

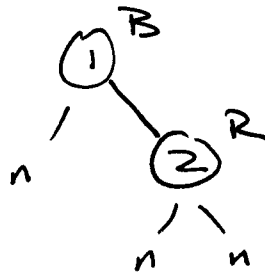
Ex RBT Insertion

Keys: 1, 2, 3, 4, 5

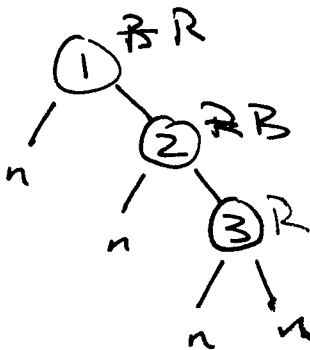
insert 1



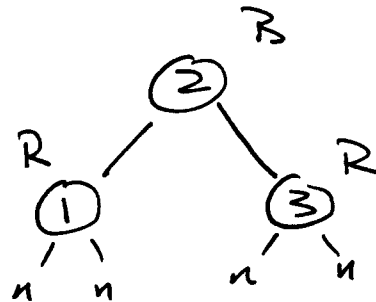
insert 2



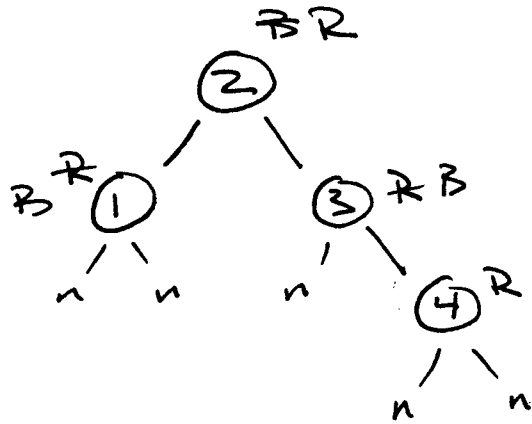
insert 3



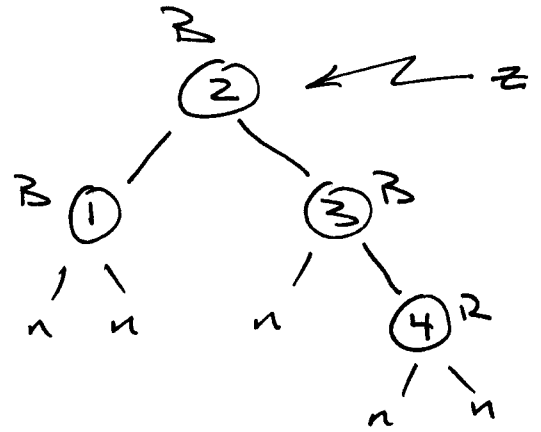
Case 6



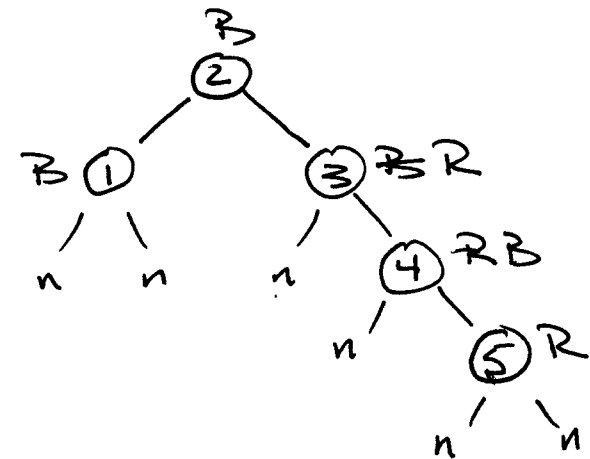
insert 4



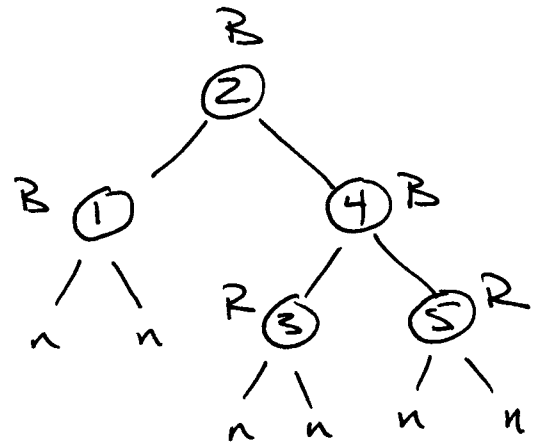
→
case 4



insert 5



→
case 6



chap. 6 : Heaps, Heapsort, Priority Queues.

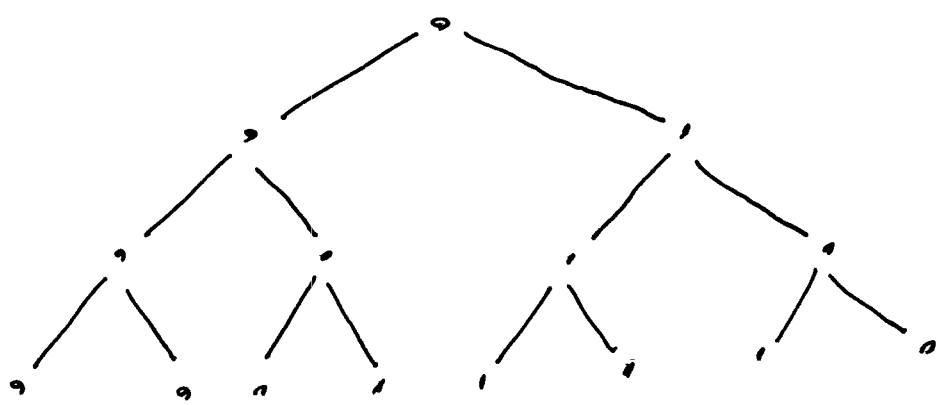
↑
Data
struct

↑
ADT

Recall

complete binary tree (CRT)

$h=3$
 $n=15$



<u>depth</u>	<u># nodes</u>
0	1
1	2
2	4
3	8

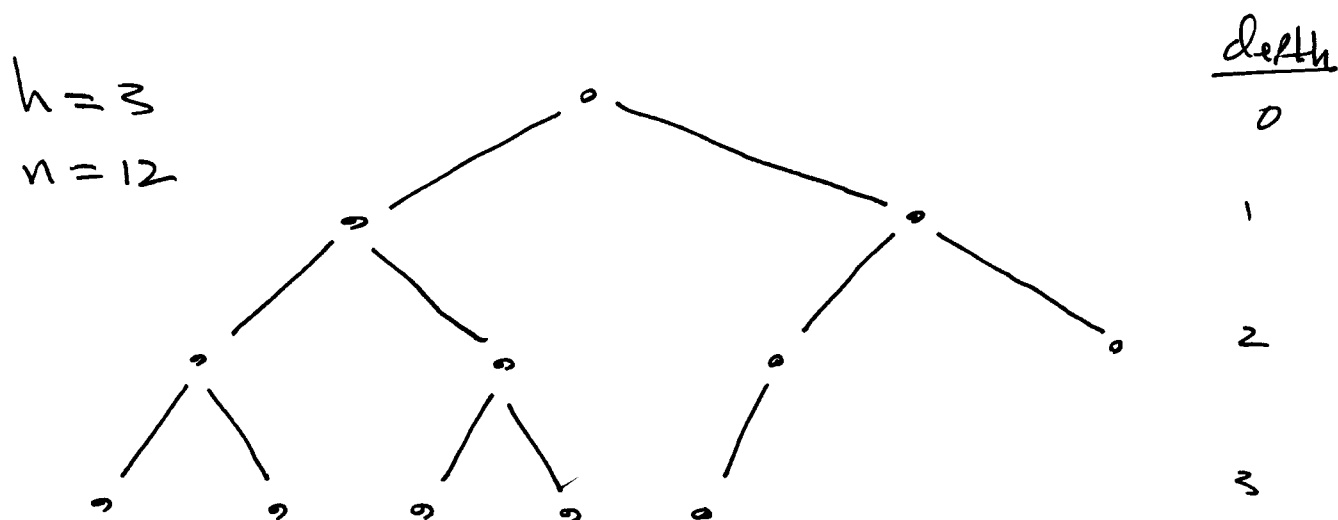
- all internal nodes have 2 children
- all leaves at same depth.

$$(\# \text{ nodes at depth } i) = 2^i$$

$$n = \sum_{d=0}^h 2^d = \frac{2^{h+1} - 1}{2 - 1} = 2^{h+1} - 1$$

almost complete binary tree (ACBT)

- all levels are filled, except (possibly) the last.
- deepest level is filled left to right



$$2^{(h-1)+1} - 1 < n \leq 2^{h+1} - 1$$

$$2^h - 1 < n \leq 2^{h+1} - 1$$

$$2^h \leq n < 2^{h+1}$$

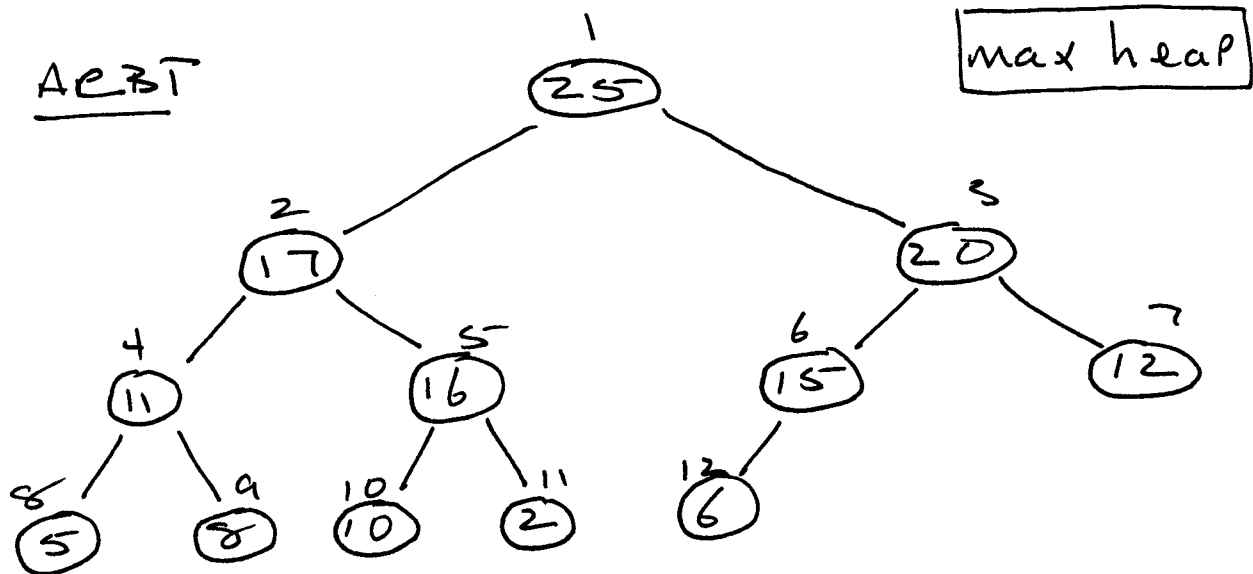
$$h \leq \lg(n) < h+1$$

$$\therefore h = \lfloor \lg(n) \rfloor$$

(6.1) heaps

A Binary Heap is an array that represents an ACBT, each node corresponds to an array element.
(some array elements might not correspond to nodes.)

Ex. ACBT



Array

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
A	25	17	20	11	16	15	12	5	8	10	2	6	•	•	•	•

$$\text{length}[A] = 16$$

$$\text{heapSize}[A] = 12$$

Array Attributes:

- length $[A]$
- heapSize $[A]$

Helper functions:

- $\text{Parent}(i) = \lfloor \frac{i}{2} \rfloor$ ($i \geq 2$)
- $\text{left}(i) = 2i$
- $\text{right}(i) = 2i + 1$

Heap Properties:

- max-heap Property: for $1 \leq i \leq \text{heapSize}$
 $A[\text{Parent}(i)] \geq A[i]$

- or
- min-heap Property: for $1 \leq i \leq \text{heapSize}$
 $A[\text{Parent}(i)] \leq A[i]$

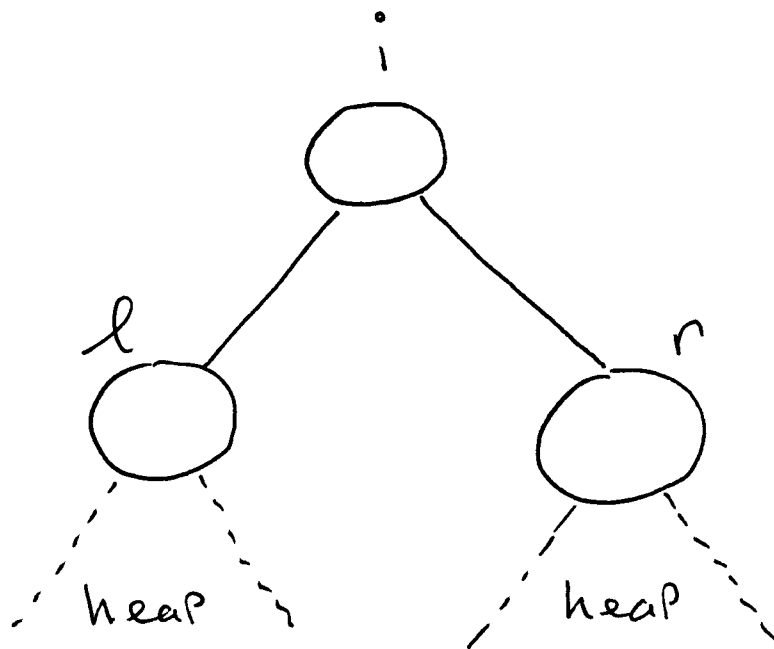
Algorithms:

- Heapify
- BuildHeap
- HeapSort
- P.Q. algorithms:
 - Insert
 - Max
 - ExtractMax
 - IncreaseKey

16.2) Heapify

Goal of Heapify(A, i) cause
max heap Prop. to be satisfied at
index i (and in left & right sub-
trees of i).

Precond: have valid heaps at subtrees
left(i) and right(i).



See Pseudocode on webpage.

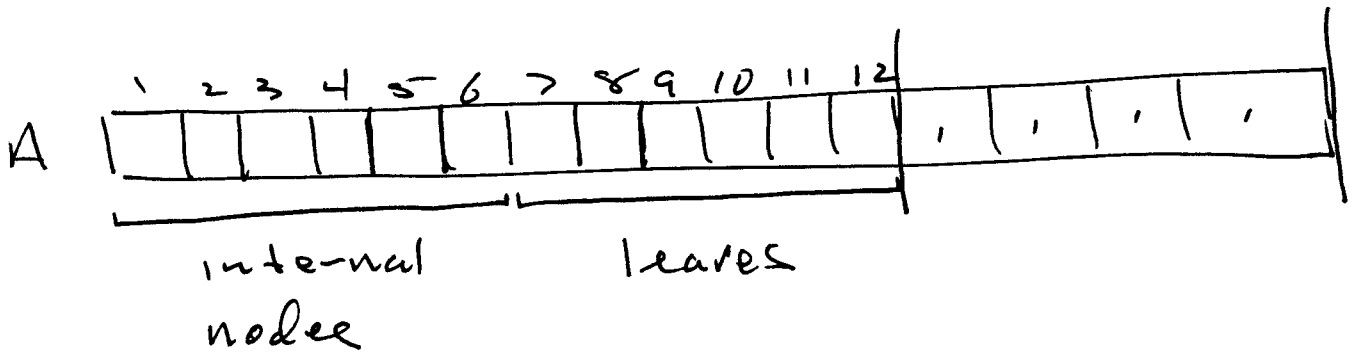
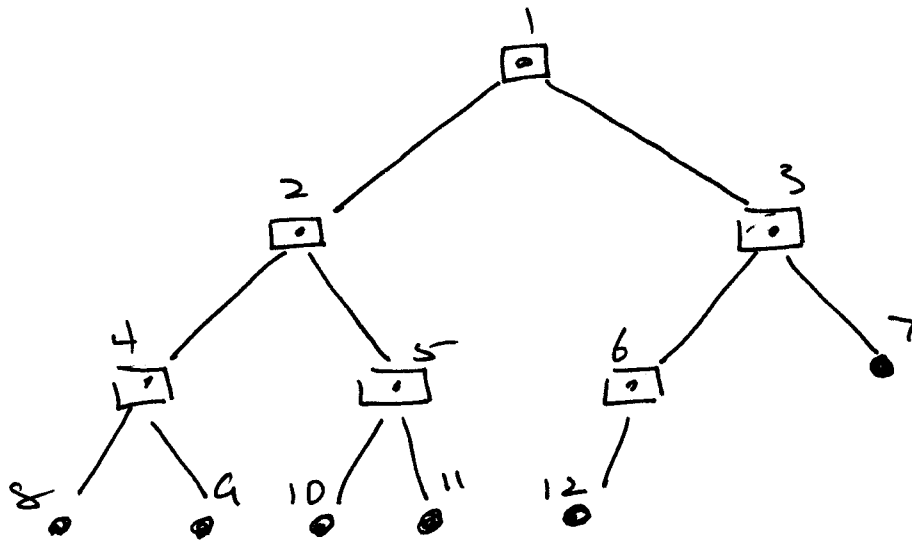
Let $T(n)$ = worst case runtime of heapify when run on a subtree (rooted at i) with n nodes, $n > 0$

$$\begin{aligned}
 T(n) &= (\text{height of sub-tree at } i) = \lfloor \lg n \rfloor \\
 &= \Theta(\log(n)) \\
 &= \Theta(h)
 \end{aligned}$$

(6.3) BuildHeap

19

Ex



see Pseudo-code .

Exercise : start with an un-ordered array, and run BuildHeap().