex. $\forall n \geq 1$:

$$\boxed{\sum_{i=1}^n i^2 = \frac{n(n+1)(2n+1)}{6}}$$

$\uparrow P(n)$

Proof.

I. $P(1)$ says $\sum_{i=1}^1 i^2 = \frac{1(1+1)(2 \cdot 1+1)}{6}$

i.e. $1^2 = \frac{1 \cdot 2 \cdot 3}{6}$, i.e. $1=1$ ✓

II a. show $\forall n \geq 1 : P(n) \rightarrow P(n+1)$.

Let $n \geq 1$ be chosen arbitrarily.

Assume $\sum_{i=1}^n i^2 = \frac{n(n+1)(2n+1)}{6}$.

we must show $\sum_{i=1}^{n+1} i^2 = \frac{(n+1)(n+2)(2n+3)}{6}$.

so $\sum_{i=1}^{n+1} i^2 = \left(\sum_{i=1}^n i^2 \right) + (n+1)^2$

$$= \frac{n(n+1)(2n+1)}{6} + (n+1)^2 \quad \left\{ \begin{array}{l} \text{by the} \\ \text{induction} \\ \text{hypothesis.} \end{array} \right.$$

$$= \frac{n(n+1)(2n+1) + 6(n+1)^2}{6}$$

$$= \frac{(n+1)[n(2n+1) + 6(n+1)]}{6}$$

$$= \frac{(n+1)[2n^2 + n + 6n + 6]}{6}$$

$$= \frac{(n+1)(2n^2 + 7n + 6)}{6}$$

$$= \frac{(n+1)(n+2)(2n+3)}{6}$$

The result follows for all $n \geq 1$
by the 1st P.M.I. . . 

Ex. let $x \in \mathbb{R}, x \neq 1$.

$$\forall n \geq 0 : \boxed{\sum_{i=0}^n x^i = \frac{x^{n+1} - 1}{x - 1}} \leftarrow P(n)$$

Proof.

I. $P(0)$ says $\sum_{i=0}^0 x^i = \frac{x^1 - 1}{x - 1}$, i.e. $1 = 1$ ✓

II b. show $\forall n \geq 0 : P(n-1) \rightarrow P(n)$.

let $n \geq 0$ be arbitrary. assume

$$\sum_{i=0}^{n-1} x^i = \frac{x^n - 1}{x - 1} \quad \text{must show that}$$

$$\sum_{i=0}^n x^i = \frac{x^{n+1} - 1}{x - 1} \quad \Rightarrow$$

$$\sum_{i=0}^n x^i = \left(\sum_{i=0}^{n-1} x^i \right) + x^n$$

$$= \left(\frac{x^n - 1}{x - 1} \right) + x^n \quad \left\{ \begin{array}{l} \text{by the} \\ \text{ind. hyp.} \end{array} \right.$$

$$= \frac{x^n - 1 + x^n(x - 1)}{x - 1}$$

$$= \frac{\cancel{x^n} - 1 + x^{n+1} - \cancel{x^n}}{x - 1}$$

$$= \frac{x^{n+1} - 1}{x - 1}$$

result follows by 1st P.M.I. ~~■~~

Always state

- let $n > n_0$ or $n \geq n_0$
- ind. hyp. explicitly
- ind. conc. explicitly
- the point(s) where the ind. hyp. is used.

Exercises

• re-do 1st ex. using III b

• " 2nd " " II a

• $\forall n \geq 1: \sum_{i=1}^n i^3 = \left(\frac{n(n+1)}{2} \right)^2$

• $\forall n \geq 1: \sum_{i=1}^n i^4 = \frac{n(n+1)(2n+1)(3n^2+3n-1)}{30}$

Ex. Define $T(n)$ by

$$T(n) = \begin{cases} 0 & \text{if } n=1 \\ T(\lfloor \frac{n}{2} \rfloor) + 1 & \text{if } n \geq 2 \end{cases}$$

n	1	2	3	4	5	6	7	8	...
$T(n)$	0	1	1	2	2	2	2	3	...

Prove that: $\forall n \geq 1 : \boxed{T(n) \leq \lg(n)} \leftarrow P(n)$

Proof:

I. $P(1)$ says $T(1) \leq \lg(1)$, i.e. $0 \leq 0$ ✓

IId. $\forall n > 1 : P(1) \wedge \dots \wedge P(n-1) \rightarrow P(n)$

Let $n > 1$ be arbitrary. Assume, for this n , that $T(k) \leq \lg(k)$ for all k in the range $1 \leq k < n$.

we must show $T(n) \leq \lg(n)$. so

$$T(n) = T\left(\left\lfloor \frac{n}{2} \right\rfloor\right) + 1 \quad \left\{ \begin{array}{l} \text{by the recurrence} \\ \text{for } T(n) \end{array} \right.$$

$$\leq \lg\left(\left\lfloor \frac{n}{2} \right\rfloor\right) + 1 \quad \left\{ \begin{array}{l} \text{by the ind. hyp.} \\ \text{with } k = \left\lfloor \frac{n}{2} \right\rfloor. \text{ note} \\ \text{that } 1 \leq \left\lfloor \frac{n}{2} \right\rfloor < n. \end{array} \right.$$

$$\leq \lg\left(\frac{n}{2}\right) + 1 \quad \left\{ \begin{array}{l} \text{since } \lfloor x \rfloor \leq x \\ \text{and } \lg(\cdot) \text{ is} \\ \text{increasing.} \end{array} \right.$$

$$= \lg(n) - \lg(2) + 1$$

$$= \lg(n) - 1 + 1$$

$$= \lg(n)$$

$\therefore T(n) \leq \lg(n)$ holds for all

$n \geq 1$ by the 2^{nd} PMI. ■