

CSE 101 3-7-25

11

Rotations:

note:

$$\text{key}(x) \leq \text{key}(y) \leq \text{key}(p) \leq \text{key}(x) \leq \text{key}(z)$$

So Rotations Preserve BST Properties

• Summarize RightRotate(T, x):

$$p \begin{cases} x.\text{left} = y.\text{right} \\ y.\text{right}.\text{Parent} = x \end{cases}$$

$$\text{Parent} \begin{cases} y.\text{Parent} = x.\text{Parent} \\ x.\text{Parent}.\text{(left or right)} = y \end{cases}$$

$$x, y \begin{cases} y.\text{right} = x \\ x.\text{Parent} = y \end{cases}$$

13.3 Insertion

2 cases

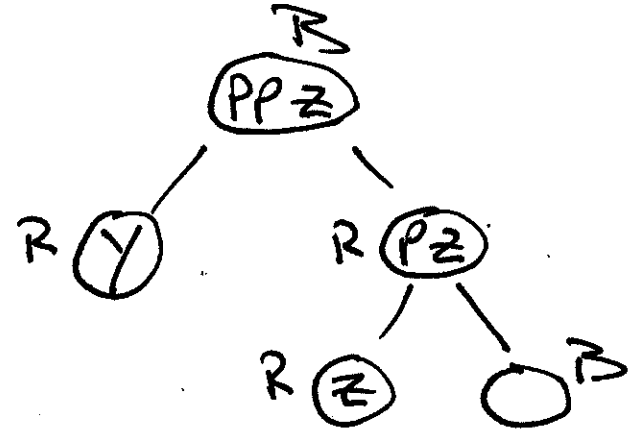
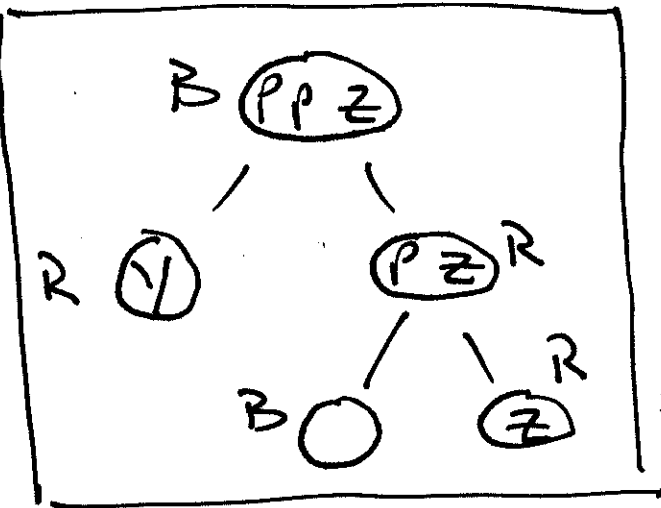
- z.Parent is left child of its Parent: cases 1, 2, 3

let y be z.Parent right sibling,
i.e. z's uncle.

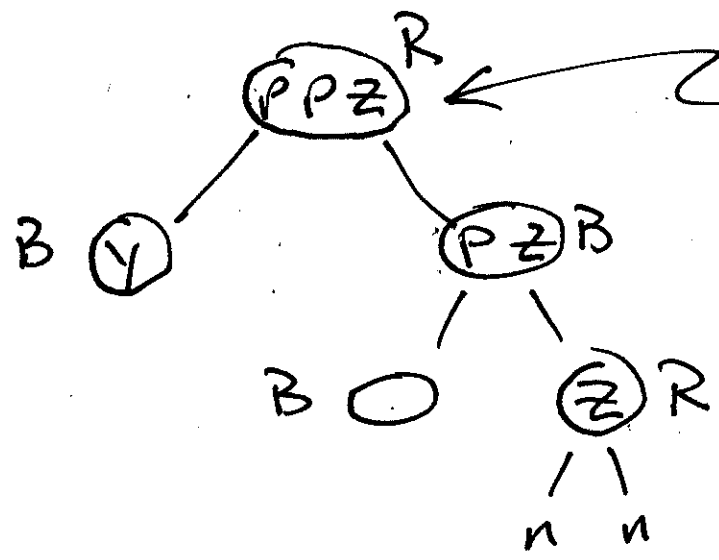
- z.Parent is right child of its Parent: cases 4, 5, 6

let y be z.Parent left sibling,
i.e. z uncle.

Case 4 : y is Red

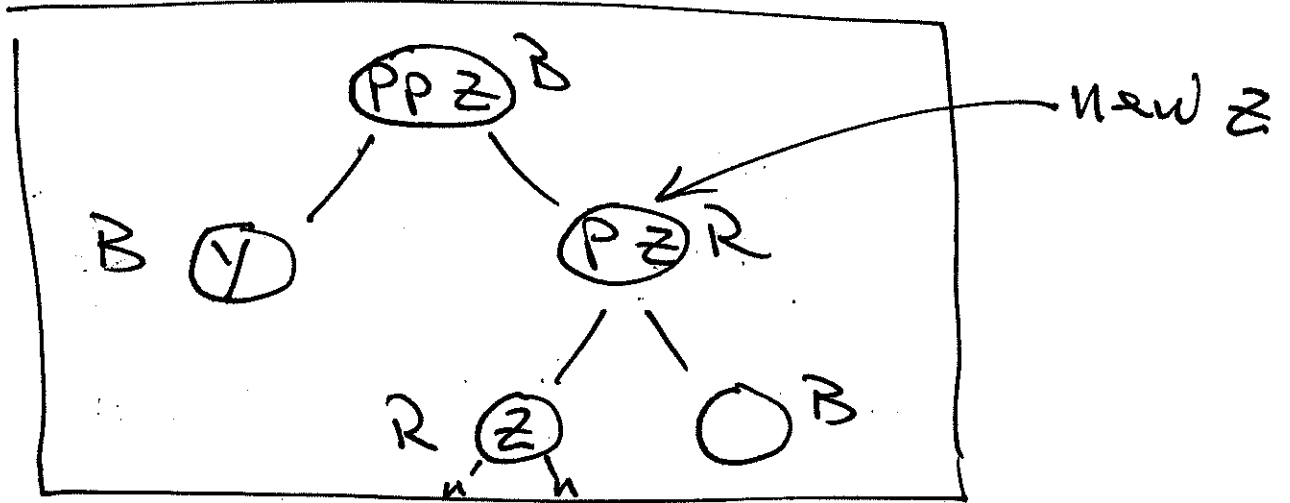


fix colors of y, pz, ppz. let local variable z 'climb up' to ppz.

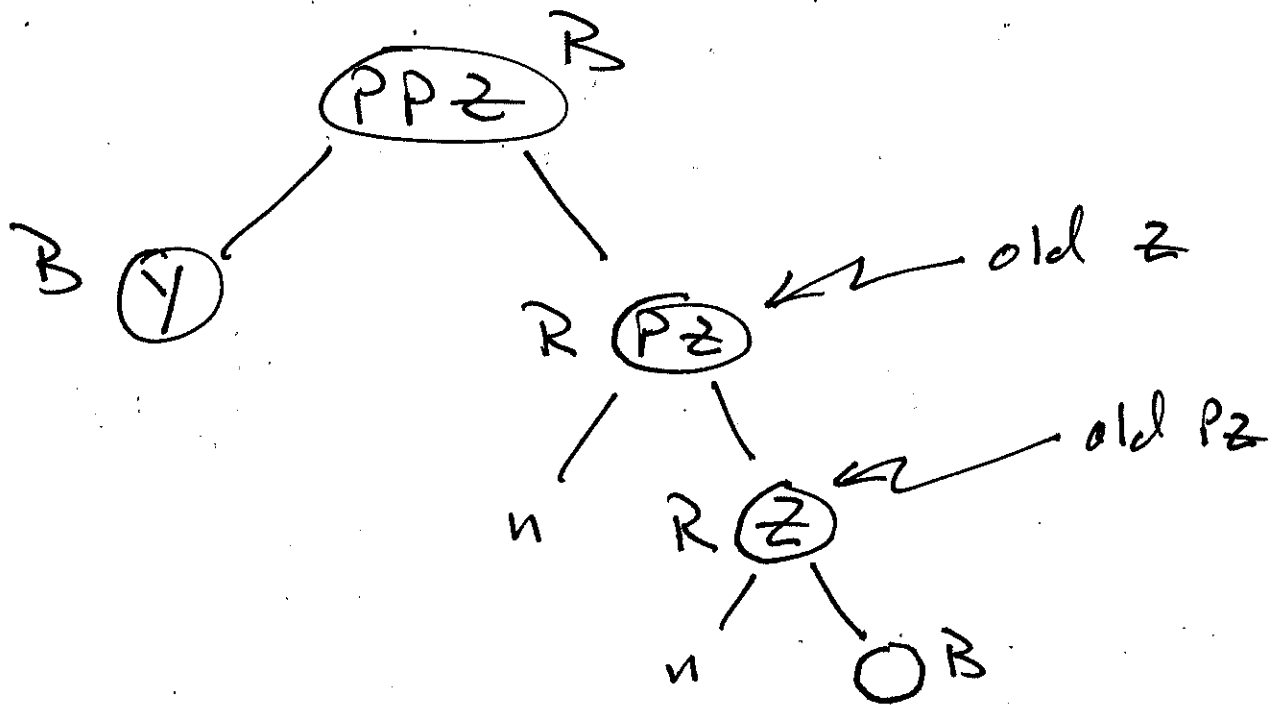


new z
may execute another iteration of while loop.

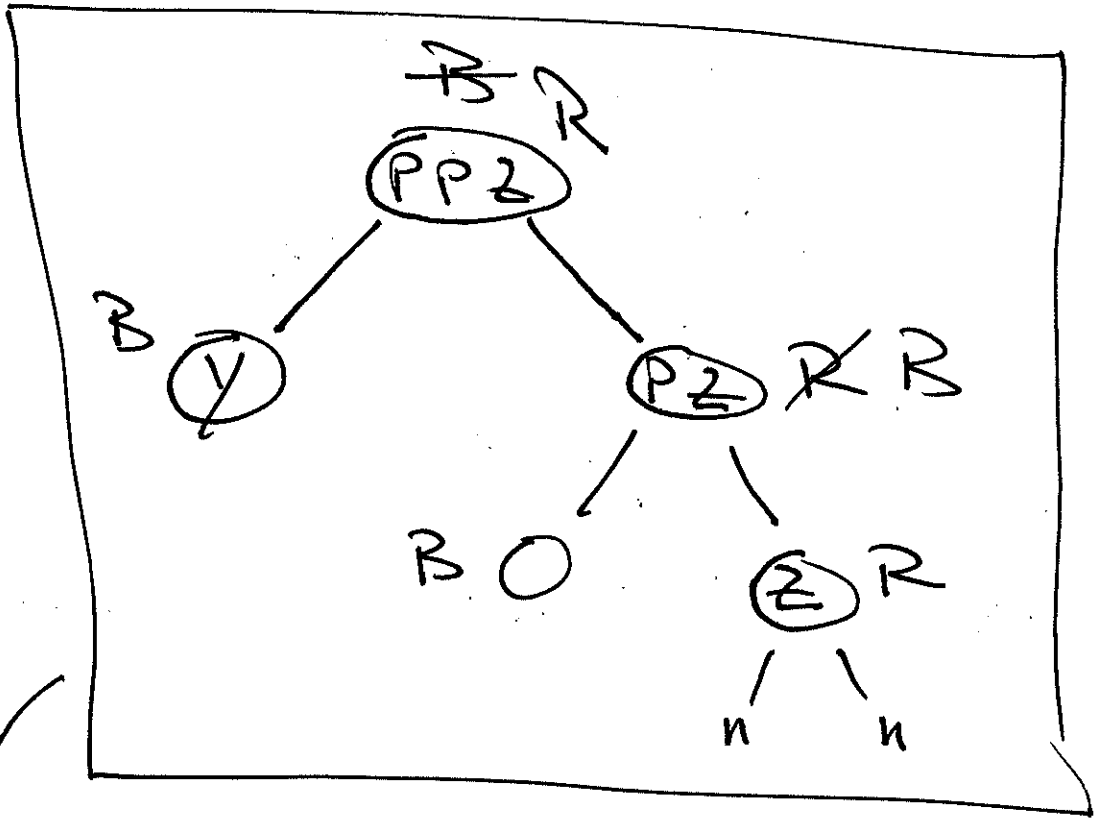
Case 5: y is Black, $z == z.parent.left$



convert to case 6 ($z == z.parent.right$)



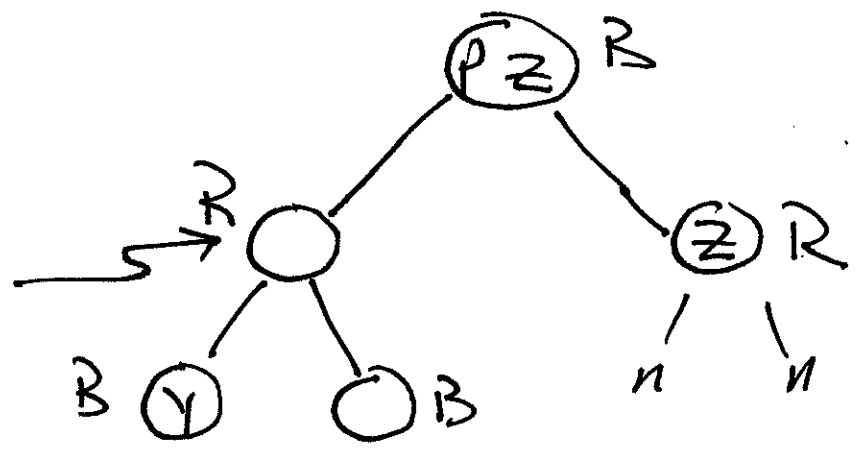
Case 6: y is Black, z == z.Parent.right



color Pz Black, PPz Red, then left Rotate about PPz.



old PPz

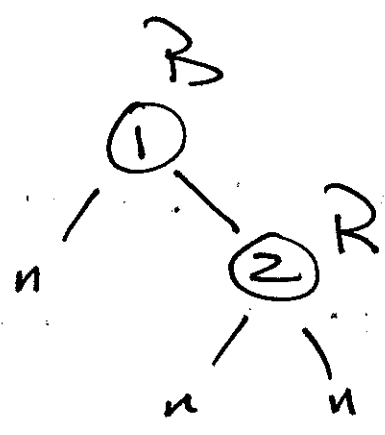


EX Insert 1, 2, 3, 4, 5 into an initially empty RBT.

Insert 1:

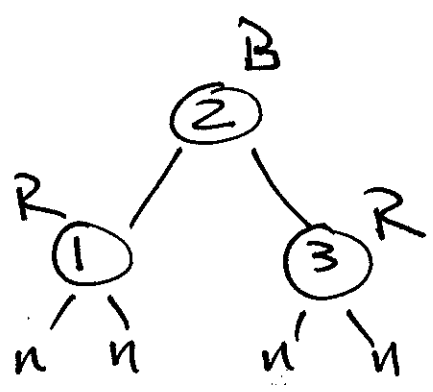
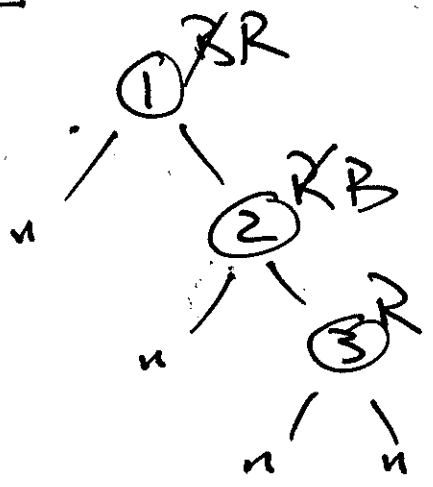


insert 2:



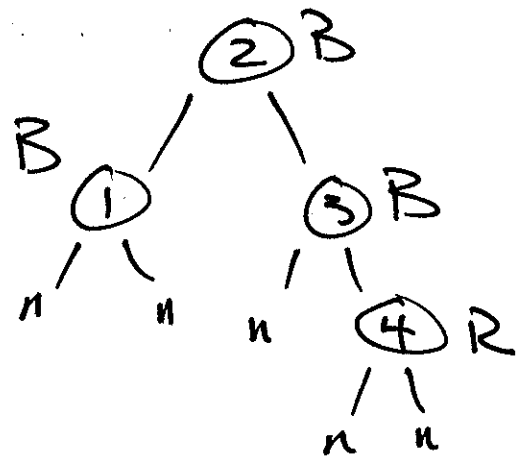
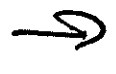
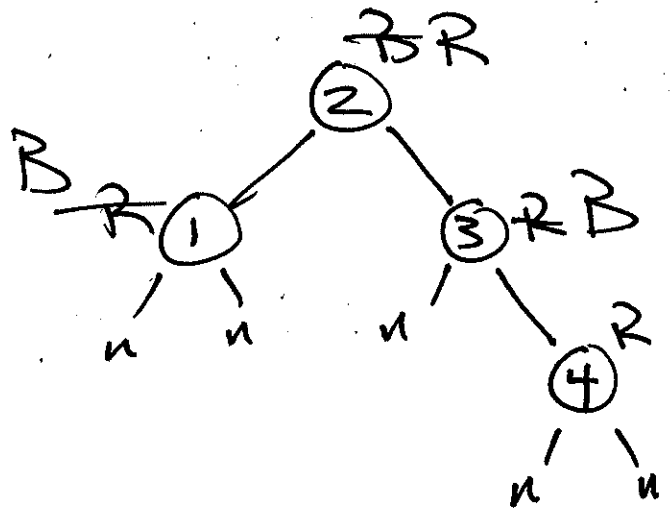
Insert 3:

Case 6

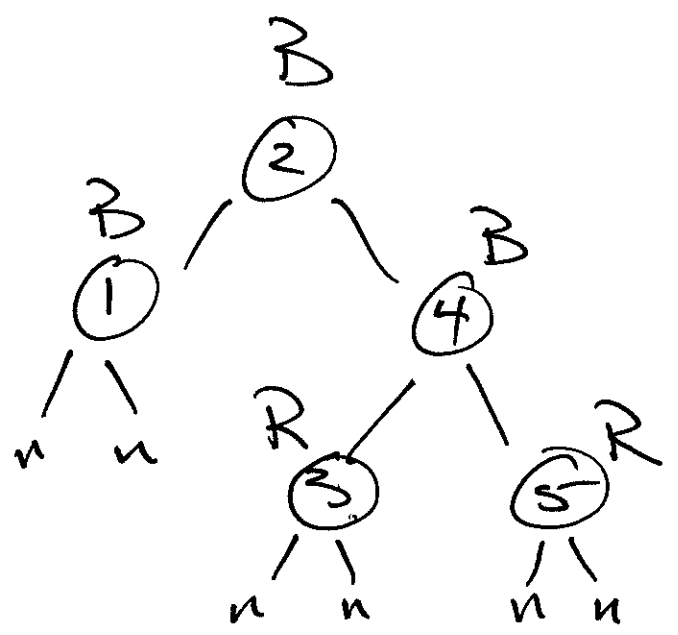
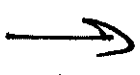
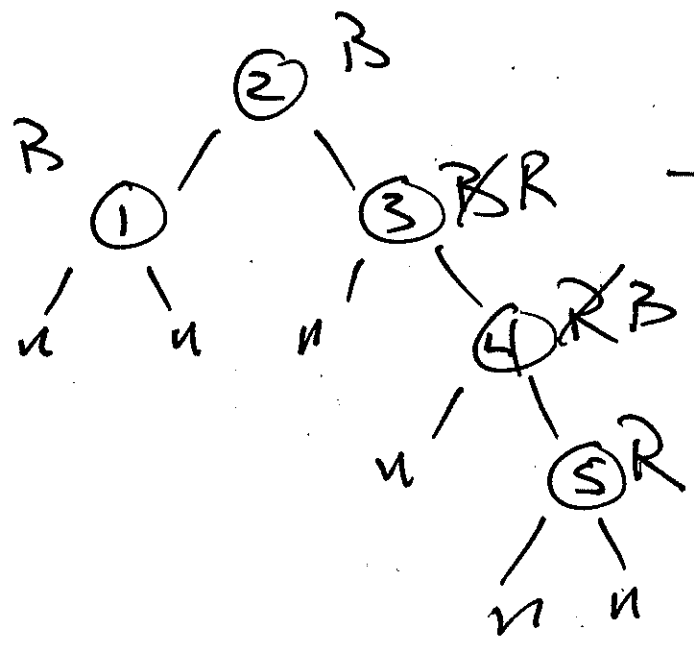


Insert 4

Case 4



Insert 5



Case 6