

- Quiz 2: next Tuesday
- Pa3: Posted
- Pa2: ext 1 last day

DFS continue

Predecessor Subgraph.

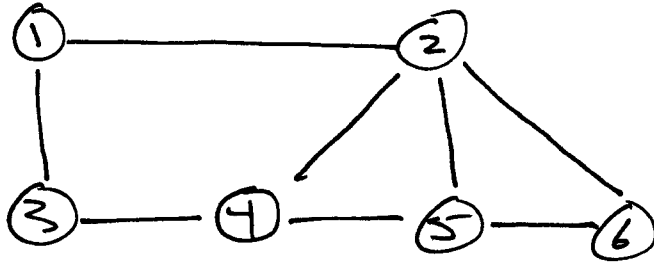
$$G_p = (V_p, E_p)$$

$$V_p = V(G)$$

$$E_p = \{ \underline{(p[x], x)} \mid p[x] \neq \text{nil} \}$$

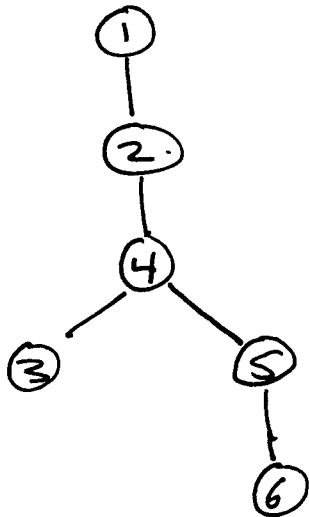
directed edge if G is a dir. graph
undirected edge if G a undir. graph

Ex.

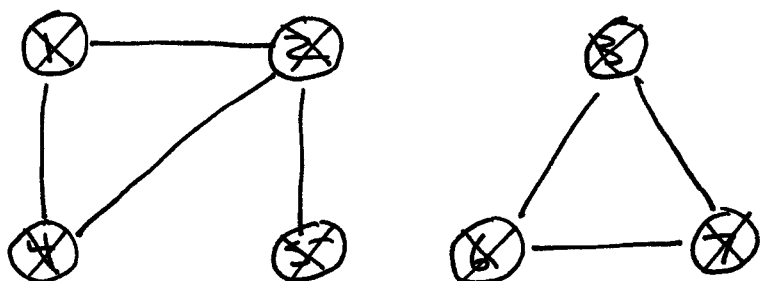


	adj	color	d	f	p
1	<u>2</u> <u>3</u>	w g b	1	12	4
2	<u>1</u> <u>4</u> <u>5</u> <u>6</u>	w g b	2	11	4 1
3	<u>1</u> <u>4</u>	w g b	4	5	4 4
4	<u>2</u> <u>3</u> <u>5</u>	w g b	3	10	4 2
5	<u>2</u> <u>4</u> <u>6</u>	w g b	6	9	4 4
6	<u>2</u> <u>5</u>	w g b	7	8	4 5

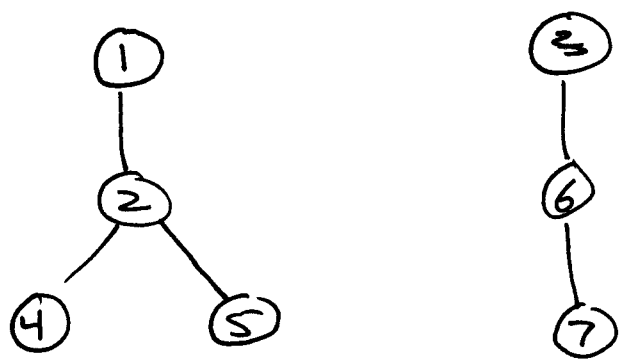
DFS
forest



Ex.



Dfs
forest



	adj	color	d	f	p
1	2 4	b	1	8	n
2	1 4 5	b	2	7	1
3	6 7	b	9	14	n
4	1 2	b	3	4	2
5	2	b	5	6	2
6	3 7	b	10	13	3
7	3 6	b	11	12	6

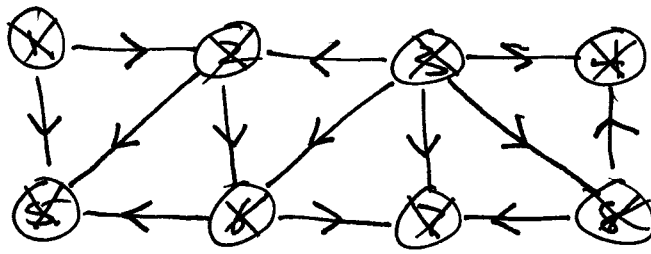
Exercise (G an undir. graph)
 add an attribute $cc[x]$ to
 each vertex x . Introduce a
 new variable $compNum$ in DFS
 (local to DFS, static over all
 recursive calls to $visit()$, like
 time.)

figure out what to do with
 $cc[x]$ and $compNum$ such that
 after $DFS(G)$, each $x \in V(G)$ has
 a cc -value and

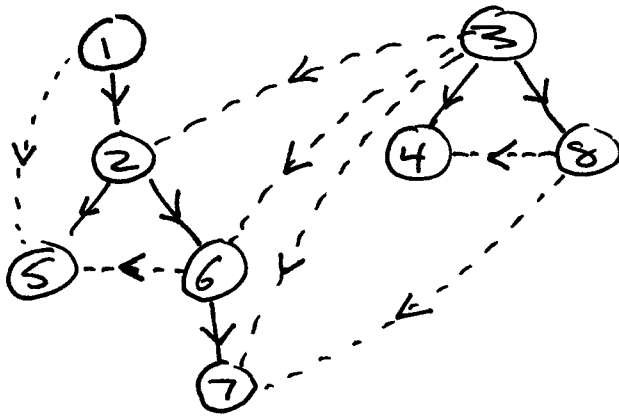
$$cc[x] == cc[y]$$

iff x, y lie in same connected
 component.

Ex.



DFS forest



Tree edges

(1, 2), (2, 5), (2, 6), (6, 7)

(3, 4), (3, 8)

Back edges $\emptyset$

Forward edges

(1, 5)

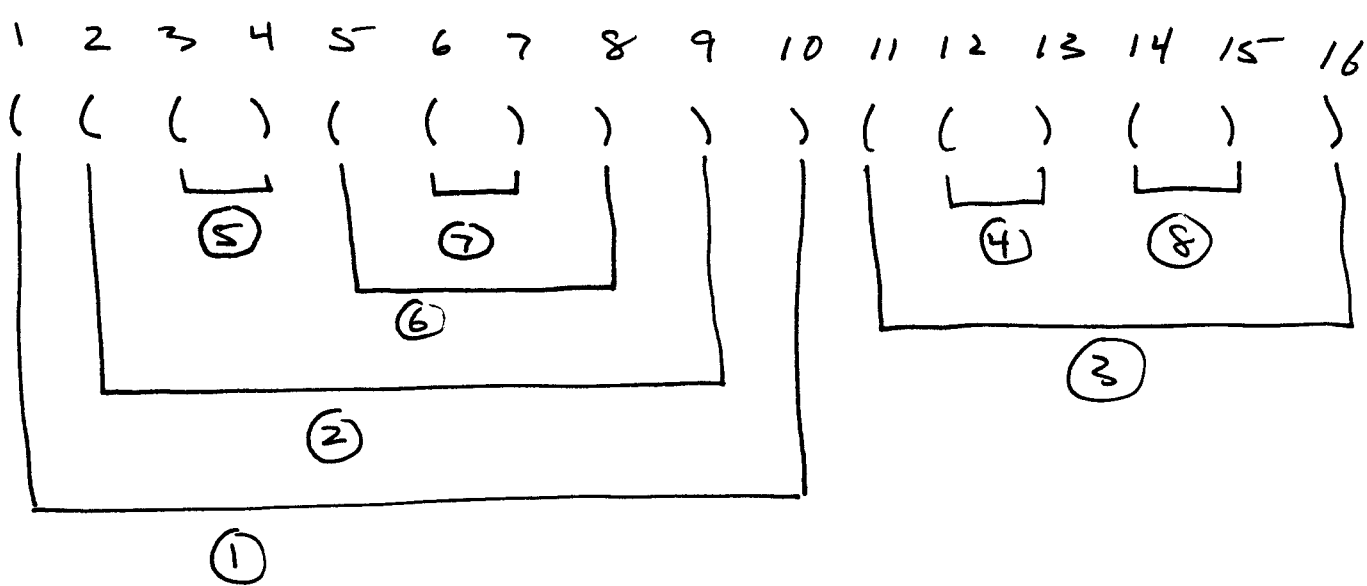
Cross edges

(6, 5), (8, 4)

(3, 4), (3, 6), (3, 7), (8, 7)

	adj	d	f	p
1	2 5	1	10	n
2	5 6	2	9	1
3	2 4 6 7 8	11	16	n
4		12	13	3
5		3	4	2
6	5 7	5	8	2
7		6	7	6
8	4 7	14	15	3

Parenthesis string



Thm

Suppose $x, y \in V(G)$ and $d[x] < d[y]$. Then either

$$\textcircled{1} \quad \begin{matrix} d[x] < f[x] < d[y] < f[y] \\ (\quad) & (\quad) \end{matrix} \quad \left\{ \begin{array}{l} \bullet y \text{ is not a} \\ \text{descendant} \\ \text{of } x \end{array} \right.$$

or

$$\textcircled{2} \quad \begin{matrix} d[x] < d[y] < f[y] < f[x] \\ (\quad (\quad) \quad) \\ \quad \quad \quad \underbrace{\hspace{2cm}} \\ \quad \quad \quad y \text{ gray} \\ \underbrace{\hspace{4cm}} \\ x \text{ gray} \end{matrix} \quad \left\{ \begin{array}{l} \bullet y \text{ disc. while } x \text{ gray} \\ \bullet y \text{ is a descendant} \\ \text{of } x \end{array} \right.$$

Thm (white path)

Let $x, y \in V(G)$. Then y is a descendant of x (i.e. ② holds) iff at time $d[x]$, G contains an x - y Path consisting of white vertices.

Edge classification

Run DFS. Then

- Tree edges: Belong to DFS forest G_p
- Back edges: join a vertex to one of its ancestors.
- Forward edges: join vertex to one of its descendants
- Cross edges: all else
 - tree to tree
 - cousin to cousin