

Defn

Let $f(n), g(n)$ be positive func.
we write $f(n) = O(g(n))$ iff
there exists $B > 0, n_0 > 0$ s.t.

$$\frac{f(n)}{g(n)} \leq B$$

for all $n \geq n_0$.

We say: $g(n)$ is an asymptotic
upper bound for $f(n)$.

Ex. $f(n) = 2n + 5$, $g(n) = n$.

claim: $2n + 5 = O(n)$ ✓

check:

$$\frac{2n+5}{n} \leq 3 \quad \forall n \geq 5 \quad \checkmark$$

$\exists$ $\downarrow$ n_0 $\downarrow$

i.e. $2n + 5 \leq 3n \quad \forall n \geq 5 \quad \checkmark$

i.e. $5 \leq n \quad \forall n \geq 5 \quad \checkmark$

Exercise:

show $an + b = O(n)$ for any $a > 0$
and b

Defn

write $f(n) = \Omega(g(n))$ iff

there exist $B > 0, n_0 > 0$ s.t.

$$B \leq \frac{f(n)}{g(n)}$$

for all $n \geq n_0$. we say

$g(n)$ is asymptotic lower bound

for $f(n)$.

Ex. $f(n) = 6n^3 + 4n, g(n) = 2n^2$

claim: $6n^3 + 4n = \Omega(2n^2)$

need

$$\frac{7}{2} \leq \frac{6n^3 + 4n}{2n^2}$$

for all $n \geq 2$

$$\text{i.e. } 14n^2 \leq 6n^3 + 4n \quad \forall n \geq 2$$

$$\text{i.e. } 7n \leq 3n^2 + 2 \quad \forall n \geq 2$$

$$\text{i.e. } 3n^2 - 7n + 2 \geq 0 \quad \forall n \geq 2$$

Exercise Show

$$an^3 + bn^2 + cn + d = \Omega(en^2 + fn + g)$$

for any $a > 0, e > 0, b, c, d, f, g$

Note:

$$\frac{f(n)}{g(n)} \leq B \quad \text{iff} \quad \frac{g(n)}{f(n)} \geq \frac{1}{B}$$

hence

$$f(n) = O(g(n)) \quad \text{iff} \quad g(n) = \Omega(f(n))$$

Defn

write $f(n) = \Theta(g(n))$ iff there exist $B_1 > 0$, $B_2 > 0$, $n_0 > 0$ st.

$$B_1 \leq \frac{f(n)}{g(n)} \leq B_2$$

for all $n \geq n_0$

Equivalently:

$$f(n) = \Theta(g(n))$$

iff both $f(n) = O(g(n))$ and $f(n) = \Omega(g(n))$

Ex. $f(n) = 5n^2 + 11n - 24$, $g(n) = n^2$

claim: $5n^2 + 11n - 24 = \Theta(n^2)$

need:

$$4 \leq \frac{5n^2 + 11n - 24}{n^2} \leq 6$$

check
this

for all $n \geq 8$

note:

we have

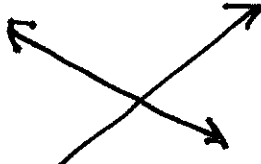
$$f(n) = \Theta(g(n)) \text{ iff } g(n) = \Theta(f(n))$$



$$f(n) = O(g(n))$$

and

$$f(n) = \Omega(g(n))$$



$$g(n) = O(f(n))$$

and

$$g(n) = \Omega(f(n))$$

Defn

write $f(n) = o(g(n))$ iff

$$\lim_{n \rightarrow \infty} \left(\frac{f(n)}{g(n)} \right) = 0$$

We say: $g(n)$ is a strict asymptotic upper bound for $f(n)$.

Ex. $f(n) = \ln(n)$, $g(n) = n$

claim: $\ln(n) = o(n)$ ✓

$$\lim_{n \rightarrow \infty} \left(\frac{\ln(n)}{n} \right) = \lim_{n \rightarrow \infty} \left(\frac{1/n}{1} \right)$$

$$= \lim_{n \rightarrow \infty} \left(\frac{1}{n} \right) = 0$$

Exercise: (a) $\log_b(n) = o(n)$ for any $b > 1$

(b) $\ln(n) = o(n^\alpha)$ for $\alpha > 0$

(c) $\log_b(n) = o(n^\alpha)$ for any $\alpha > 0$, $b > 1$.

Ex. $f(n) = n^k$ ($k \geq 0$ ~~integer~~), $g(n) = e^n$

claim: $n^k = o(e^n)$ ✓

i.e. $\lim_{n \rightarrow \infty} \left(\frac{n^k}{e^n} \right) = \lim_{n \rightarrow \infty} \left(\frac{k n^{k-1}}{e^n} \right)$

$$= \lim_{n \rightarrow \infty} \left(\frac{k(k-1)n^{k-2}}{e^n} \right)$$

⋮

$$= 0 \quad \checkmark$$

Exercise replace e by $b > 1$.

Defn

$f(n) = o(g(n))$ iff

$$\lim_{n \rightarrow \infty} \left(\frac{f(n)}{g(n)} \right) = 0$$

note :

$$\lim_{n \rightarrow \infty} \left(\frac{f(n)}{g(n)} \right) = 0 \quad \text{iff} \quad \lim_{n \rightarrow \infty} \left(\frac{g(n)}{f(n)} \right) = \infty$$

hence

$$f(n) = o(g(n)) \quad \text{iff} \quad g(n) = \omega(f(n))$$

Analogy

is
analogous
to

||

$$f(n) = O(g(n)) \longleftrightarrow x \leq y$$

$$f(n) = \Omega(g(n)) \longleftrightarrow x \geq y$$

$$f(n) = \Theta(g(n)) \longleftrightarrow x = y$$

$$f(n) = o(g(n)) \longleftrightarrow x < y$$

$$f(n) = \omega(g(n)) \longleftrightarrow x > y$$

Part 4: Sparse Matrices

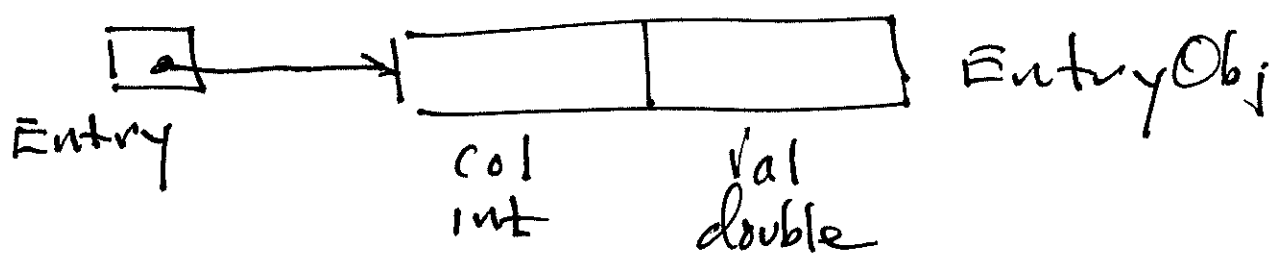
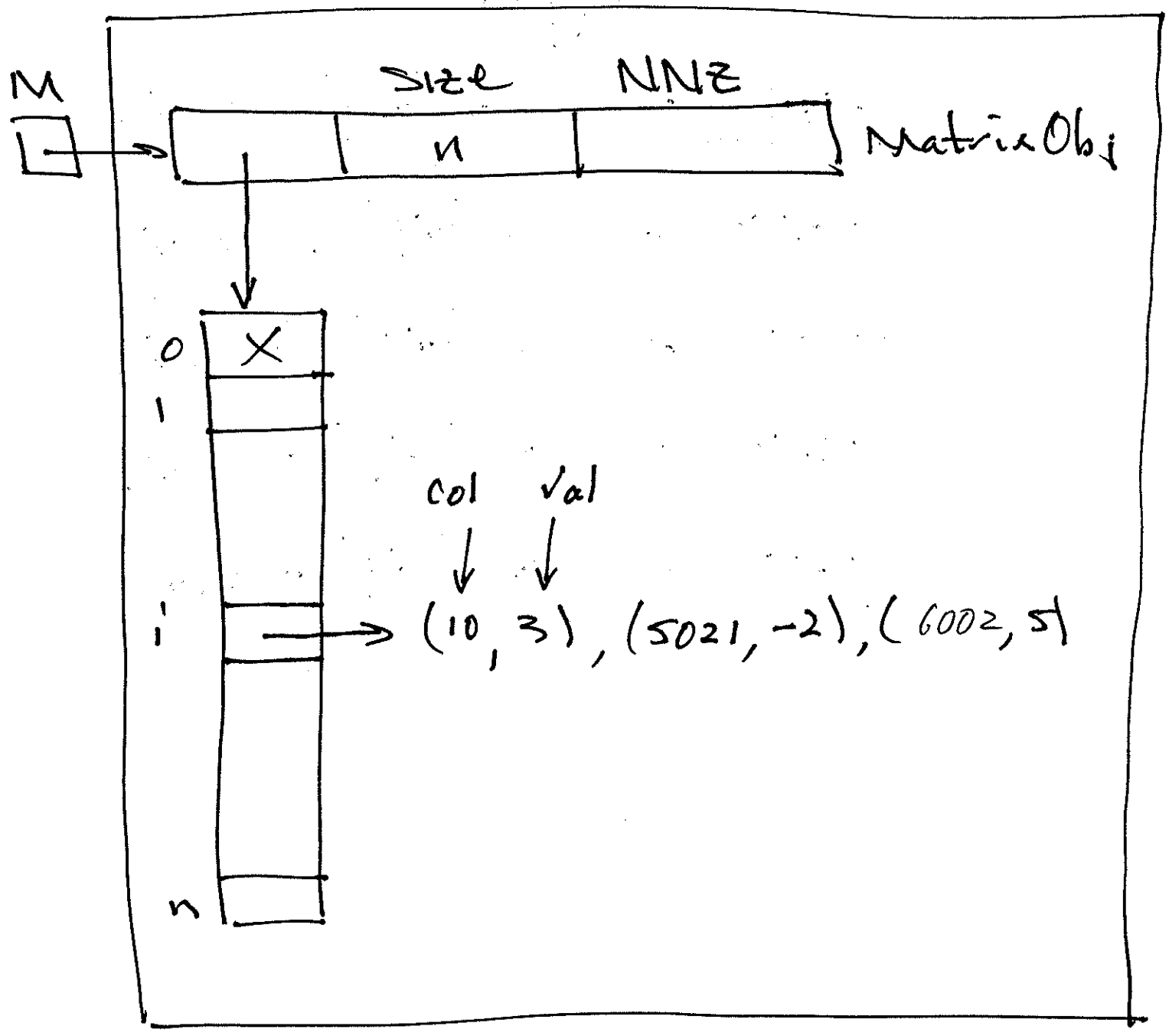


chart view

$$\begin{pmatrix}
 & 10 & & 5021 & & 6002 \\
 & \vdots & & \vdots & & \vdots \\
 & \vdots & & \vdots & & \vdots \\
 & \vdots & & \vdots & & \vdots \\
 & \vdots & & \vdots & & \vdots \\
 & \vdots & & \vdots & & \vdots \\
 0 \dots 0 & 3 & 0 \dots 0 & -2 & 0 \dots 0 & 5 & 0 \dots
 \end{pmatrix}$$