

CSE 101-02 5-12-22

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- mid 2: Thurs May 19
- Pag: ext 1 day (Tue.)

• Pag: Question

from-string constructor

Dictionary ADT (also called Map)

state: a set of ordered Pairs

(key, value)

i.e. a state is

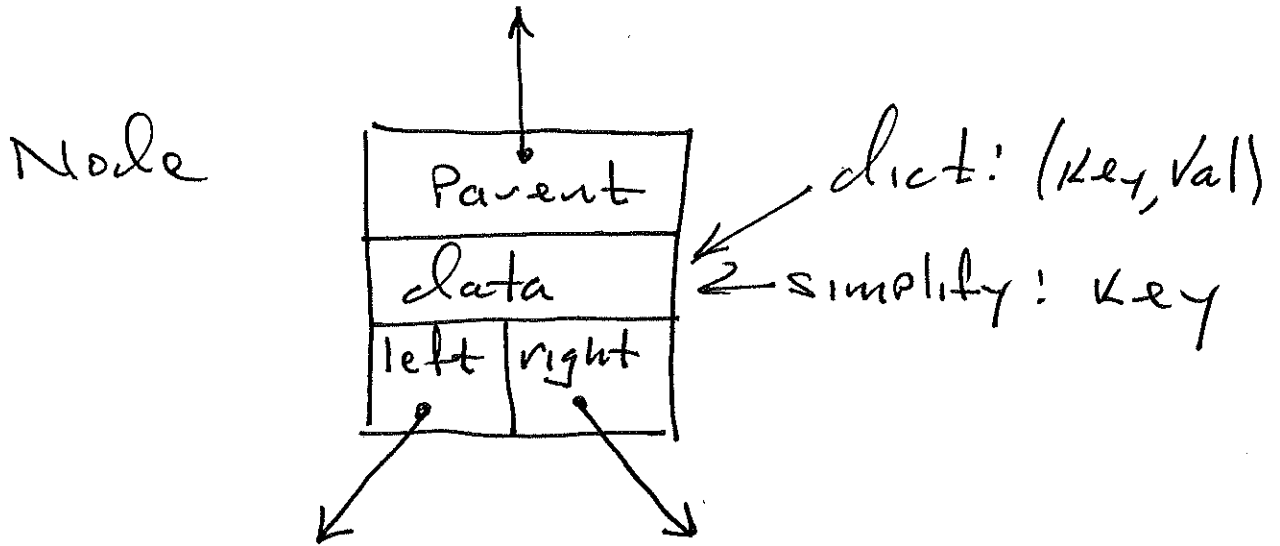
$$\{(k_1, v_1), (k_2, v_2), \dots, (k_n, v_n)\}$$

option: keys are unique

operations: main 3

- $\text{insert}(k, v)$: inserts a new pair in dictionary. Pre: k does not exist in dict. ($\text{lookup} = \text{nil}$)
- $\text{lookup}(k)$: returns v s.t. (k, v) is in dict., or nil if no such pair exists.
- $\text{delete}(k)$: removes pair (k, v) .
Pre: $\text{lookup}(k) \neq \text{nil}$.

Binary Search Trees (BST)



BST Properties:

Let x, y be nodes.

(1) if y is in the left subtree of x

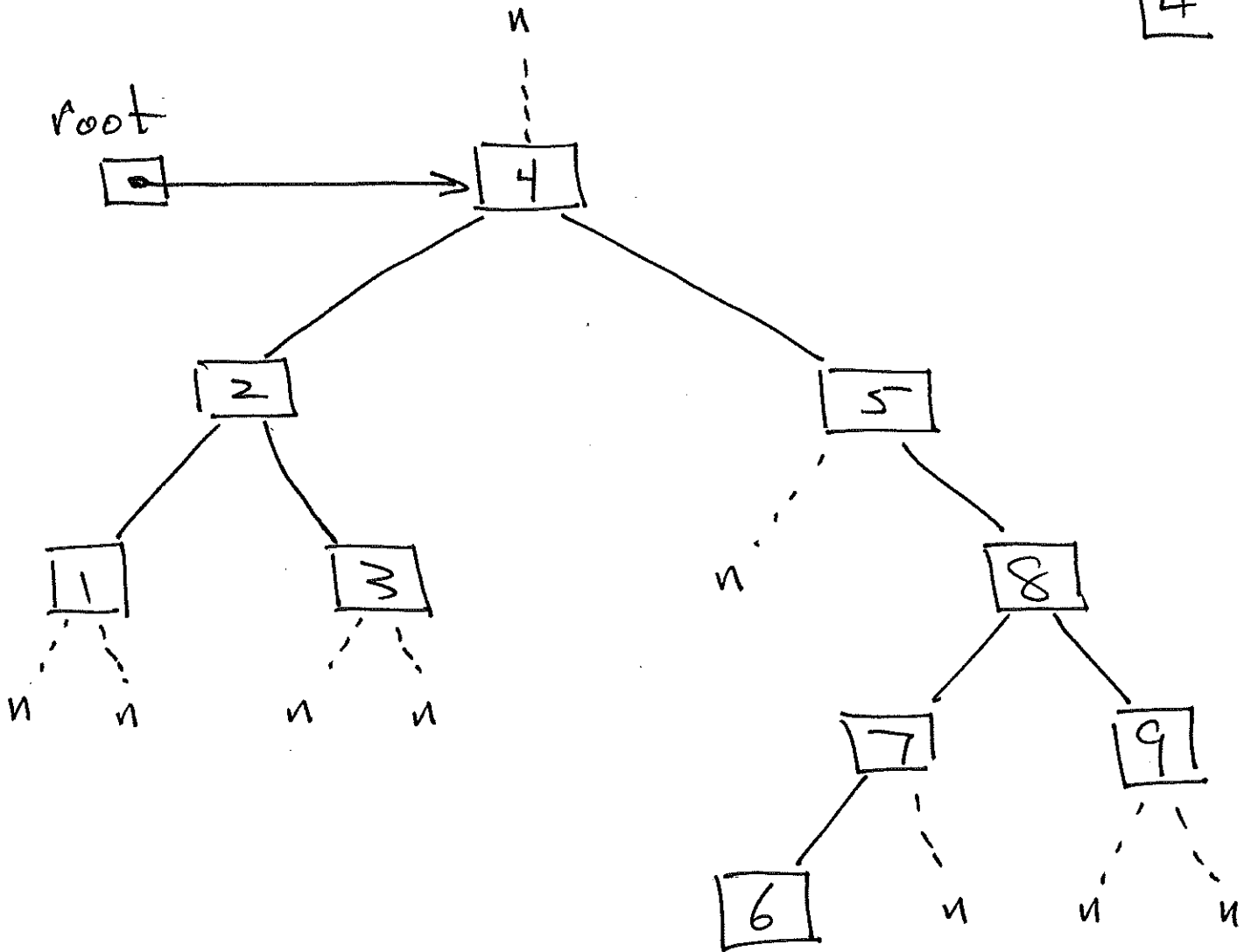
then $\text{key}[y] \leq \text{key}[x]$

(2) if y is in the right subtree of x

then $\text{key}[x] \leq \text{key}[y]$.

Ex

4



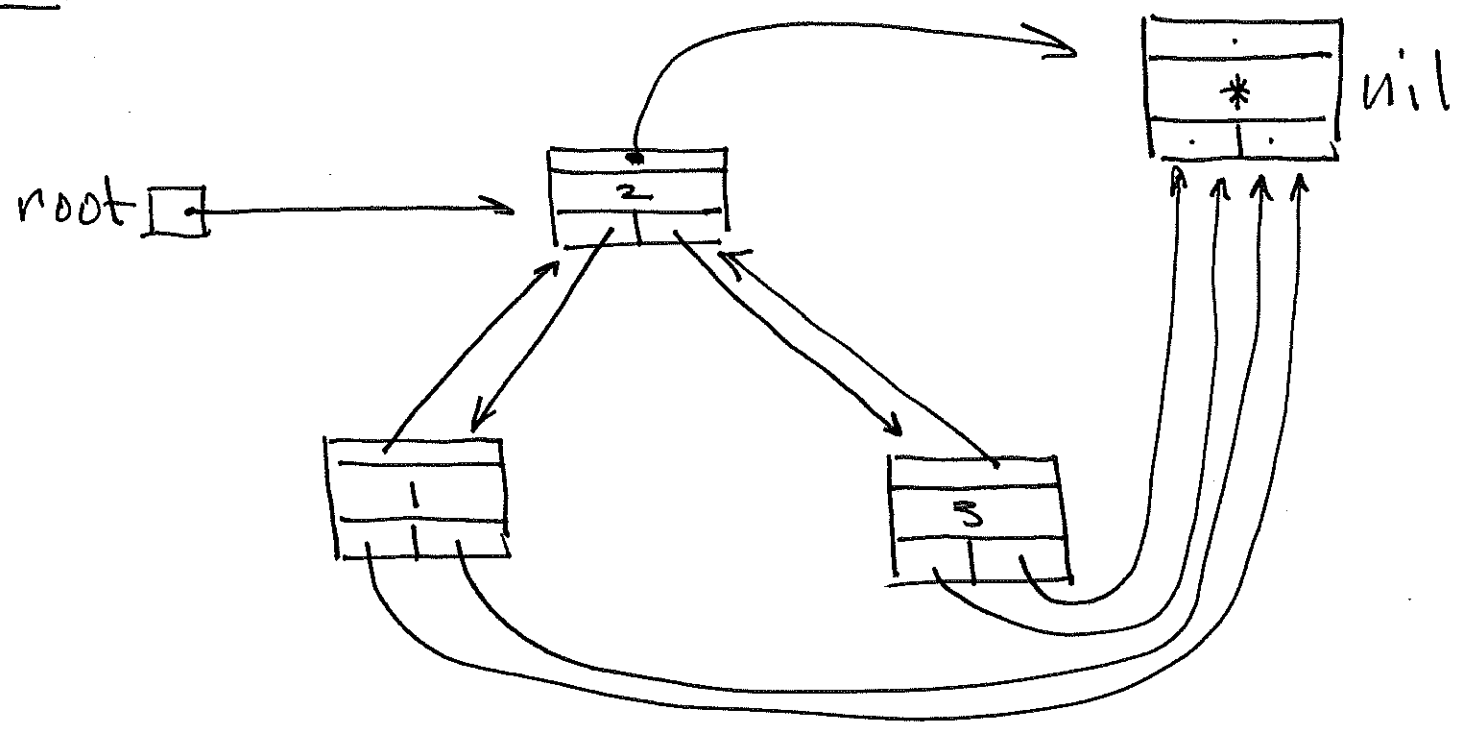
A leaf is a node with no children

An internal node is a non-leaf

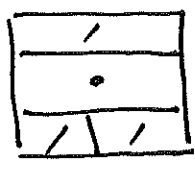
implementation option:

use a dummy (sentinel) node called nil

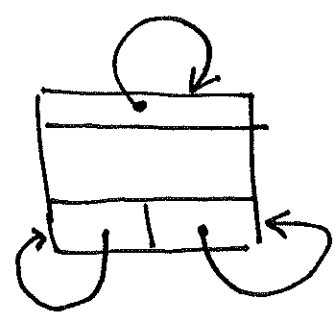
Ex.



could have



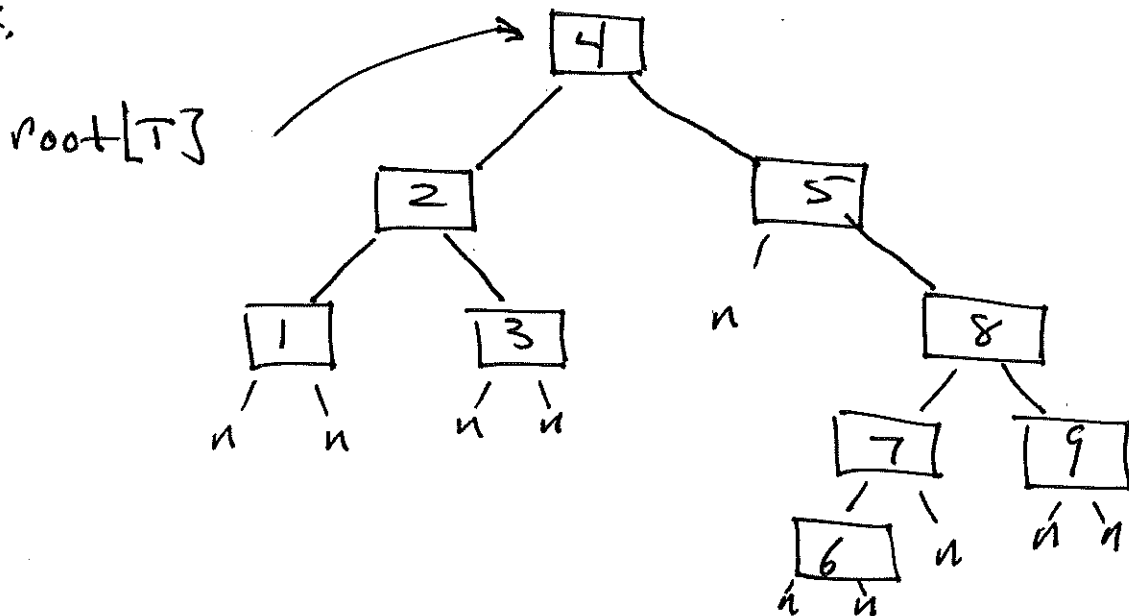
or



Tree Traversals

- InOrderTreeWalk(x) ✓
- PreOrderTreeWalk(x)
- PostOrderTreeWalk(x)

Ex.



InOrderTreeWalk($\text{root}[T]$)

1 2 3 4 5 6 7 8 9

On same example :

PreOrderTreeWalk($\text{root}[T]$) :

4 2 1 3 5 8 7 6 9

PostOrderTreeWalk($\text{root}[T]$) :

1 3 2 6 7 9 8 5 4

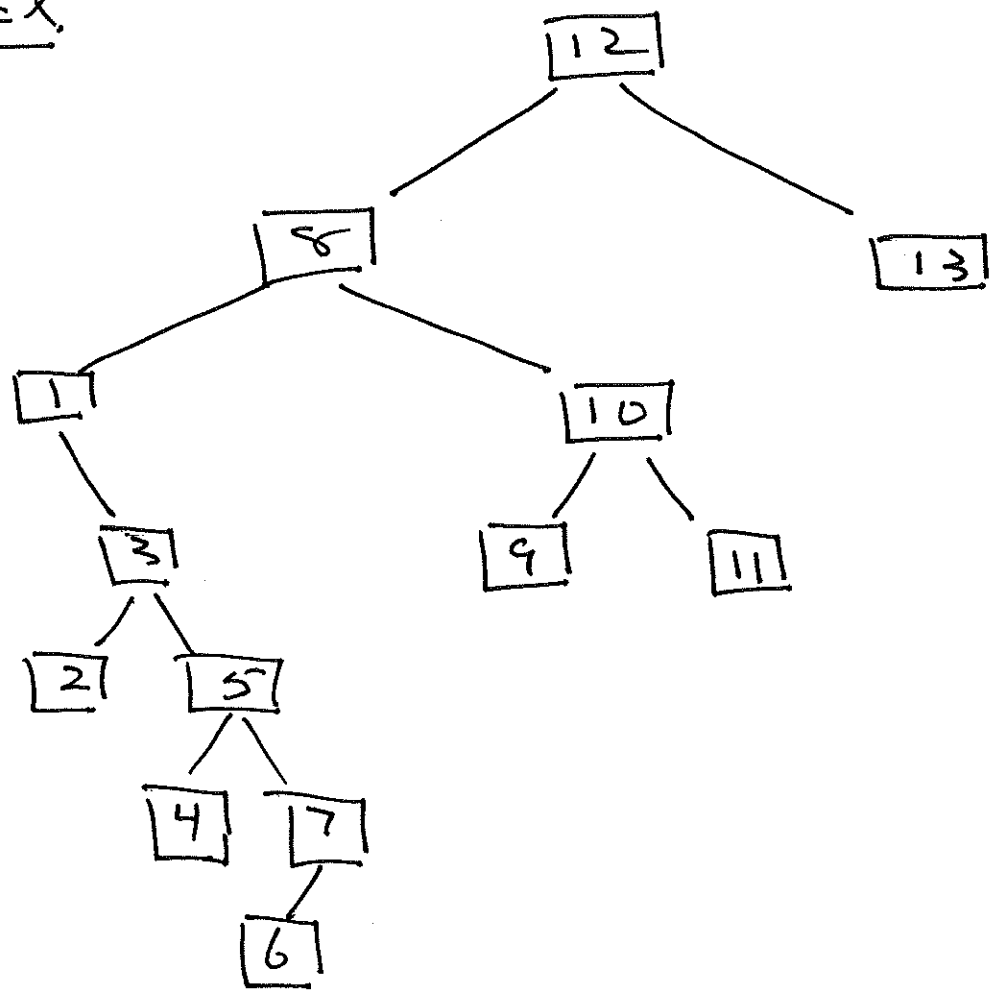
Queries

- $\text{TreeSearch}(x, k)$ ↙ node
- $\text{TreeMinimum}(x)$
- $\text{TreeMaximum}(x)$

Defn

The Successor of a node x is the next node after x to be processed in an in order tree walk.

Ex.



IOTW: 1 2 3 4 5 6 7 8 9 10 11 12 13

$$\text{Successor}(1) = 2$$

$$\text{Successor}(8) = 9$$

$$\text{Successor}(7) = 8$$

2 cases for successor:

case 1: x has a right child

Successor of x is the leftmost descendant of that child.

case 2: x has no right child

Successor of x is the lowest ancestor of x whose left child is also an ancestor of x .

Exercise

- Define Predecessor of x .
- write algorithm Predecessor(x).

Runtimes:

- Traversals: $\left. \begin{array}{l} \text{In} \\ \text{Pre} \\ \text{Post} \end{array} \right\} \Theta(n)$

where $n = \# \text{ nodes}$.

- Queues:

Defn $\text{height}(T) = \text{dist. from root to a deepest leaf}$

$\text{height}(x) = \text{height of subtree rooted at } x$

note: $\text{height}(x) \leq \text{height}(T)$