

Pass: ext. 1 day

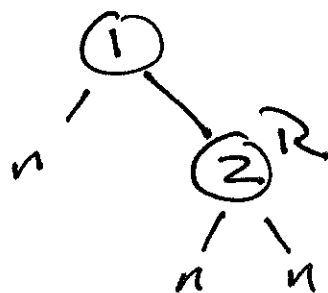
RBT insert:

Ex. Keys 1 2 3 4 5 6

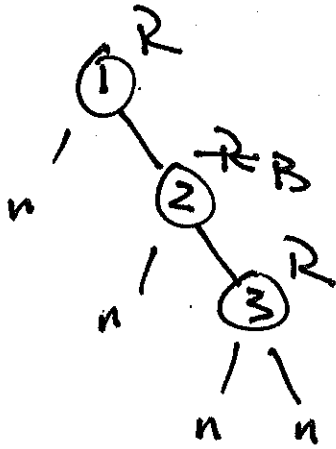
insert(1)



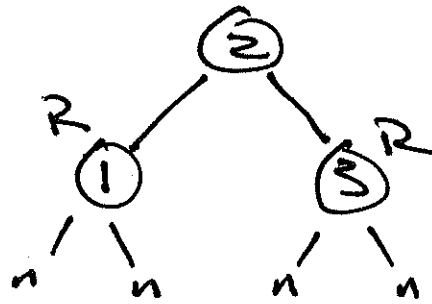
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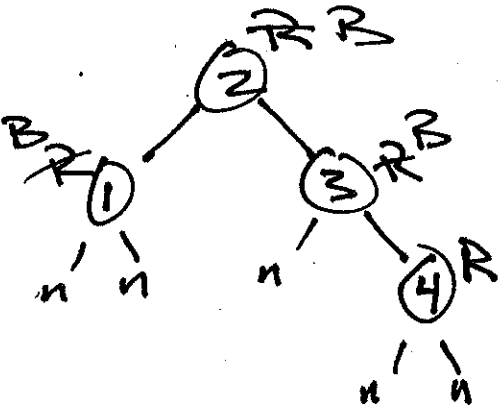
insert(3):



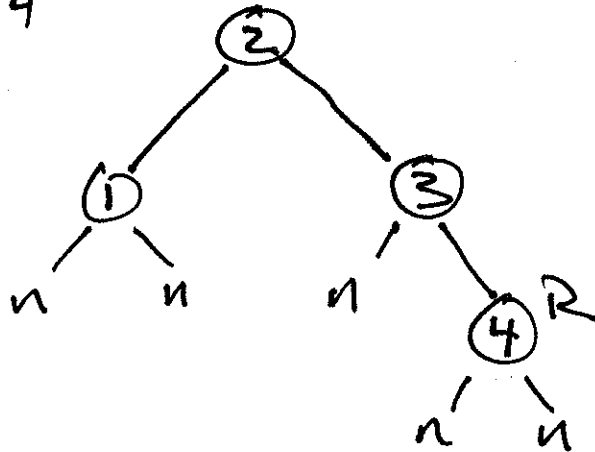
Case 6
→



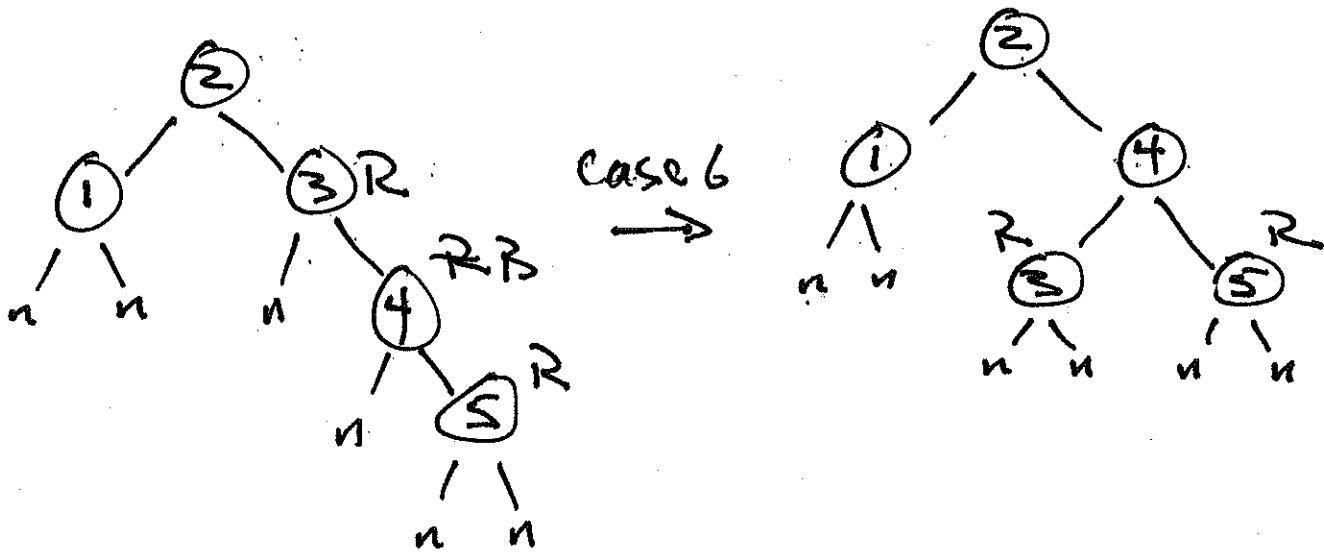
insert(4):



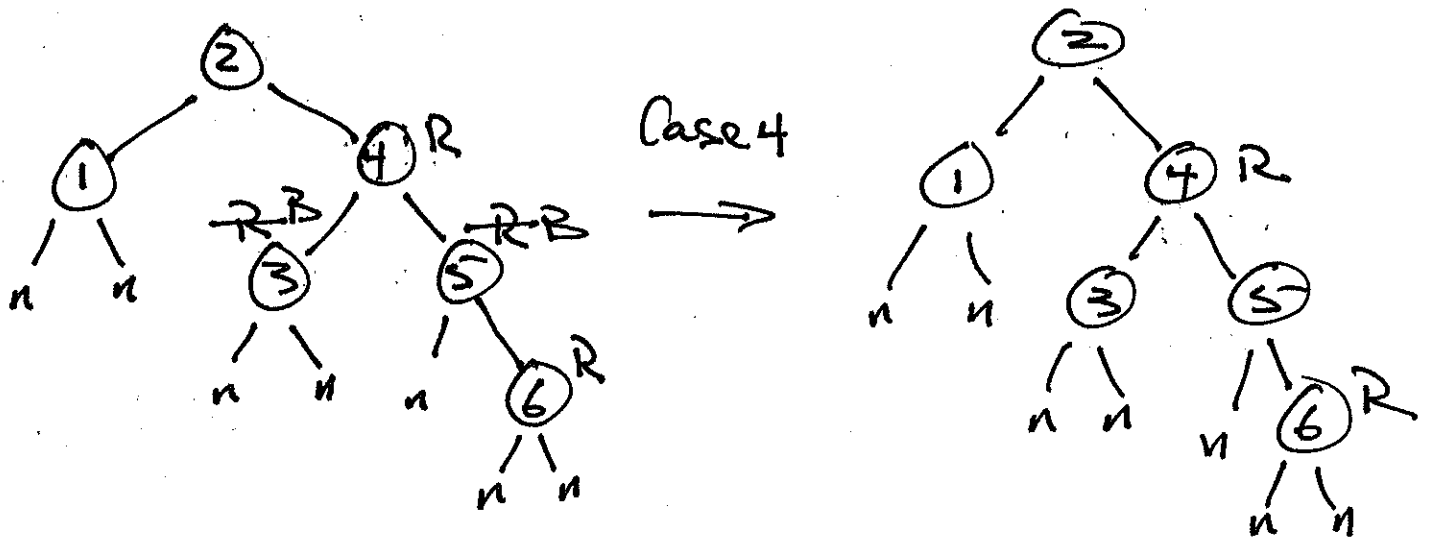
Case 4
→



insert(5)



insert(6)

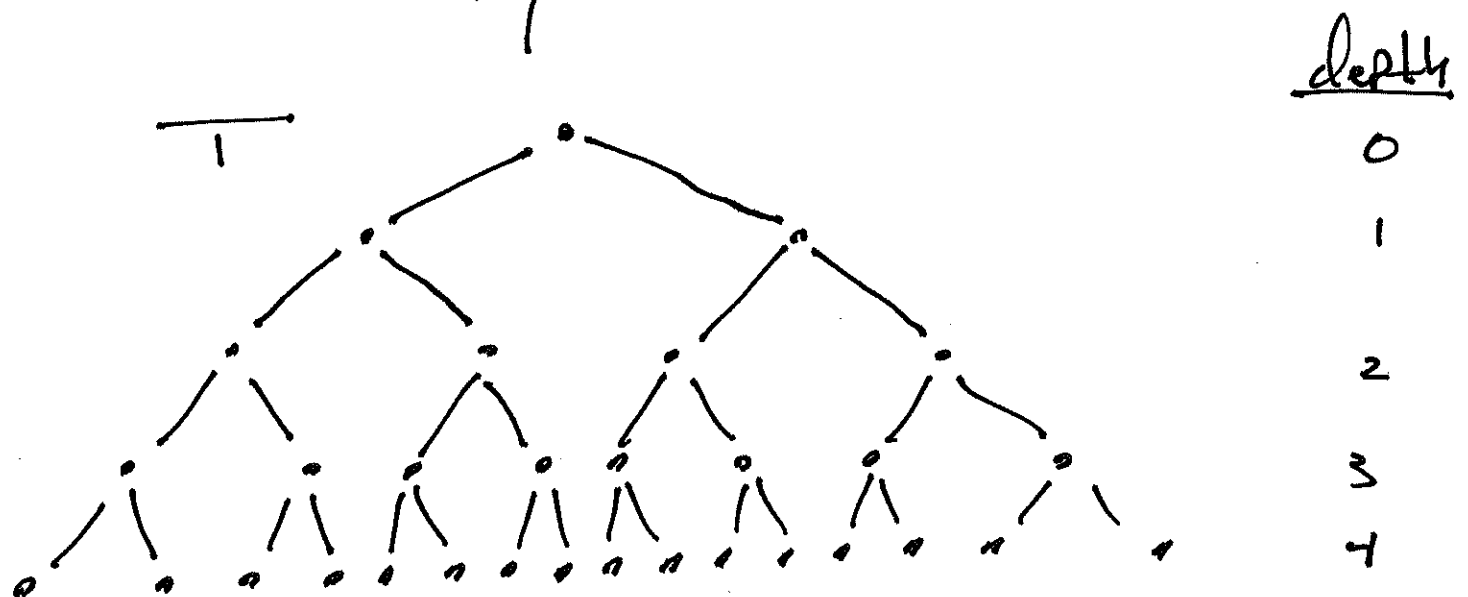


New data structure: binary heap

new ADT: Priority Queue

more on binary trees:

complete binary tree (CBT):



$$(\# \text{ nodes at depth } d) = 2^d \quad \text{if } T$$

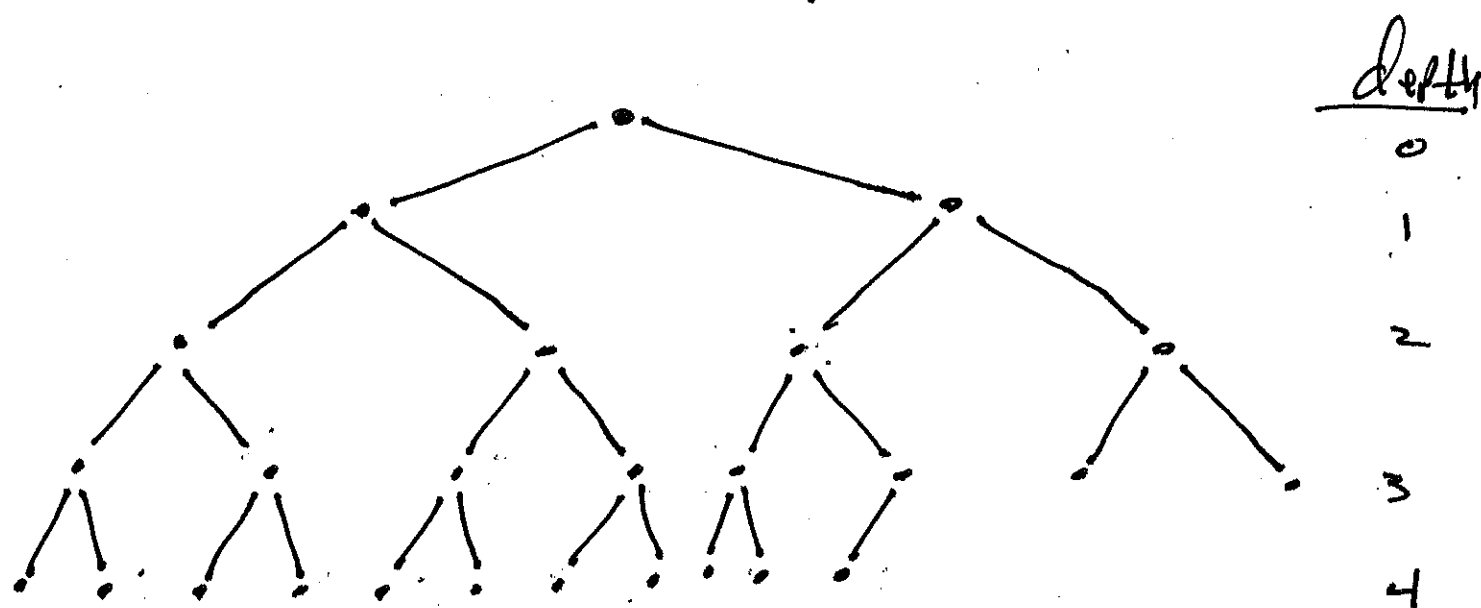
has height h , then

$$(\# \text{ nodes}) = \sum_{d=0}^{d=h} 2^d = \frac{2^{h+1} - 1}{2 - 1} = 2^{h+1} - 1$$

Thus the height of a CRT
on n nodes is

$$h = \lg(n+1) - 1$$

almost complete binary tree (ACBT)



The # of nodes n in an ACBT
of height h must satisfy

$$2^h - 1 < n \leq 2^{h+1} - 1$$

$$\therefore 2^h \leq n < 2^{h+1}$$

$$\therefore h \leq \lg(n) < h+1$$

$$\therefore h = \lfloor \lg n \rfloor$$

Recall Theorem

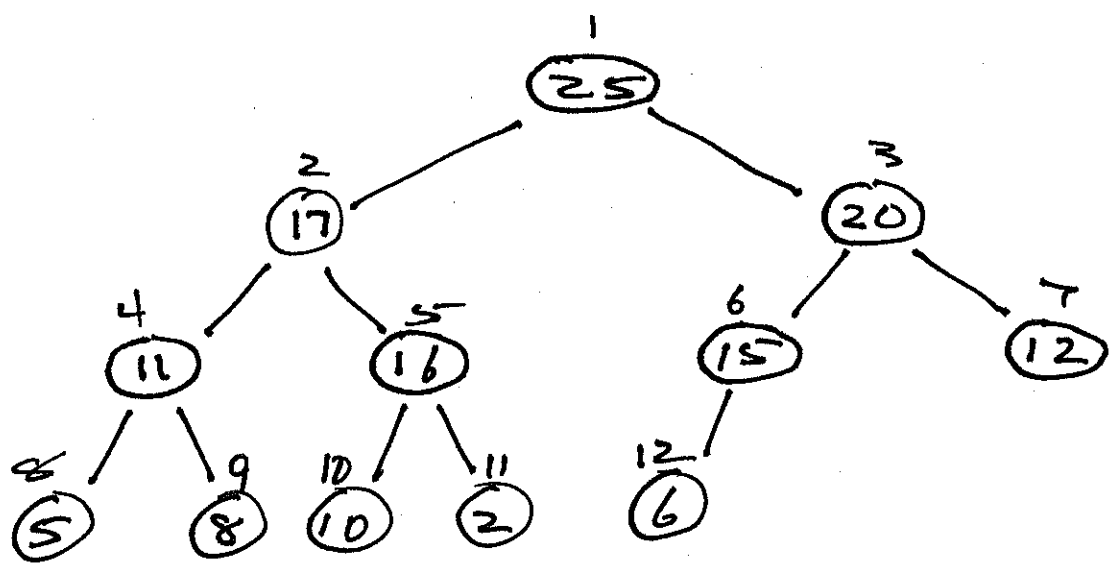
Any Binary Tree with n nodes and height h satisfies

$$h \geq \lfloor \lg n \rfloor$$

6.1 Heaps

A Binary Heap is an array object which represents an ACBT.
each node in tree corresponds to an array element. (some array elements may not correspond to nodes.)

EX



A

1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
25	17	20	11	16	15	12	5	8	10	2	6	0	0	0	0

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Array attributes:

- $\text{length}[A]$

- $\text{heapSize}[A]$

heap functions:

$$\text{Parent}(i) = \left\lfloor \frac{i}{2} \right\rfloor \quad (i \geq 2)$$

$$\text{left}(i) = 2i$$

$$\text{right}(i) = 2i + 1$$

Array values satisfy:

- max-heap Property: $A[\text{Parent}(i)] \geq A[i]$

or

- min-heap Property: $A[\text{Parent}(i)] \leq A[i]$

we assume our heaps are all max heaps.

Heap operations

- Heapify : maintain the heap property
- BuildHeap : create a heap from an un-ordered array
- Heapsort : sort an array
- Priority Queue ops :
 :
 :
 next time ..