

Ex. Selection Sort: $\Theta(n^2)$ Runtime
 Quick Sort: $\Theta(n \log n)$

which is better?

in fact $n \ln(n) = o(n^2)$ ✓

$$\begin{aligned} \lim_{n \rightarrow \infty} \left(\frac{n \ln n}{n^2} \right) &= \lim_{n \rightarrow \infty} \left(\frac{\ln n}{n} \right) \\ &= \lim_{n \rightarrow \infty} \left(\frac{1/n}{1} \right) = 0 \end{aligned}$$

Ex $n \ln(n) = \omega(n)$ ✓

$$\lim_{n \rightarrow \infty} \left(\frac{n \ln(n)}{n} \right) = \lim_{n \rightarrow \infty} (\ln(n)) = \infty$$

what about BFS & DFS?

• BFS

size of input instance

$$n = |V(G)| \text{ and } m = |E(G)|$$

Initialization $\Theta(3n+5) = \Theta(n)$

every $x \in V$ goes through queue at most once. $\therefore$ queue ops cost at most $\Theta(n)$. inner loop executes, in total

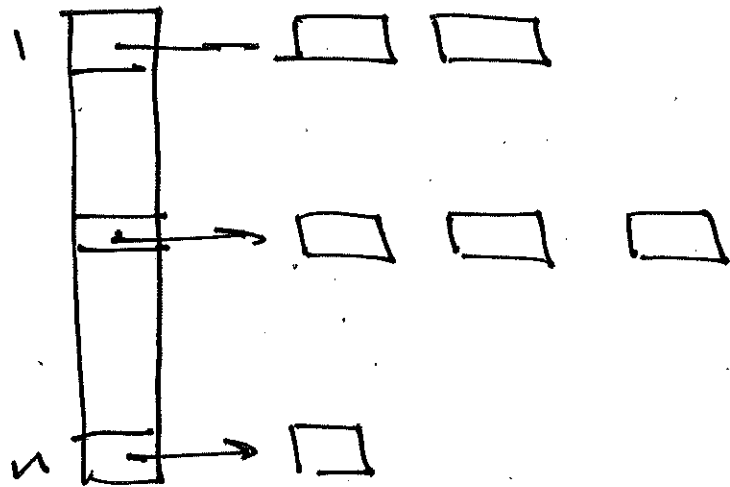
$$\sum_{x \in V} |\text{adj}(x)| = \sum_{x \in V} \text{deg}(x)$$

$$= \begin{cases} 2m & \text{undir.} \\ m & \text{dir.} \end{cases}$$

so that loop has total cost $\Theta(m)$.

$$\therefore \text{Total cost of BFS} = \mathcal{O}(n+m) \quad \boxed{3}$$

note: $n+m \approx$ a const. multiple of # of bytes needed to store the adj list representation



So $\mathcal{O}(n+m)$ is linear in size of input instance.

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DFS $n = |V|, m = |E|$

Initialization: $\Theta(n)$

main loop (other than visit): $\Theta(n)$

every vertex is visited exactly once, $\Rightarrow$ const cost.

ops in visit cost: $\Theta(n)$

for-loop in visit processes each adj list once, $\Rightarrow$ its cost is

$$\sum_{x \in V} |\text{adj}(x)| = \sum_{x \in V} \deg(x) = \begin{cases} 2m & \text{undir.} \\ m & \text{dir.} \end{cases}$$

$$= \Theta(m)$$

(total cost) = $\Theta(n+m)$

again linear in # of
bytes needed to store
an input instance.

ADT = in C++

Pass: ext $\geq$ days

A class in C++ is the
Program implementation of
an ADT.

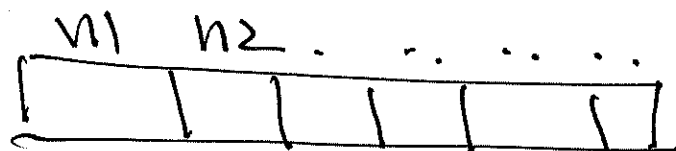
A struct in C++ is practically the same thing as a class.

major difference!

- struct members are default Public
- class " " " " Private

most C++ programmers use structs for PODs (member vars only, apart from constructors.)

A struct in C is more like an array with field names in place of indices



In C :

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```
struct Blah {
```

```
    int foo;
```

```
    int bar;
```

```
};
```

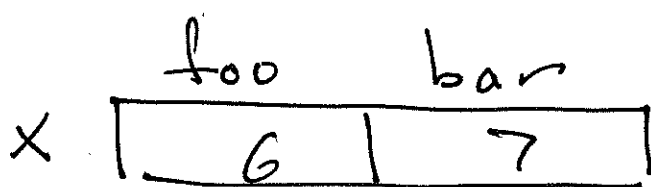
```
...
```

```
struct Blah x;
```

```
x.foo = 6
```

```
x.bar = 7
```

keyword
struct is
necessary.



stack
memory

In C++

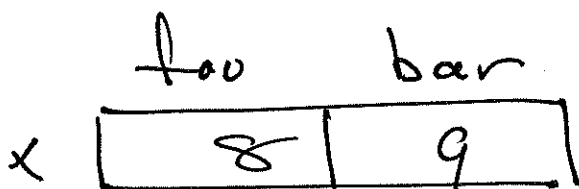
```
struct Blah {
    int foo;
    int bar;
```

```
    Blah (int a, int b) {
        foo = a;
        bar = b;
    }
```

```
}
```

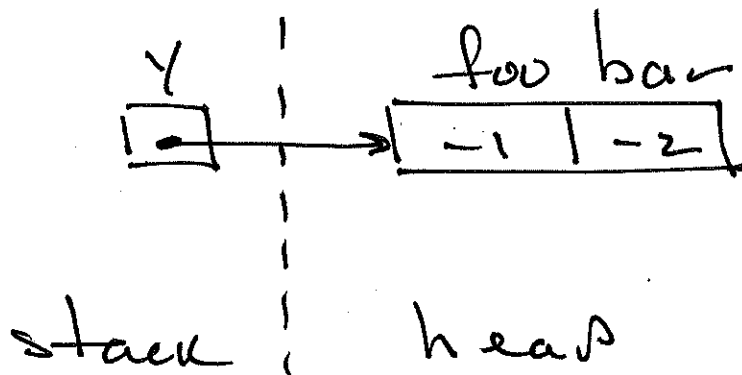
calls constructor

```
Blah x(8, 9);
```



} stack
memory

`Blah* y = new Blah(-1, -2);`



Basic skeleton for a class in C++:

```
class Blah {
```

```
private:
```

```
    // inner structs
```

```
    // member variables
```

```
    // helper functions
```

Public :

- // constructor (standard/copy)
- // destructor
- // member methods
- // overloaded operators

};

To build an ADT : 2 files

- Blah.h contains class defn and function prototypes.
- Blah.cpp contains function definitions

In Blah.h

class Blah {

Private:

// inner structs

// member var

// helper fn. Prototypes

Public:

// Prototype for

// const. & destruct,

// member methods

// overloaded operators

A class defines a namespace in C++. To access that namespace from outside the class defn, use the namespace resolution operator

::

In Blah.cpp

```
type Blah::fcn1 (... Params ...) {  
    // code  
}
```

Example: Queue.h
 .cpp
QueueTest.cpp
Makefile.