

CSE 101 3-8-22

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• SETs: evals

due: Sun. Mar 13 11:59 PM

incentive: ?

• Par: 2 more days, that's it!

• Quiz 5 tonight

continue... heaps...

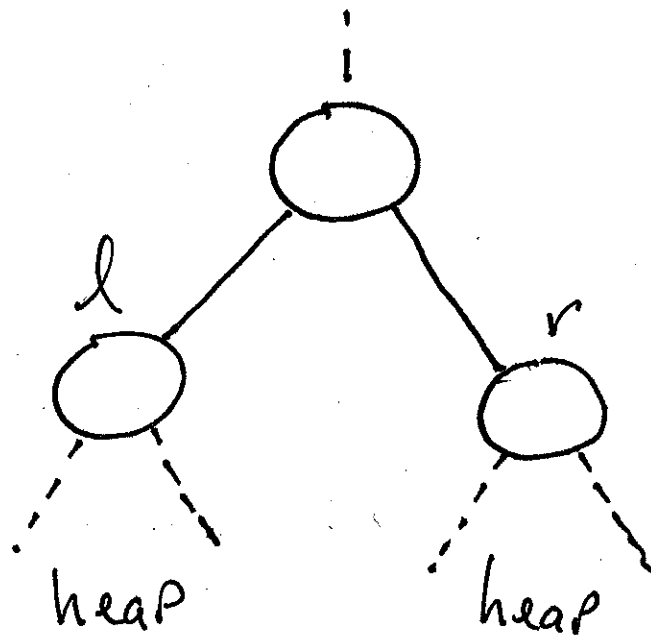
Priority Queue OPS.

- ExtractMax
- IncreaseKey
- Insert
- Maximum

G.2 Heapify

Heapify(A, i) fixes ^(max) heap
Property at index i

Precondition: Subtrees at
 $\text{left}(i)$ and $\text{right}(i)$ are valid
heaps



see Pseudo-code

routine :

$$T(n) \stackrel{\text{def}}{=} \left[\begin{array}{l} \text{worst case \# of} \\ \text{recursive invocations} \\ \text{of Heuristics() on} \\ \text{a subtree (rooted at } i) \\ \text{consisting of } n \text{ nodes} \end{array} \right]$$

Then

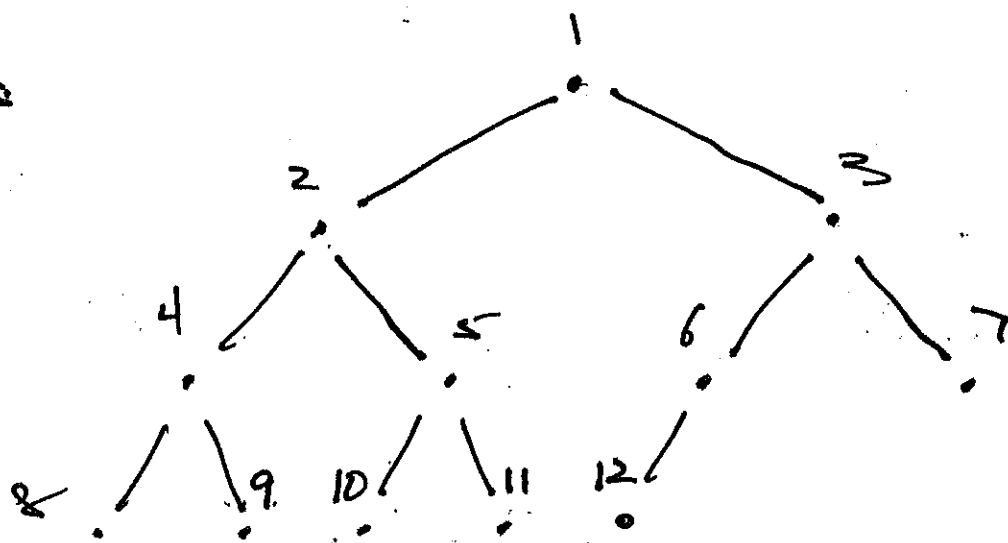
$$\begin{aligned}
 T(n) &= \overbrace{(\text{height of subtree at } i)}^h \\
 &= \lfloor \lg n \rfloor \\
 &= \Theta(\log n) \\
 &= \Theta(h)
 \end{aligned}$$

6.3 Build Heap

Goal: create a max-heap out of an unordered array.
let:

$$n = \text{heapSize}[A] = \text{length}[A]$$

EX



$\underbrace{A_1 \dots A_6}_{\text{internal nodes}} \quad \underbrace{A_7 \dots A_{12}}_{\text{leaves}}$

note:

rightmost internal node is
Parent of rightmost leaf

$$(\text{index of rightmost leaf}) = n$$

$$(\text{index of rightmost internal node}) = \left\lfloor \frac{n}{2} \right\rfloor$$

see pseudo-code

Routine:

let $T(n) \stackrel{\text{def}}{=} \left\{ \begin{array}{l} \text{worst case \# of calls} \\ \text{to heapify() on an} \\ \text{array of length } n \end{array} \right\}$

$$= \Theta(n)$$

↑ see text.

6.4 Heapsort

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See Pseudo-code

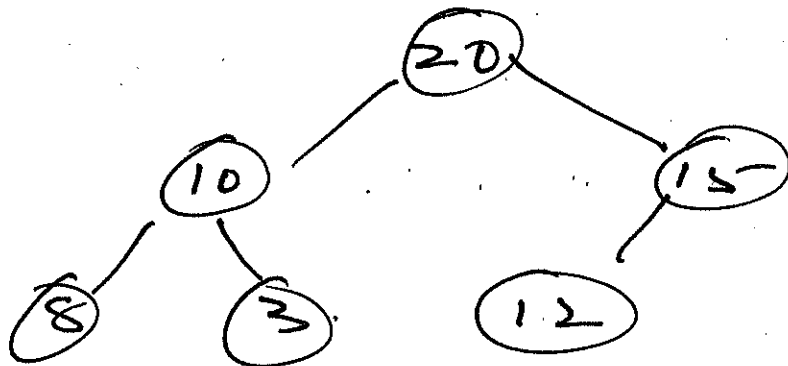
Ex. Perform Heapsort on

	1	2	3	4	5	6
A	10	8	12	20	3	15

Build Heap

heapSize

	1	2	3	4	5	6
A	20	10	15	8	3	12



Runtime:

$$T(n) = \Theta(n) + (n-1) \cdot \Theta(\log n) \\ = \Theta(n \log n)$$

Priority Queue

a Priority Queue (PQ) is an ADT that maintains a finite set S of elements (called records) each with an attribute called key.

record

$$x = (\underbrace{\cdot, \cdot, \cdot, \cdot}_{\text{satellite data}}, \text{key}) \in S$$

$$S = \{ \dots, x, \dots \}$$

$\text{key}[x]$ defines Priority of x in S .

states: $\mathcal{S} = \{ \dots, S, \dots \}$

A (max) P.Q. supports following operations,

- $\text{Insert}(S, x)$
inserts x into S .
- $\text{Max}(S)$:
return element with largest
key
- $\text{ExtractMax}(S)$
return & delete element
with largest key.
- $\text{IncreaseKey}(S, x, K)$:
change $\text{key}[x]$ to K if $K > \text{key}[x]$,
else do nothing

Heaps implemented P.Q.

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See Pseudo-code.

Keys only, no satellite data

Runtimes

HeapMaximum(A) : $\Theta(1)$

HeapDeleteMax(A) : $\Theta(\log n)$

HeapExtractMax(A) : $\Theta(\log n)$

HeapIncreaseKey(A, i, k) : $\Theta(\log n)$

HeapInsert(A, k) : $\Theta(\log n)$

Exercise

write

Min()

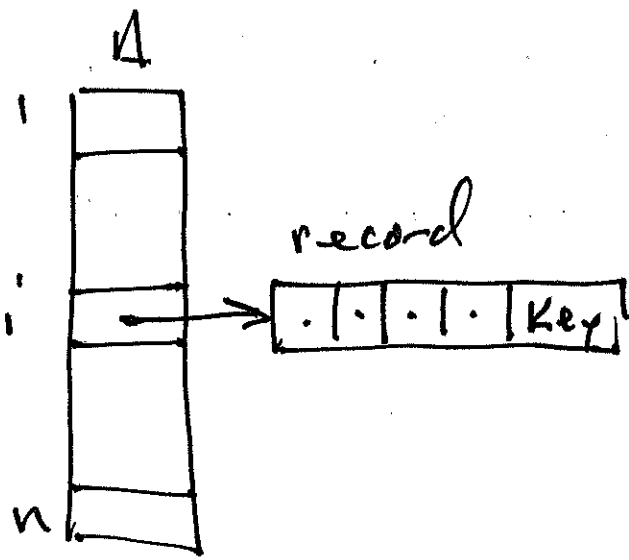
DeleteMin()

ExtractMin()

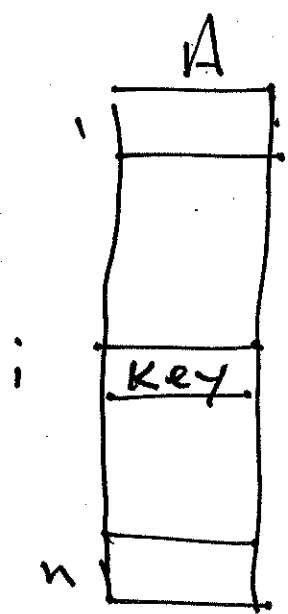
DecreaseKey()

for a min P.Q. implemented
as a min heap.

General Picture



Our Picture



Exercise

write P.Q. algorithm for
general Picture.

Exercise

implement a min P.Q. ADT
in C or C++ (or Python),

Read Chap 24:

SSSP in a weighted graph.

- Bellman Ford
- Dijkstra