

CSE 101-02 4-12-22

1

• Pa2 : ext. 2 days

Question: $\text{addEdge}()$ vs. $\text{addArc}()$

• Pa3 is posted.

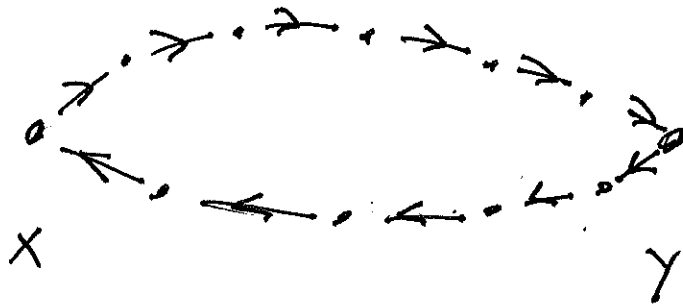
Let G be a digraph.

we say $y \in V(G)$ is reachable from $x \in V(G)$
iff G contains a directed x - y path.

Defn

A digraph G is called strongly connected iff for all $x, y \in V(G)$

x is reachable from y (and y from x .)



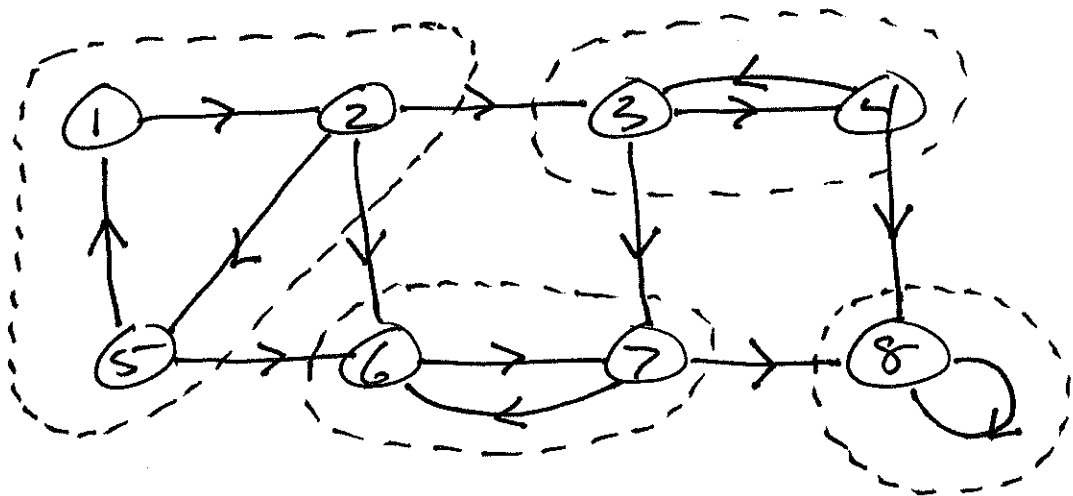
Also $S \subseteq V(G)$ is called strongly connected iff for all $x, y \in S$

x is reachable from y , and y from x .

A subset $S \subseteq V(G)$ is called a strongly connected component (scc) of G iff

- (i) S is strongly connected, and
- (ii) S is maximal w.r.t. (i).

Ex.

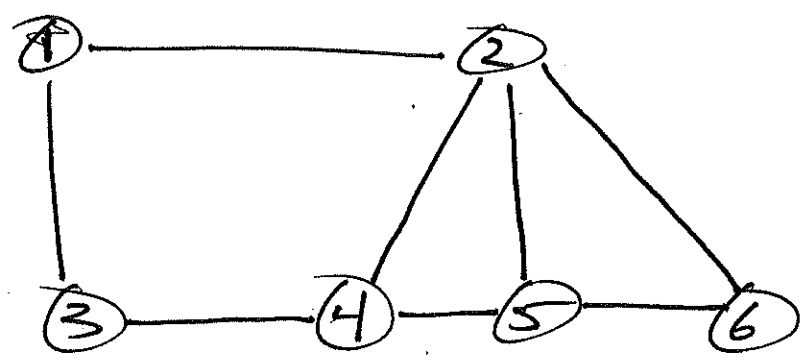


Depth First Search (DFS)

each vertex $x \in V$ has attributes

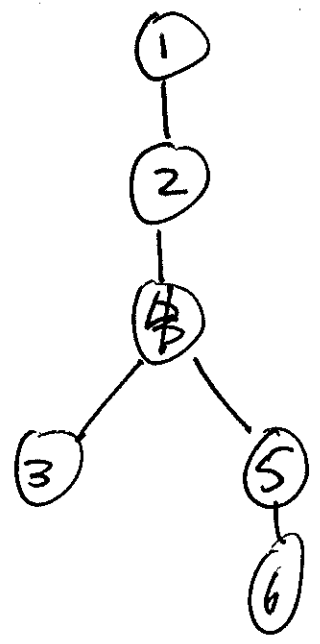
- $color[x]$: white, gray, black
- $P[x]$: Parent (Predecessor) of x .
- $d[x]$: discover time of x
- $f[x]$: finish time of x .
- recursive part: $visit(x)$
- DFS has a variable called time static over all calls to $visit$.

Ex.



	adj	color	d	f	p
✓ 1	2 3	w g b	1	12	11
✓ 2	1 4 5 6	w g b	2	11	1
✓ 3	1 4	w g b	4	5	4
✓ 4	2 3 5	w g b	3	10	2
✓ 5	2 4 6	w g b	6	9	4
✓ 6	2 5	w g b	7	8	5

DFS Forest



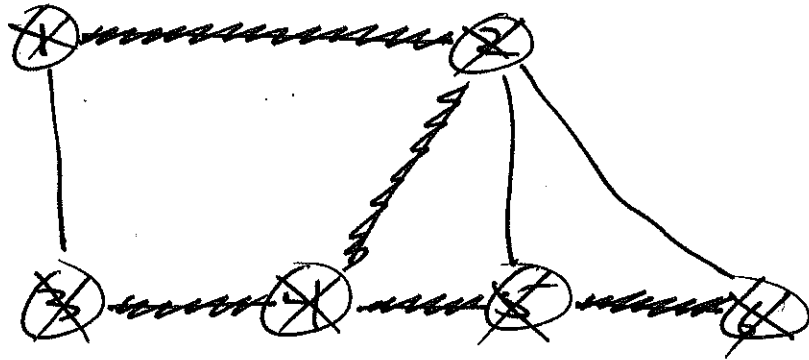
time

10
11
12

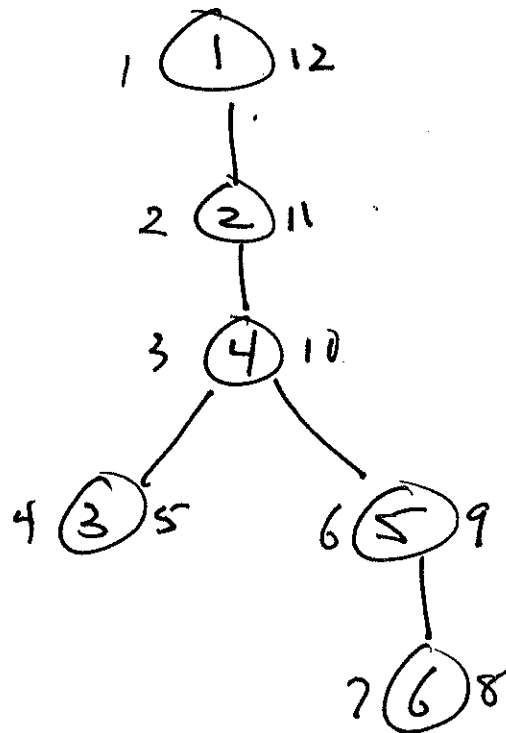
0
1
2
3
4
5
6
7
8
9

Quick way!

Ex.



DFS Forest!

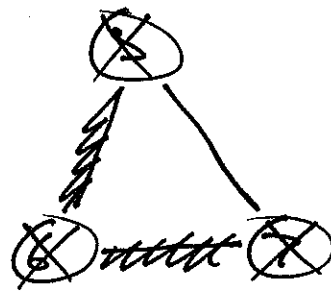
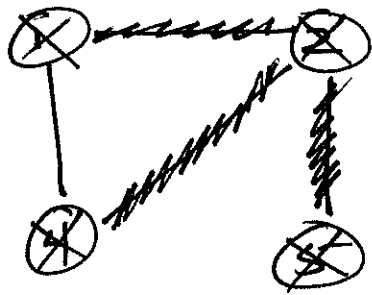


Predecessor subgraph $G_p = (V_p, E_p)$

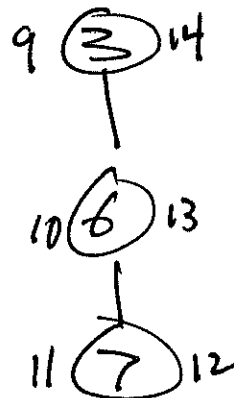
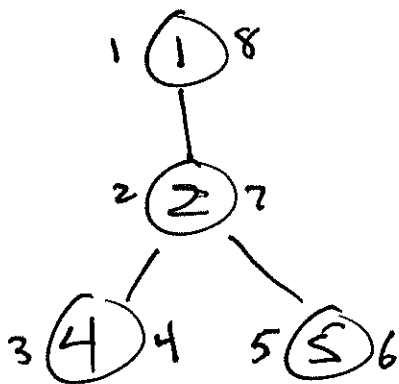
$$V_p = V(G)$$

$$E_p = \{ (p[x], x) \mid p[x] \neq \text{nil} \}$$

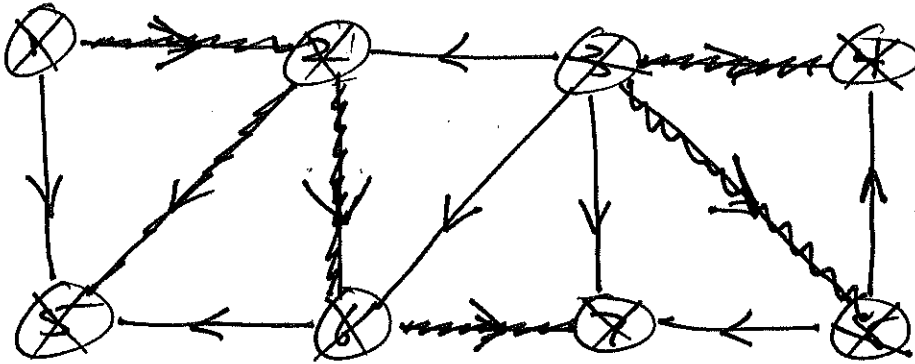
Ex.



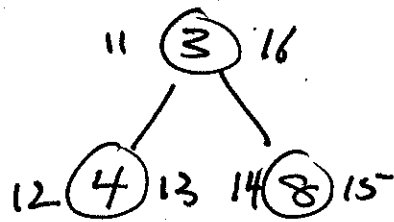
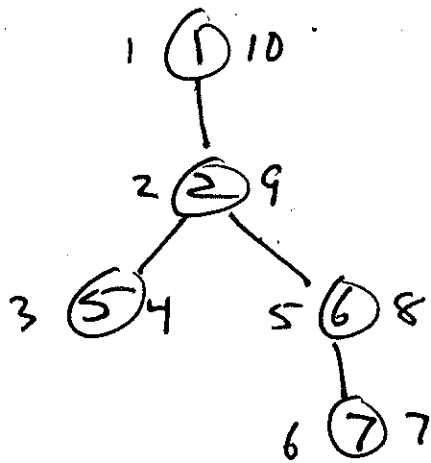
DFS Forest



Ex.



Forest:



Parenthesis string!

time: 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15-16
string: ((() (()))) (() ())