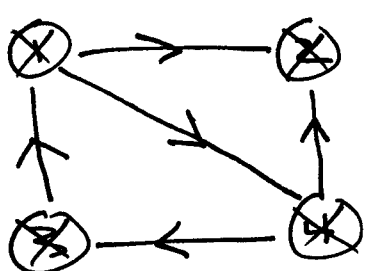


CSE 101 1-19-21

Pa2: ext. 2 days

Pa3: Posted tonight.

Ex.

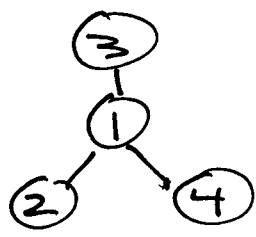


run BFS with $s = 3$

	adj	color	dist	Parent
1	2 4	black	1	3
2		black	2	1
3	1	black	0	nil
4	2 3	black	2	1

Q: $3 \neq 4$

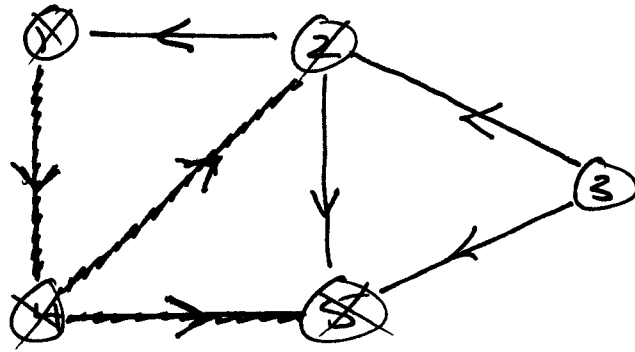
Tree



dist
0
1
2

Ex.

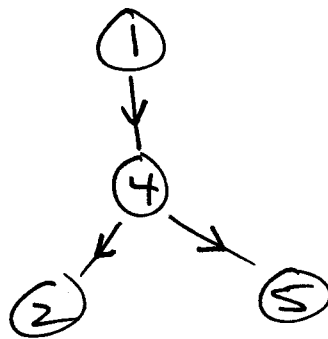
S=1



	adj	color	dist	Parent
1	4	b	0	nil
2	1 5	b	2	4
3	2 5	w	∞	nil
4	2 5	b	1	1
5		b	2	4

Q: $X \neq \neq$

Tree



dist

0

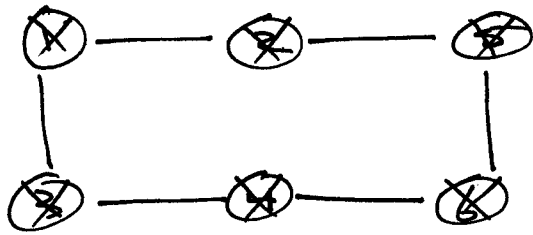
1

2

$\text{PrintPath}(G, s, x)$ prints a shortest $s-x$ Path in G .

Precondition: $\text{BFS}(G, s)$ must be run first.

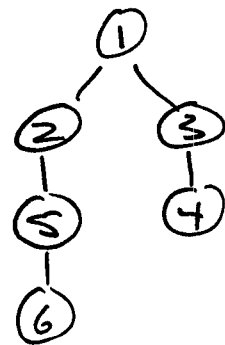
Ex.

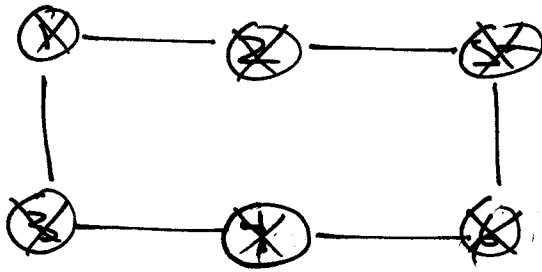


$s=1$ $Q: 1 \neq 2 \neq 3 \neq 4 \neq 6$

Shortest 1-6 Path:
1, 2, 5, 6

Tree



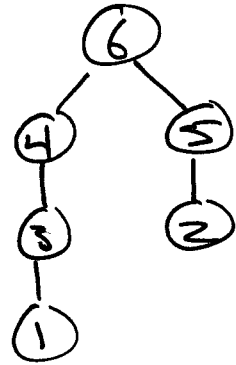
Ex.

$S = 6$ $Q: 6 \ 4 \ 3 \ 2 \ 1$

Shortest 6-1 Path:

6, 4, 3, 1

Tree



Exercise(s)

- Run BFS on many examples.
- defn The eccentricity of $x \in V(G)$, denoted $e(x)$ is the max. poss. distance from x .

$$e(x) = \max_{y \in V} d(x, y)$$

use Graph ADT ops to write a client fn. that computes $e(x)$, for any x .

- defn The radius of G is the minimum eccentricity, denoted $\text{rad}(G)$. The central vertices in G are those for which $e(x) = \text{rad}(G)$.

using Graph ADT, write a client fn that computes (and returns) $\text{rad}(G)$, and appends to L (2nd arg) the central vertices.

note

$$\text{rad}(G) = \min_{x \in V} e(x)$$

$$= \min_{x \in V} \left(\max_{y \in V} d(x, y) \right)$$

• defn The diameter of G 16
is the maximum eccentricity,
denoted $\text{diam}(G)$.

$$\text{diam}(G) = \max_{x \in V} e(x)$$

$$= \max_{x \in V} \left(\max_{y \in V} d(x, y) \right)$$

The Peripheral vertices are
those $x \in V$ for which
 $e(x) = \text{diam}(G)$.

Write a client fn. that finds
 $\text{diam}(G)$, and the set of
Peripheral vertices.

Recall: a g -aph H is a subgraph of a graph G iff

$$V(H) \subseteq V(G)$$

$$\text{and } E(H) \subseteq E(G)$$

Recall (G an undirected g -aph)
 G is called connected iff for all $x, y \in V(G)$: x is reachable from y (and y from x).

Defn

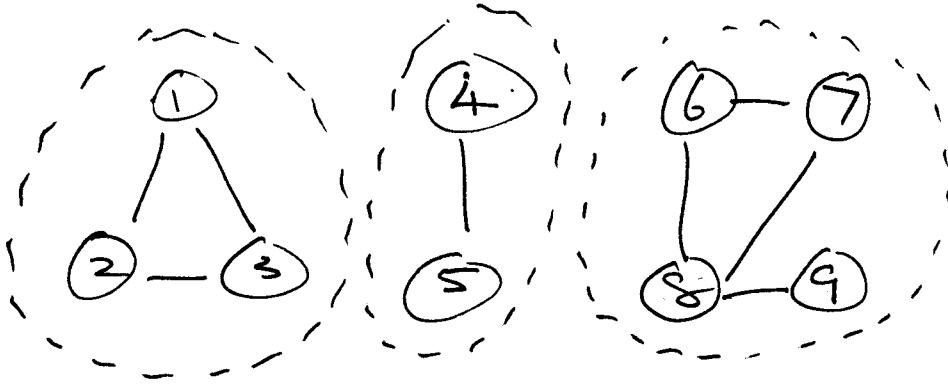
A connected component in G is a subgraph H s.t.

(1) H is connected.

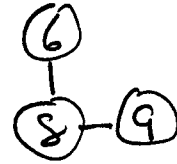
and

(2) H is maximal w.r.t. (1)

Ex.

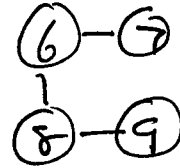


$$\bullet (\{6, 8, 9\}, \{68, 89\})$$



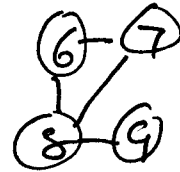
conn.
not
maximal

$$\bullet (\{6, 8, 9, 7\}, \{67, 68, 89\})$$



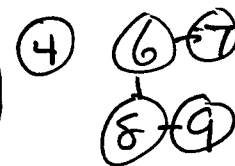
conn.
not
maximal

$$\bullet (\{6, 8, 9, 7\}, \{67, 68, 89, 78\})$$



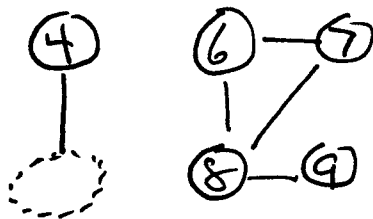
conn.
maximal

$$\bullet (\{6, 8, 9, 7, 4\}, \{67, 68, 89, 78\})$$



not conn.

$$\bullet (\{6, 8, 9, 4, 7\}, \{67, 68, 89, 78, 45\})$$



not a
graph !

Exercise

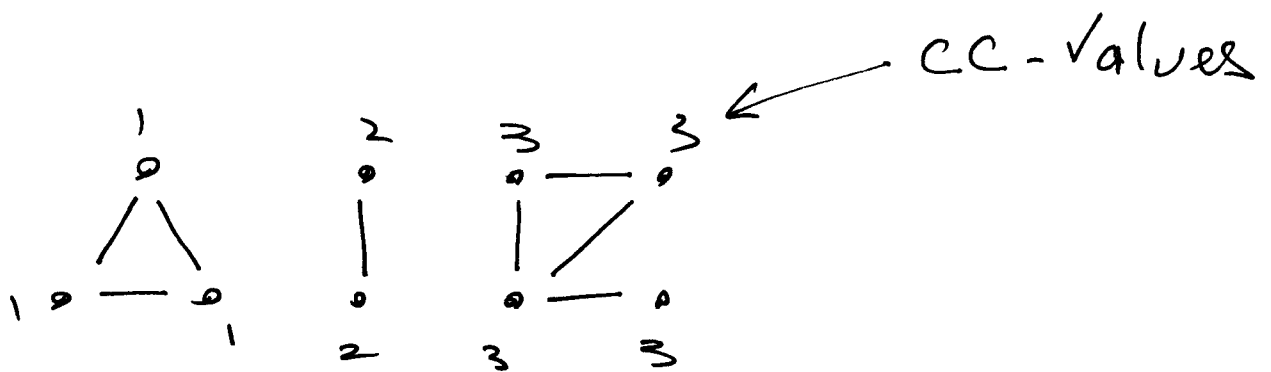
Add one new vertex attribute called $cc[x]$ for $x \in V(G)$.

Now write a client for Graph ADT that assigns value $cc[x]$ such $x \in V(G)$, such that

$$cc[x] = cc[y]$$

iff x, y lie in same Conn. Comp.

Ex.



Depth First Search (DFS)

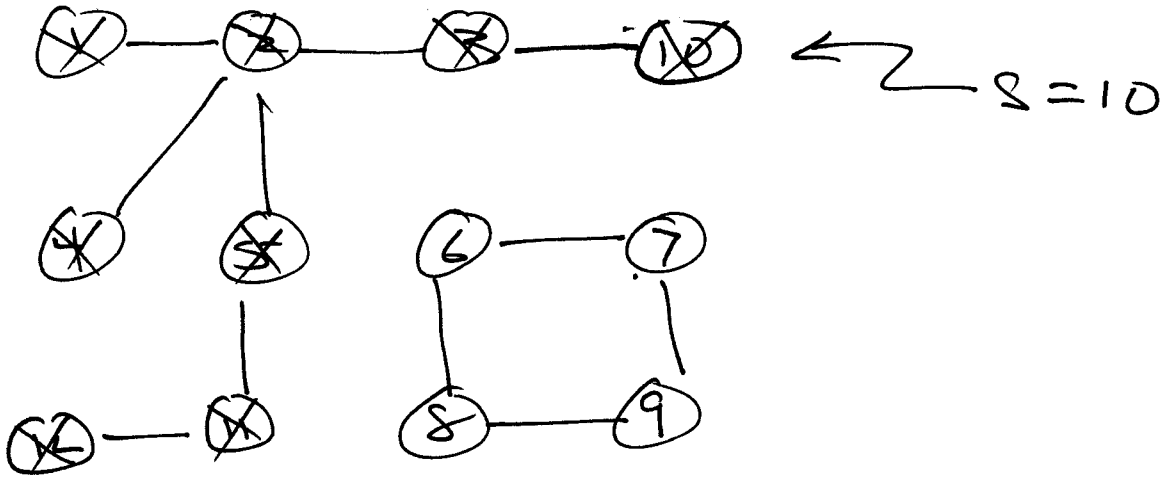
Vertex attributes

- $color[x]$: white, gray, black
- $P[x]$: Parent or Predecessor
- $d[x]$: discover time
- $f[x]$: finish time

uses a recursive subroutine
called $Visit(x)$

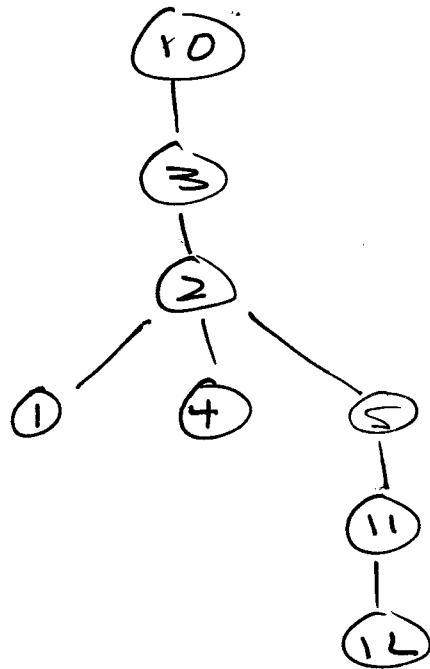
uses variable time (in DFS),
which is static over all recursive
calls to $Visit()$.

Ex,



Q: 10 3 7 1 4 5 11 12

Tree



dist

6

1

2

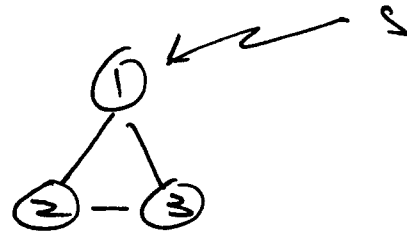
3

4

5

Graph ADT

client view



inside view

