

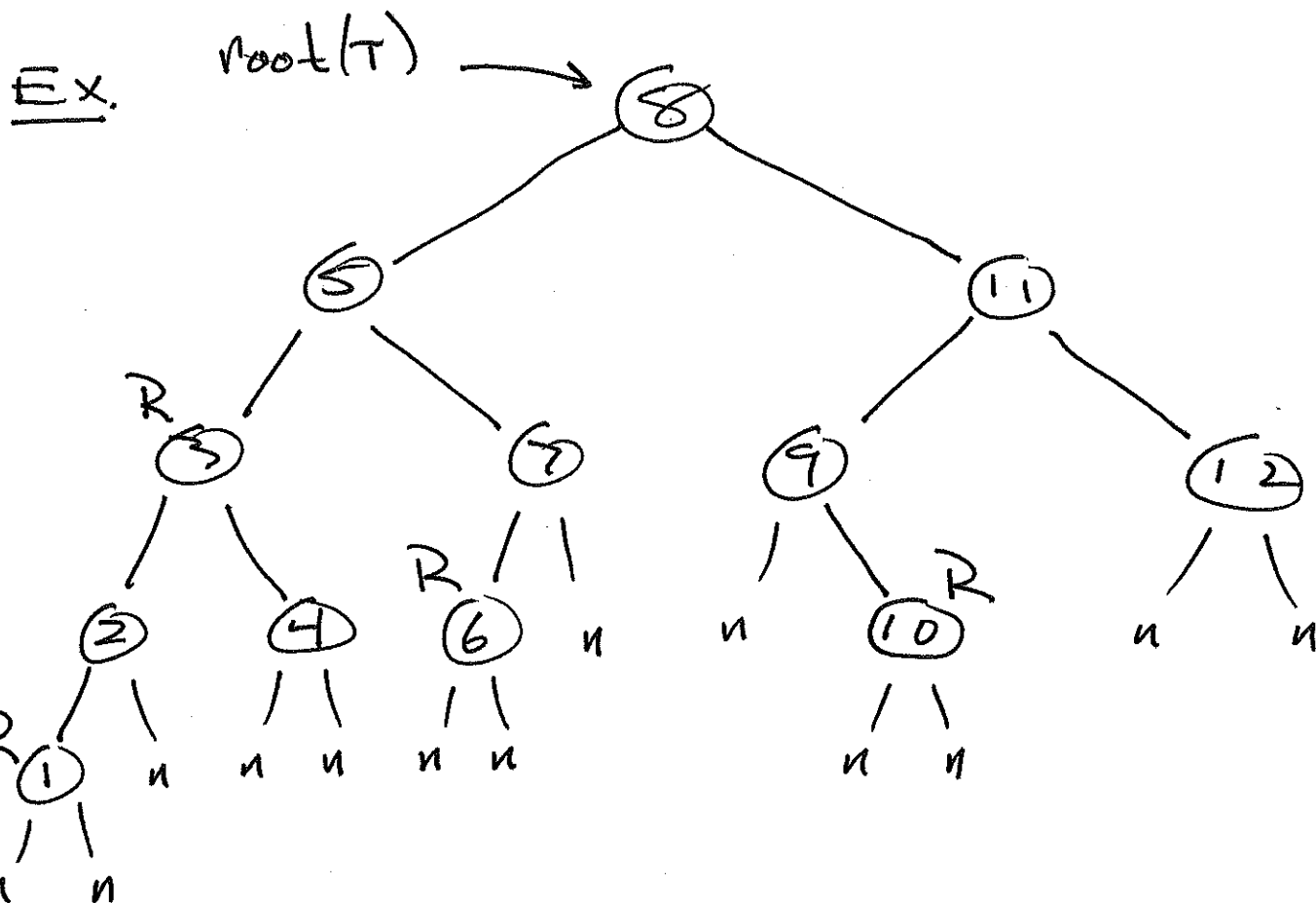
CSE 101-02 5-24-22

11

-
- mid 2 solns posted.
 - Pat 1 is posted
 - Pat 1: ext. last time (wed 10PM).

RBTs ... continue

notes: should say RBT properties.



node

8

5, 11, 3

1, 2, 4, 6, 7, 9, 10, 12

nil leaves

black height

3

2

1

0

note: $bh(x) = 0$ iff $height(x) = 0$ iff x is leaf (nil)

Theorem

An RBT with n internal (non-nil) nodes and height h satisfies

$$h \leq 2 \lg(n+1).$$

Remarks

- any BST satisfies $h \geq \lfloor \lg n \rfloor$

so an RBT satisfies

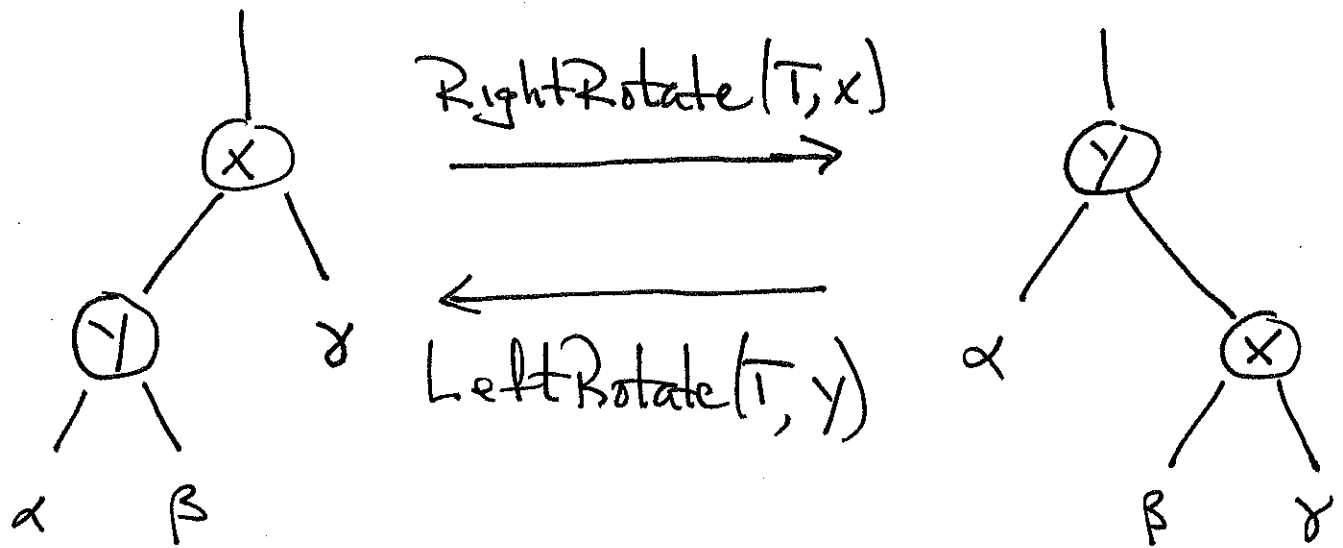
$$\Omega(\lg n) \leq \text{height}(T) \leq O(\lg n)$$

so : $\text{height}(T) = \Theta(\lg n)$.

- Queries run in time $\Theta(\lg n)$.

13.2 Rotations

In Pictures



α, β, γ are arbitrary subtrees.

Preserve BST Properties :

$$\text{key}(\alpha) \leq \text{key}(y) \leq \text{key}(\beta) \leq \text{key}(x) \leq \text{key}(\gamma)$$

Summary of $\text{RightRotate}(T, x)$

$$\beta \quad \left\{ \begin{array}{l} \bullet \text{ left}[x] = \text{right}[y] \\ \bullet \text{ p}[\text{right}[y]] = x \end{array} \right.$$

$$\text{Parent} \quad \left\{ \begin{array}{l} \bullet \text{ p}[y] = \text{p}[x] \\ \bullet \text{ p}[x]'s \text{ (left or right) child} = y \end{array} \right.$$

$$x \leftrightarrow y \quad \left\{ \begin{array}{l} \bullet \text{ right}[y] = x \\ \bullet \text{ p}[x] = y \end{array} \right.$$

See: Pseudo-code

B.3 insertion

- color new node Red
- do BST insert
- fix-up RBT Properties

Pseudo-code

RB_Insert(T, z) ✓

RB_InsertFixUP(T, z)

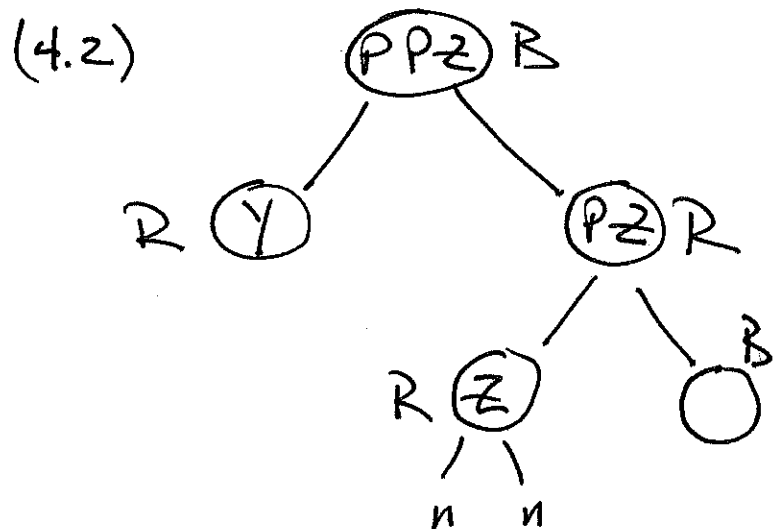
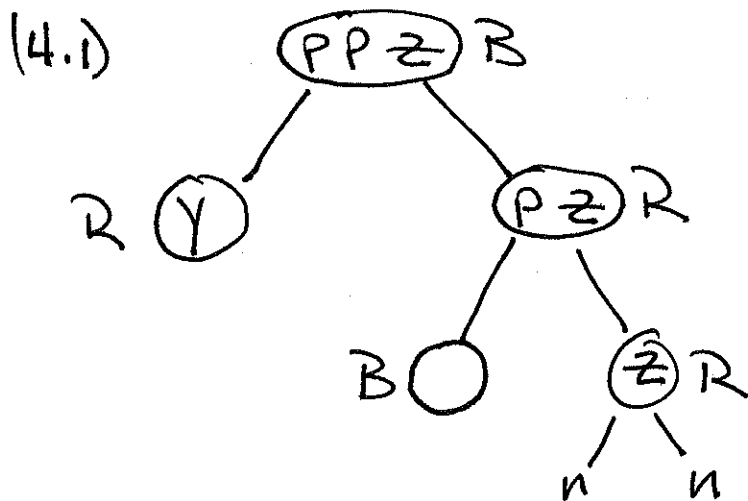
RR-Insert Fixup !

Book does cases 1, 2, 3.

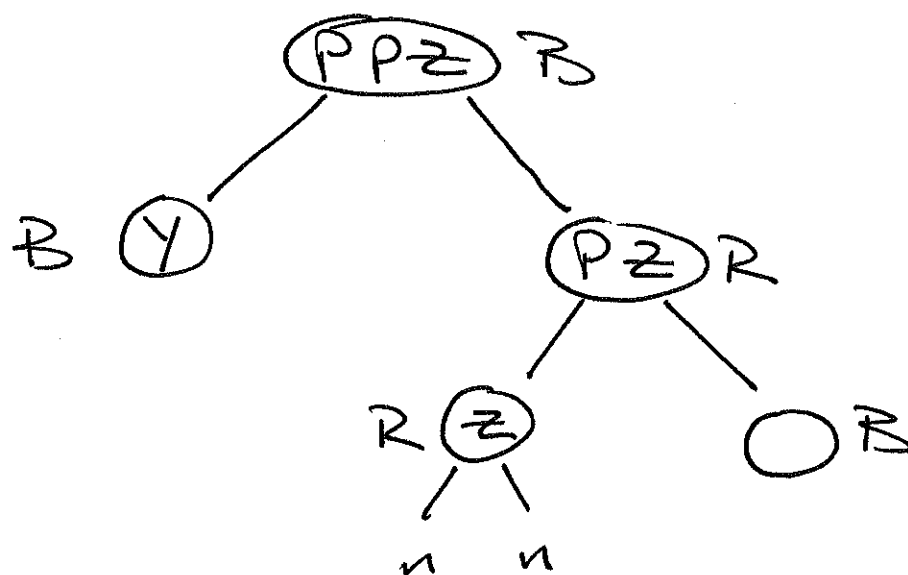
we do cases 4, 5, 6

set $y = \text{left}[P[P[z]]]$ z's uncle

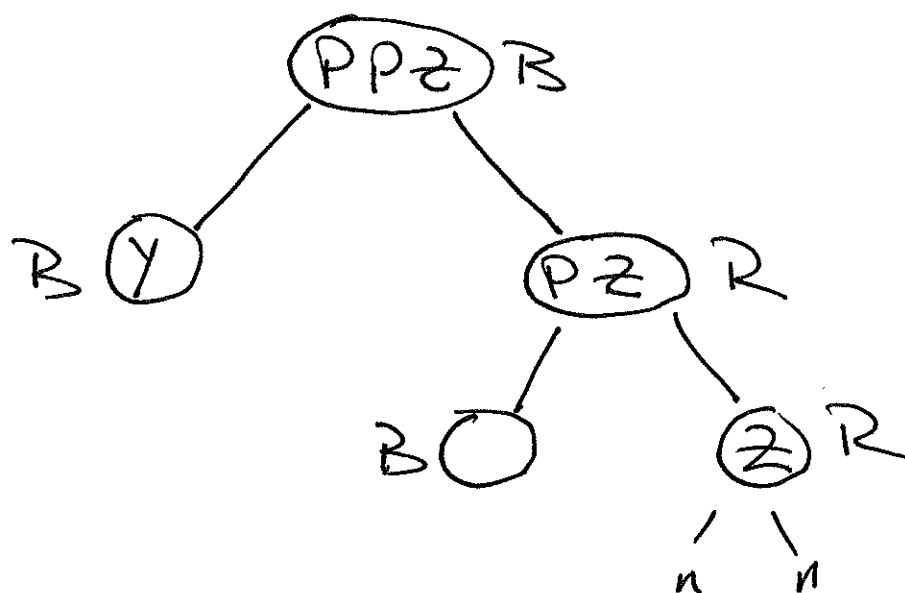
Case 4: y is R



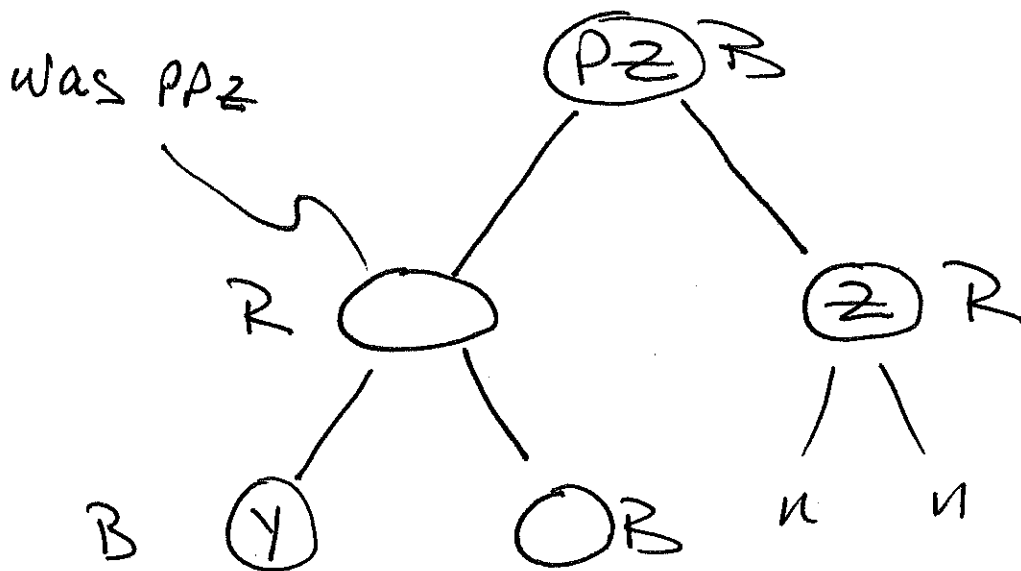
case 5: y is Black, z is left $[P[z]]$



case 6: y is Black, z is right $[P[z]]$

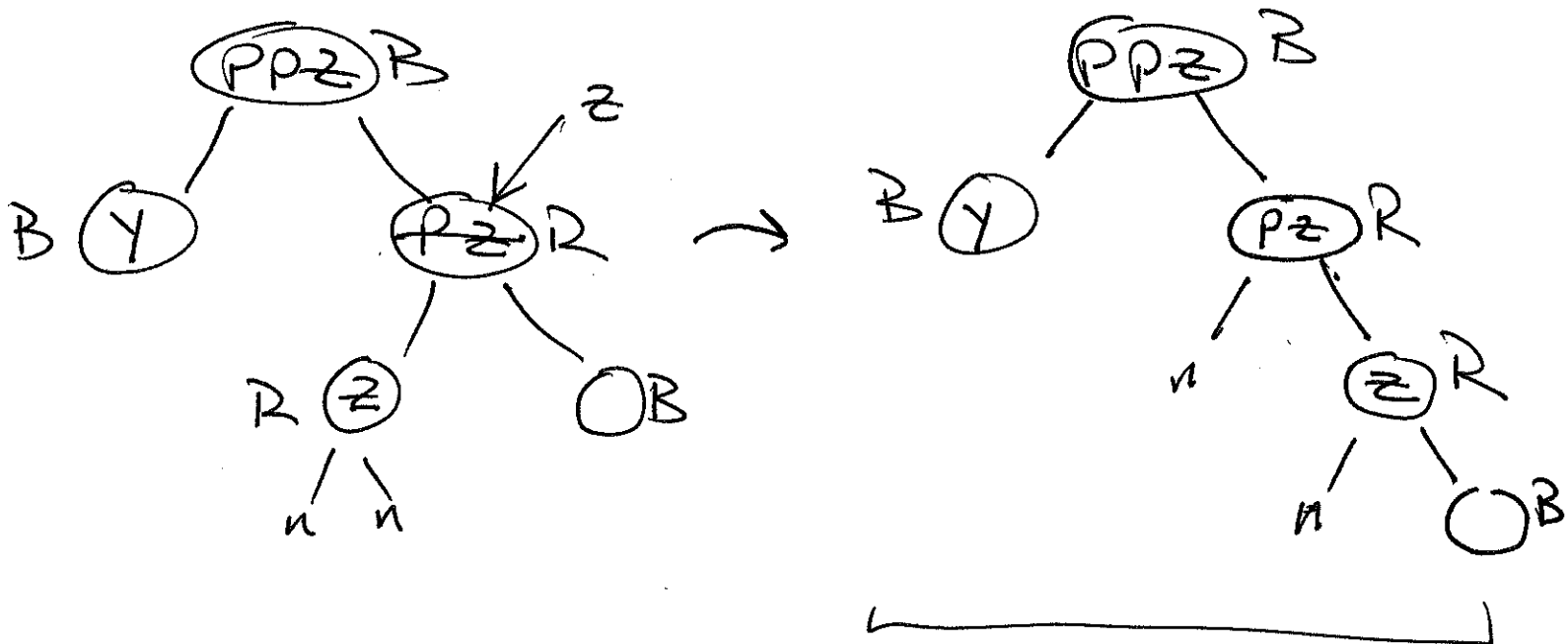


In case 6: color Pz Black,
color PPz Red, then left rotate
about PPz



In case 5: convert to case 6

by letting local $lar \rightarrow$ 'climb
up' to Pz , then right rotate
about z

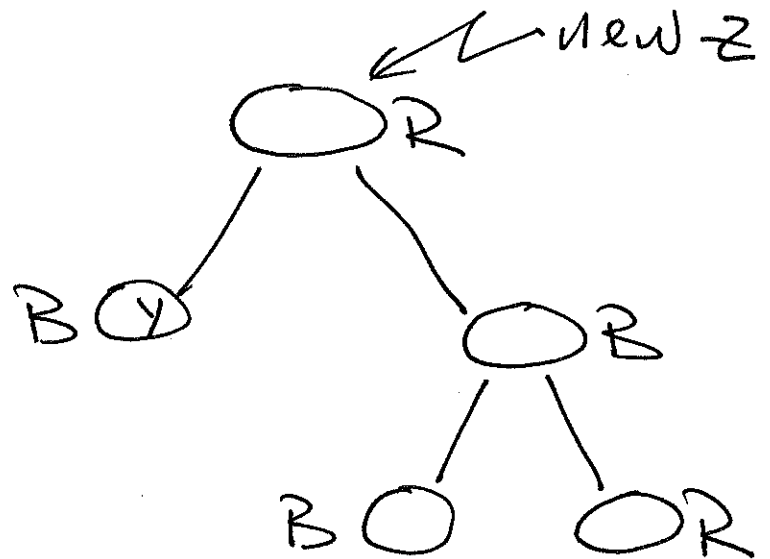
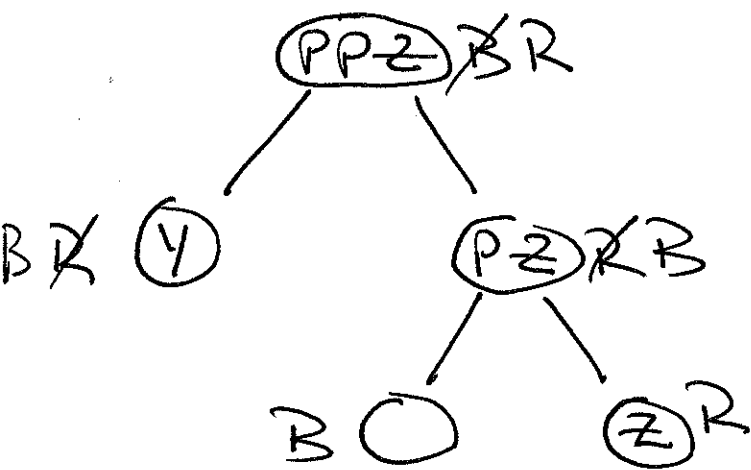


now case 6

In case 4:

fix color relations between y , p_z , and pp_z , then let z 'climb up' to pp_z

(4.1)



may have another iteration of whole loop.

Ex. insert : 1, 2, 3, 4, 5 into
an initially empty RBT.

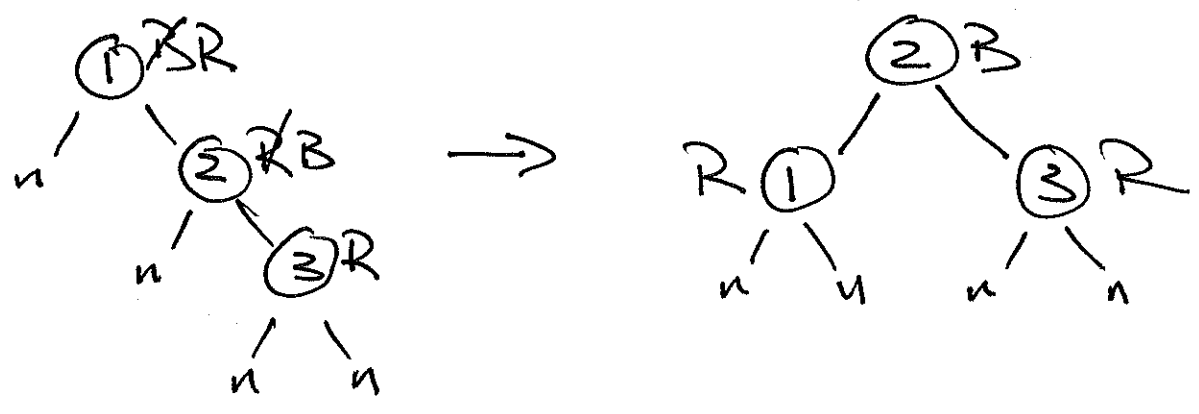
insert(1)



insert(2)

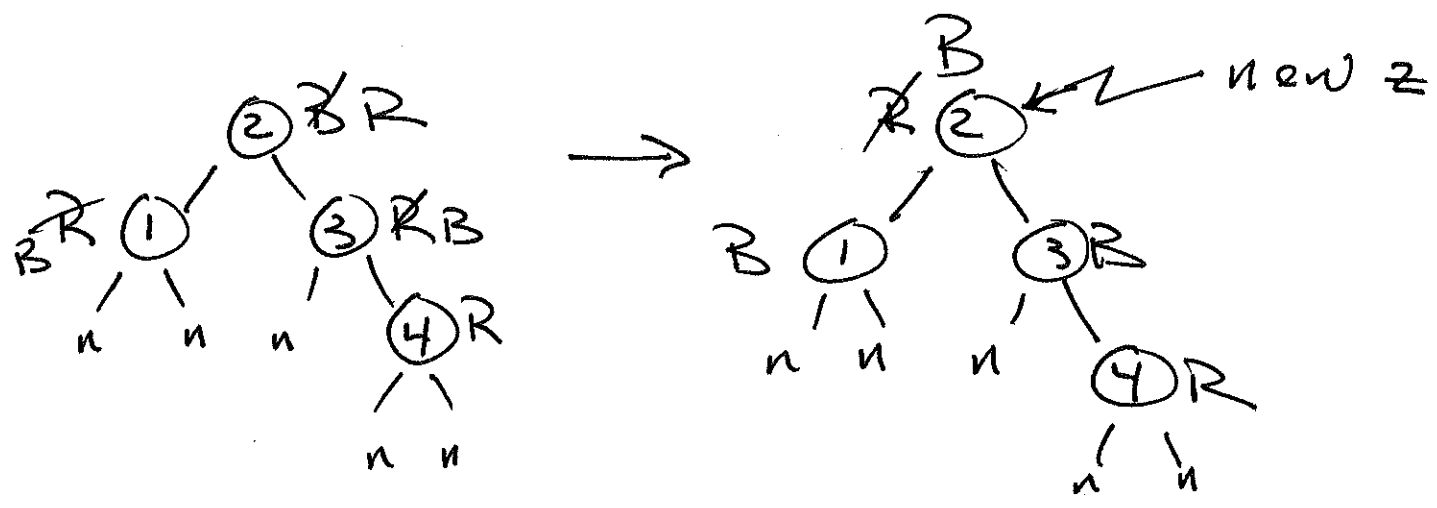


insert(3)



(case 6)

insert(4)



(Case 4)

insert(5)

