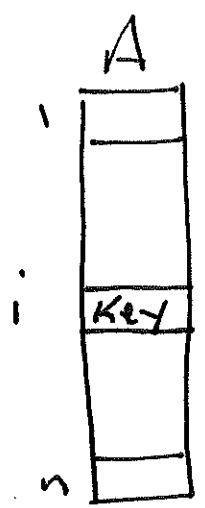


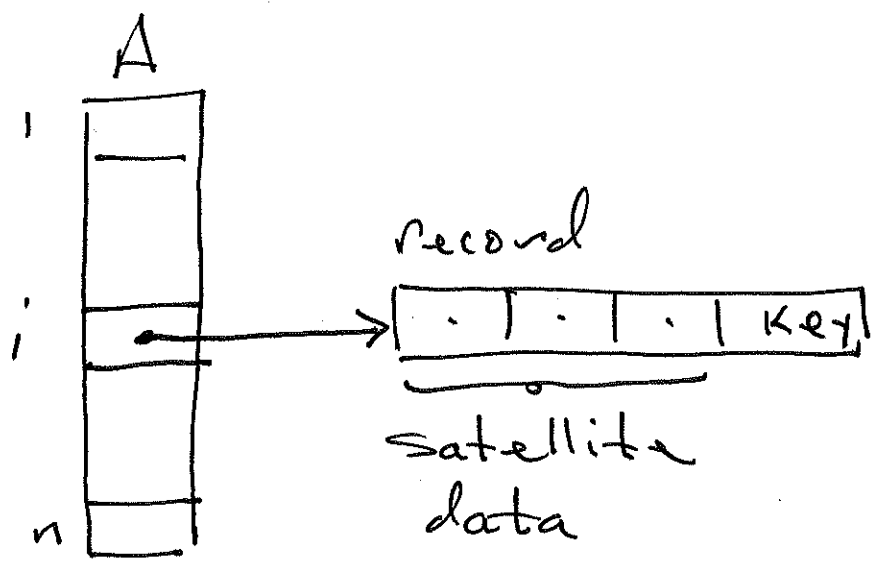


# Heap Algorithms

## our Picture:



## General Picture:



Exercise: Re-write algorithms in general Picture.

# SSSP Problem in a Weighted graph

ch. 24 of CLRS.

Defn g-graph

• a weighted <sup>g-graph</sup> is a g-graph  $G = (V, E)$  with a fcn.

$$w: E \rightarrow \mathbb{R}$$

called weight-fcn.

• weight of a path:  $x = v_0, v_1, \dots, v_k = y$  (P)

$$w(P) = \sum_{i=1}^k w(v_{i-1}, v_i)$$

• Shortest Path Weight  $x \rightarrow y$

$$\delta(x, y) = \begin{cases} \min\{w(P) \mid P \text{ is an } x-y \text{ Path}\} \\ \text{if such a path exists} \\ \infty \text{ if no such path exists.} \end{cases}$$

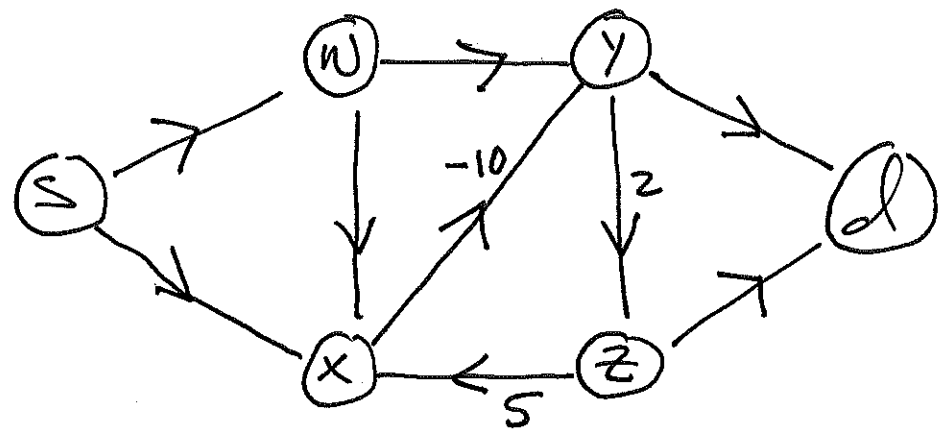
• a shortest  $x-y$  Path is an  $x-y$  path  $P$  s.t.  $w(P) = \delta(x, y)$ .

SSSP Problem

Given a weighted graph  $G$ , and a source  $s \in V(G)$ : (1) determine  $\delta(s, x)$  for all  $x \in V(G)$ , and (2) for all  $x$  with  $\delta(s, x) < \infty$ , find a shortest  $s-x$  path.

# Problem with SSSP:

Ex



cycle: x, y, z, x weight = -3

Dijkstra: outlaws negative weight edges.

Infrastructure for Dijkstra:

for each vertex:

- $P[x]$ : encodes a shortest path tree
- $d[x]$ : estimate of  $\delta(s, x)$
- Predecessor subgraph  $G_p = (V_p, E_p)$ 

$$V_p = \{x \in V(G) \mid P[x] \neq \text{nil} \text{ or } x = s\}$$

$$E_p = \{(P[x], x) \mid P[x] \neq \text{nil}\}$$

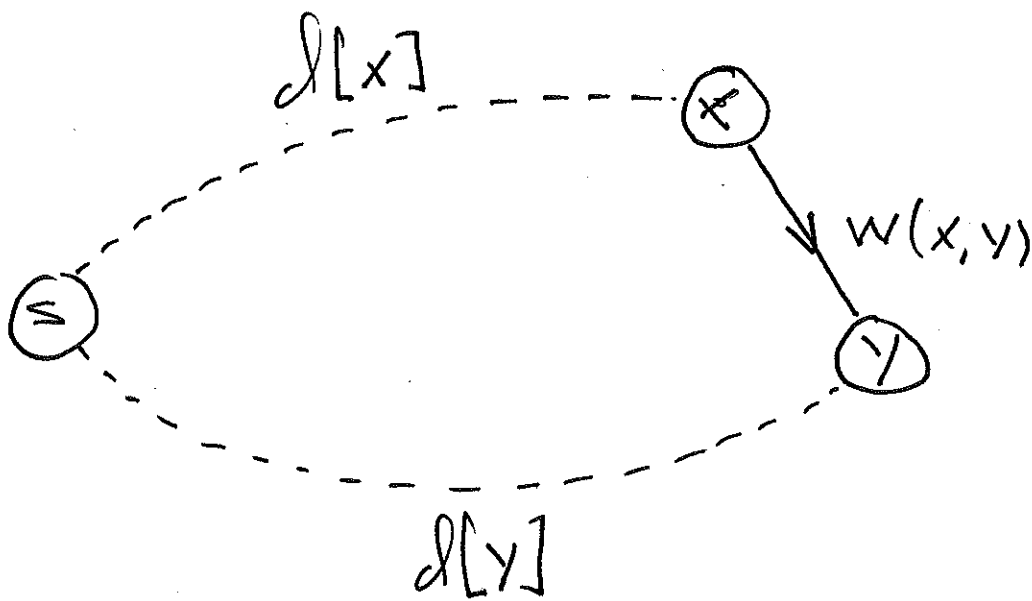
To print out a s.p. use

$\text{PrintPath}(G, s, x)$

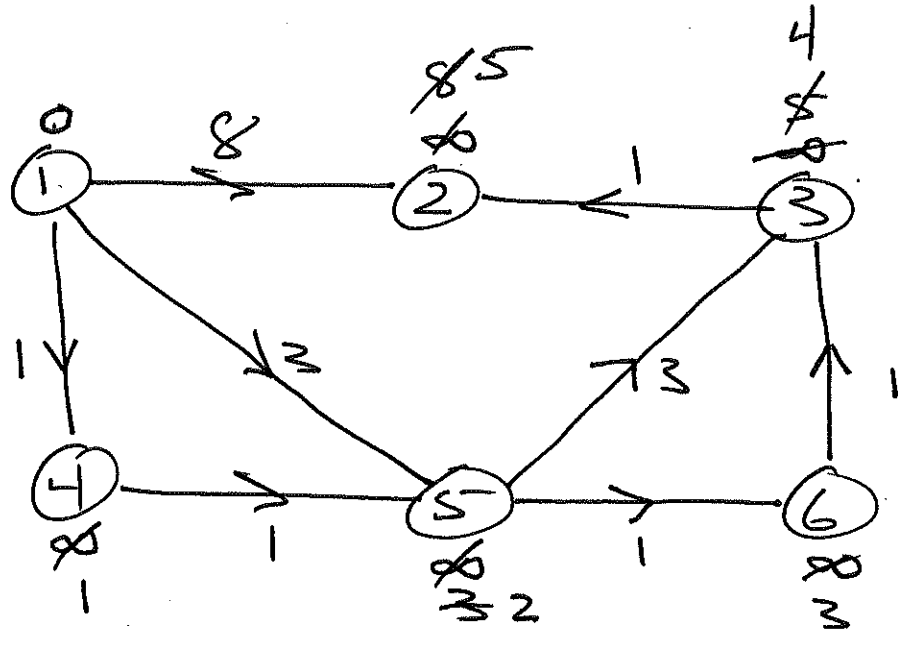
Two helper funcs.

• Initialize( $G, s$ )

• Relax( $x, y$ ) Pre:  $(x, y) \in E(G)$



Ex.  $S=1$



Extract order

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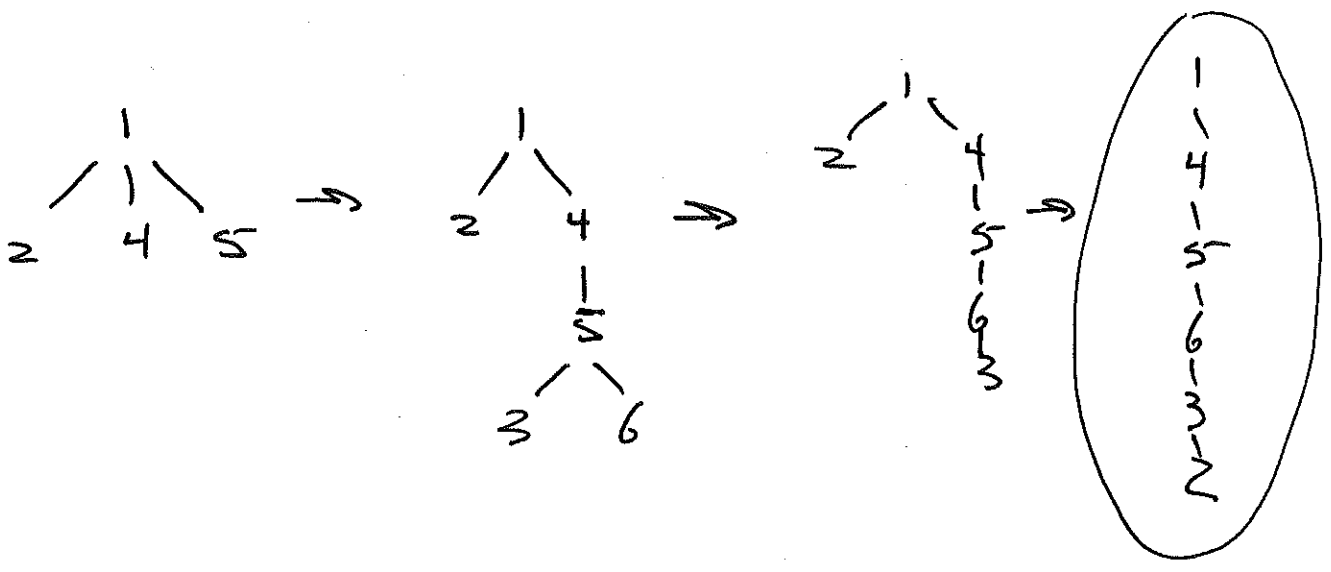
1  
4  
5  
6  
3  
2

PQ: ~~1~~ ~~2~~ ~~3~~ ~~4~~ ~~5~~ ~~6~~

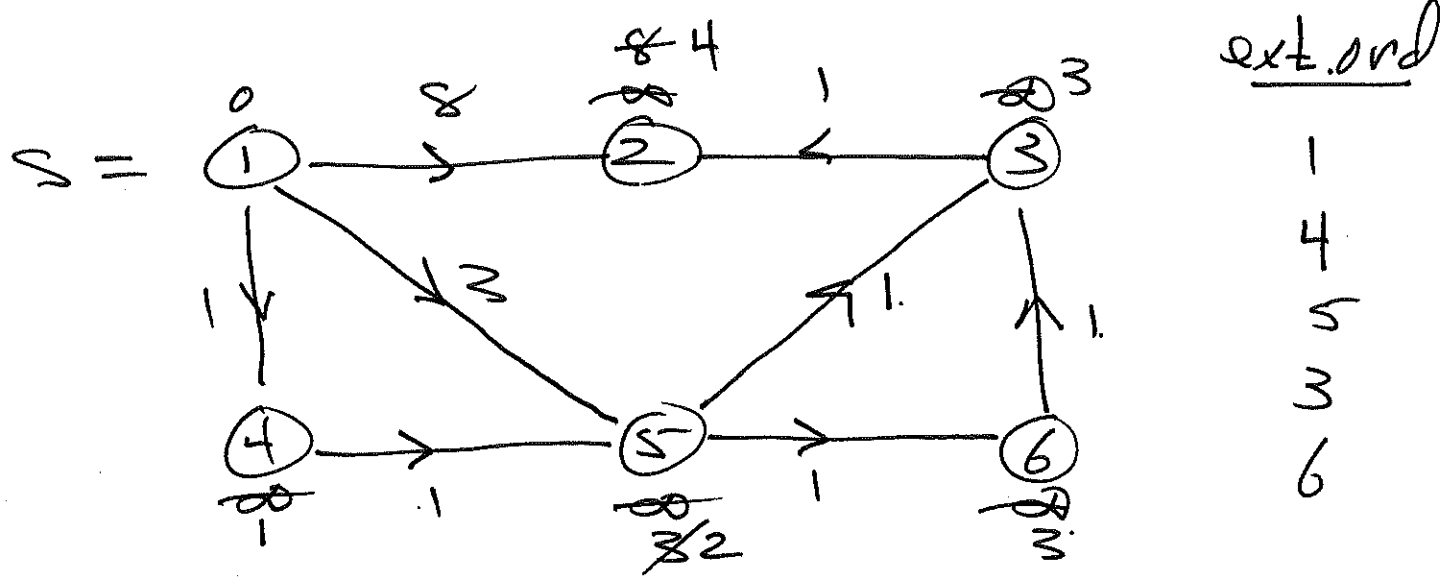
d: 0 ~~8~~5 ~~8~~4 1 ~~8~~2 3

P: n ~~x~~3 ~~8~~6 1 ~~x~~4 5

Tree:



Ex. if mult. id. keys, extract vertex with smallest label.



PQ: ~~1~~ ~~2~~ ~~3~~ ~~4~~ ~~5~~ ~~6~~

d: 0 ~~4~~ 3 1 ~~2~~ 3

P: n ~~3~~ 5 1 ~~4~~ 5

Tree!

