

css 101 11-19-20

1

Pass: ext 2 days

Pass: question

cleanUp() :

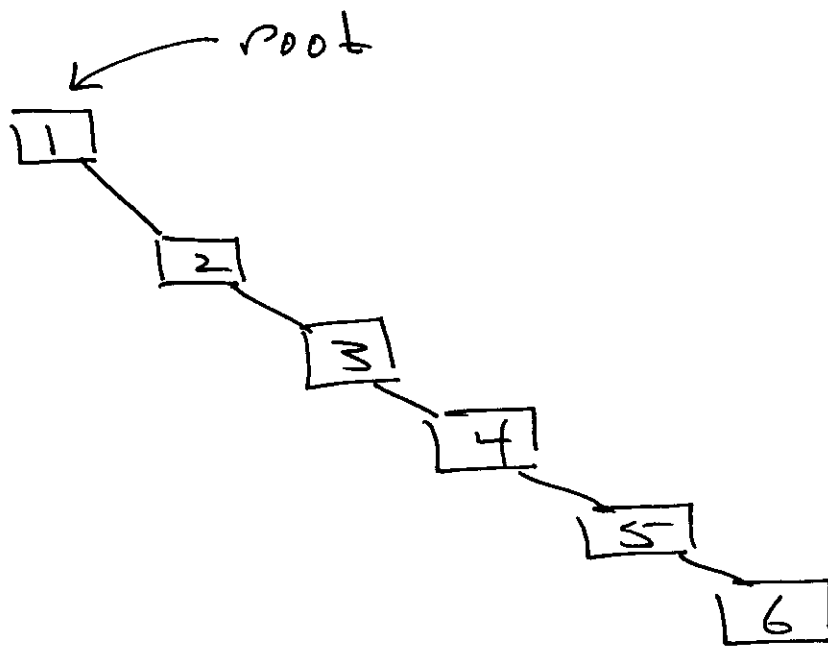
client view :

$$\textcircled{1} + \textcircled{2} = \mid = = \textcircled{3}$$

Quiz 4 : next Tue.

Page: Posted

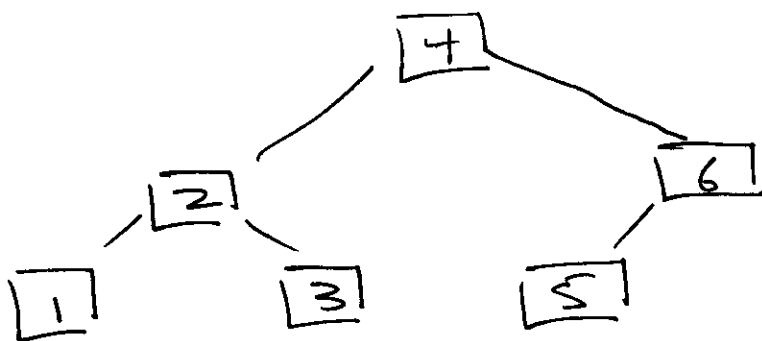
worst possible BST for
queries .



every node has only one child.

cost of query : $\Theta(n) = 6$

Best Possible BST :



cost of query : $\Theta(\text{height}(T)) = 2$
 general : $\Theta(\log n)$

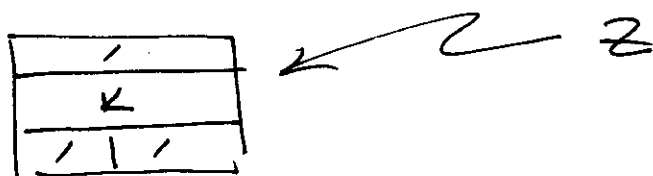
12.3 Insertion & Deletion

4

new node z

To insert 1 and maintain BST
Properties

- set $key[z] = k$
- set $P[z] = left[z] = right[z] = nil$.



- find a node to adopt z
as left or right child.

strategy:

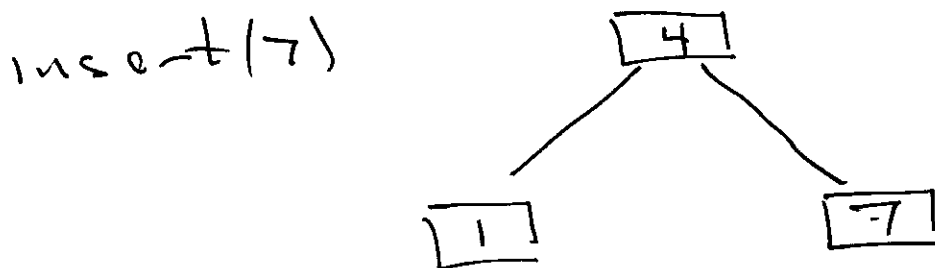
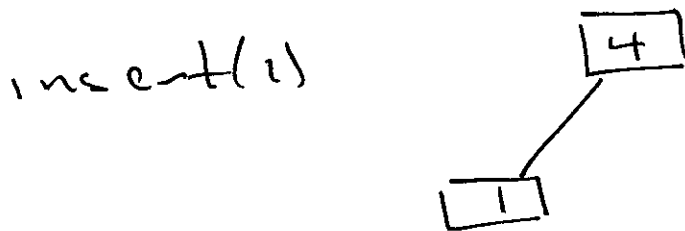
Do a search for key k in
 T , find nil child where k
would reside, if it were in T .

Ex. TreeInsert(T, z)

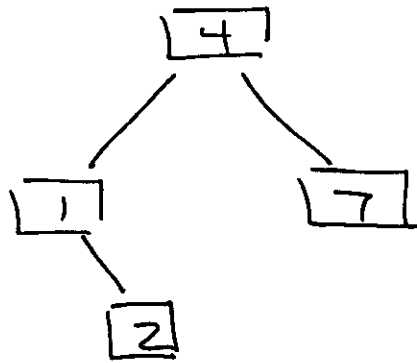
see pseudo-code

keys: ~~4~~ ~~1~~ ~~7~~ ~~2~~ ~~8~~ ~~3~~ ~~6~~ ~~5~~

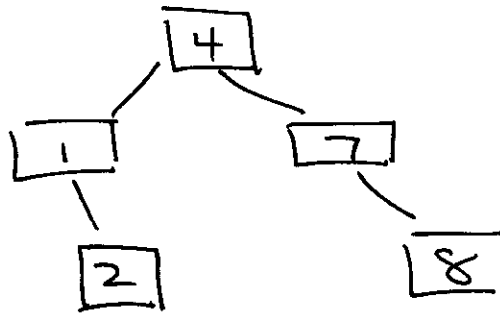
$T.root = nil$: empty state



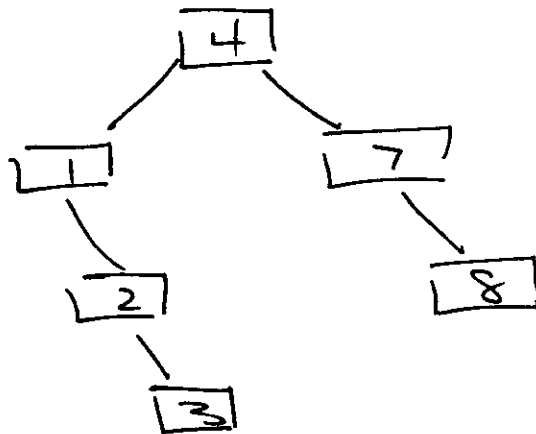
insert(2)



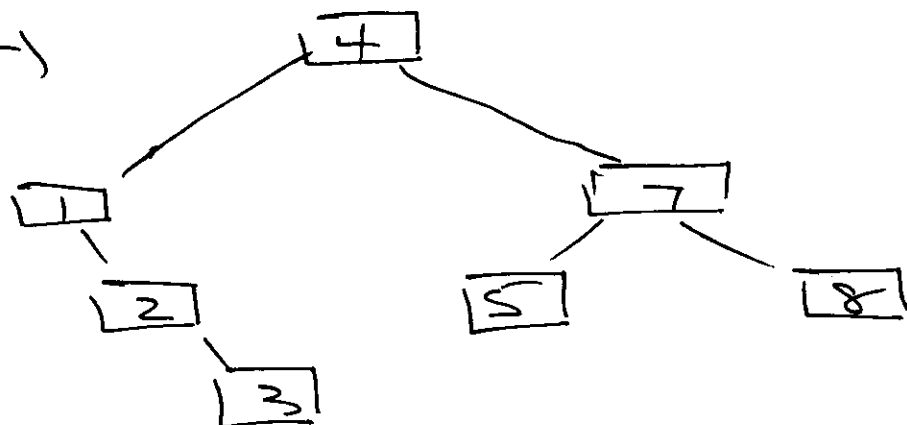
insert(8)



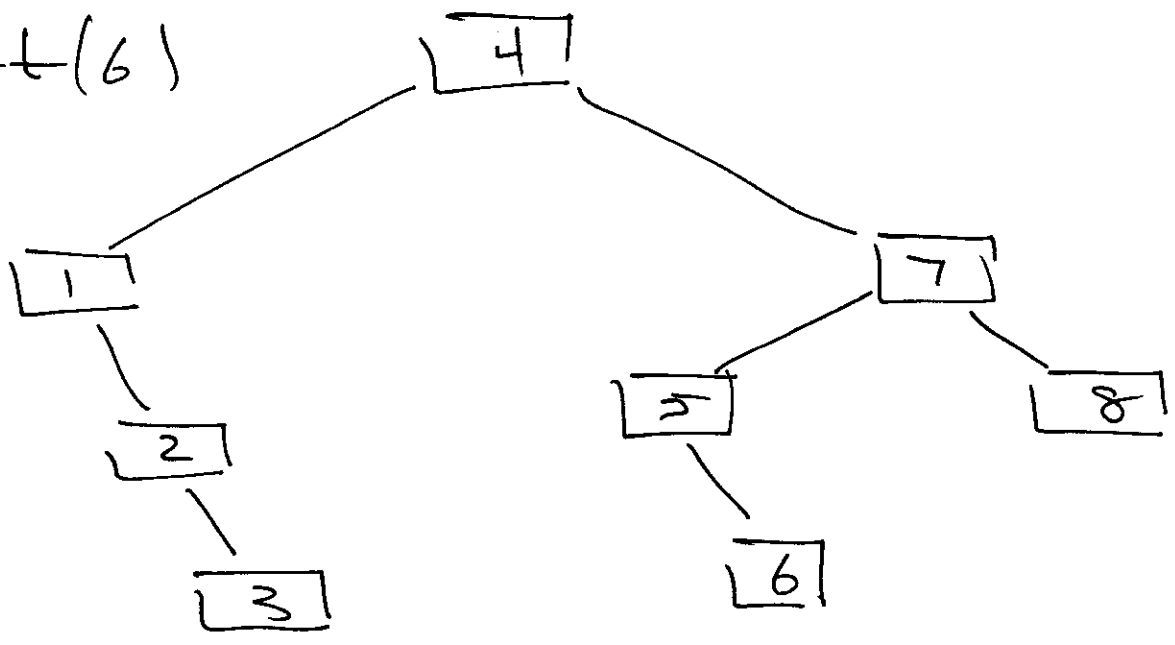
insert(3)



insert(5)



insert(6)



Runtime:

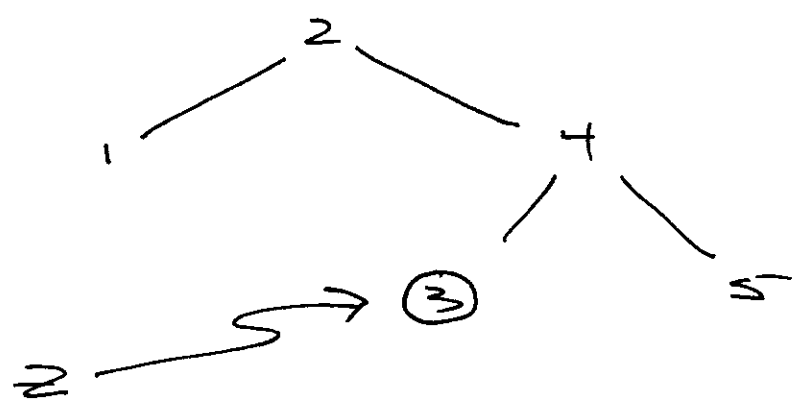
worst case cost : $\Theta(\text{height}(T))$.

Deletion: delete z from T

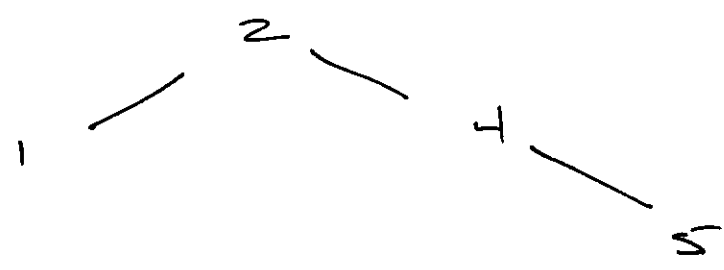
z case:

case 1: z is a leaf

detach z from its Parent



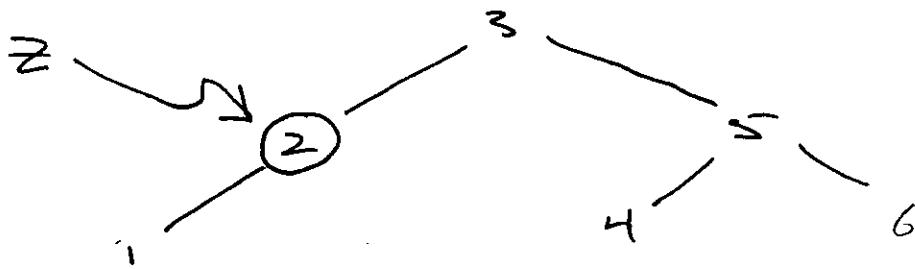
delete(z)



Case 2: z has 1 child.

Subcase 2.1 z has right but no left

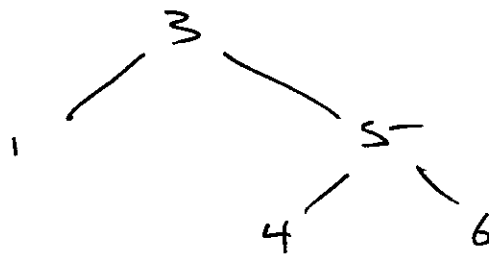
Subcase 2.2 z has left but no right



splice z out of 'local' linked list

delete(z):

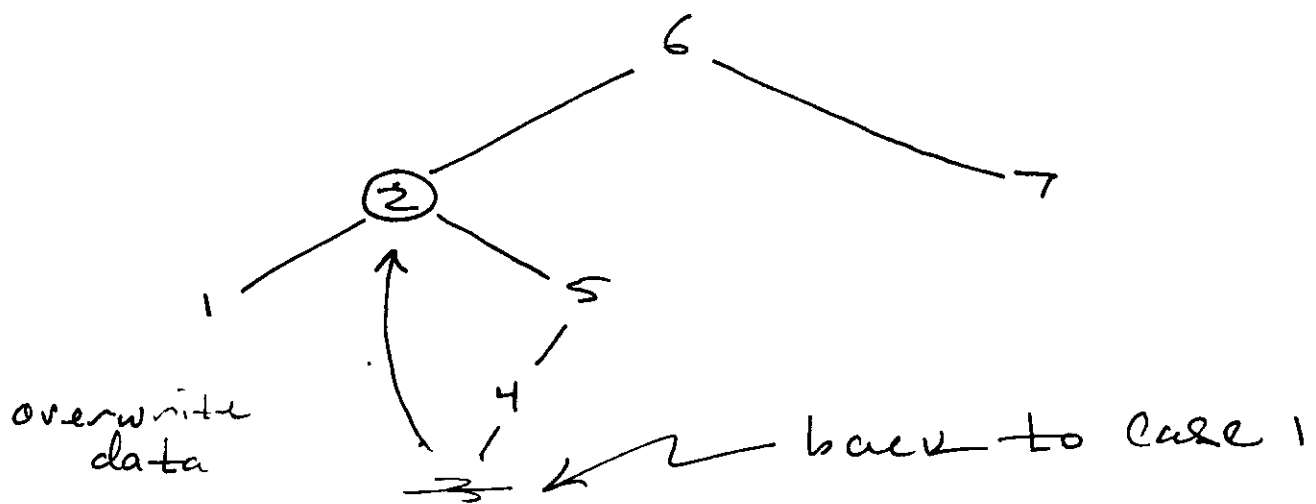
(Case 2.2)



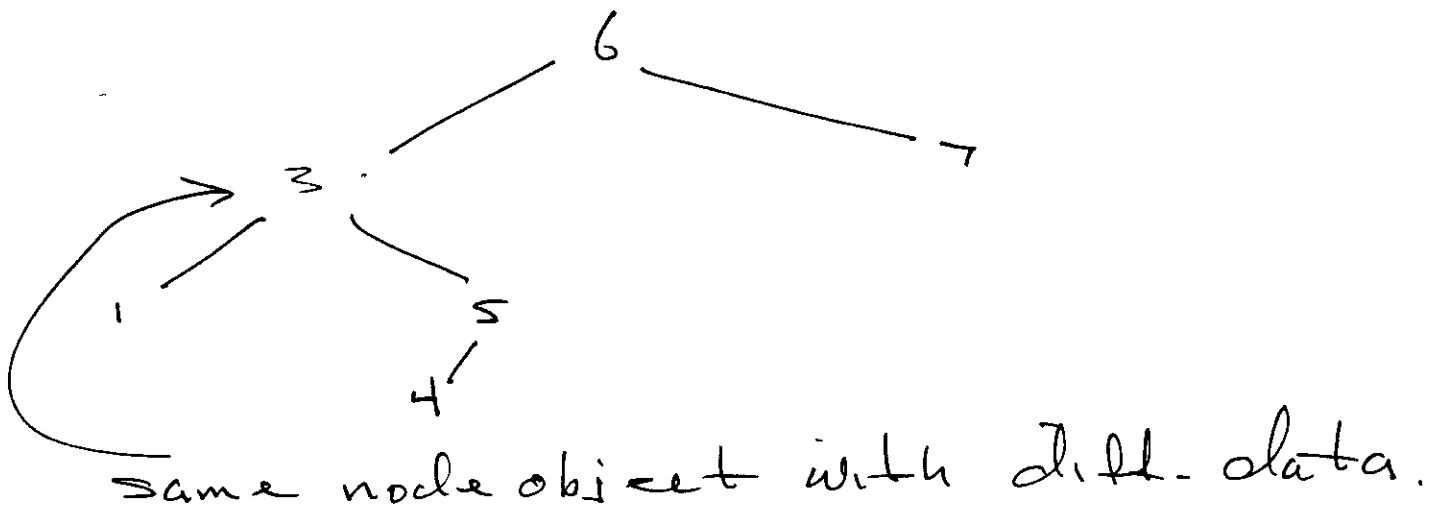
Exercise: draw example of subcase 2.1

Case 3: z has 2 children.

splice out z's successor y
(which has no left child),
then copy data of y into z.

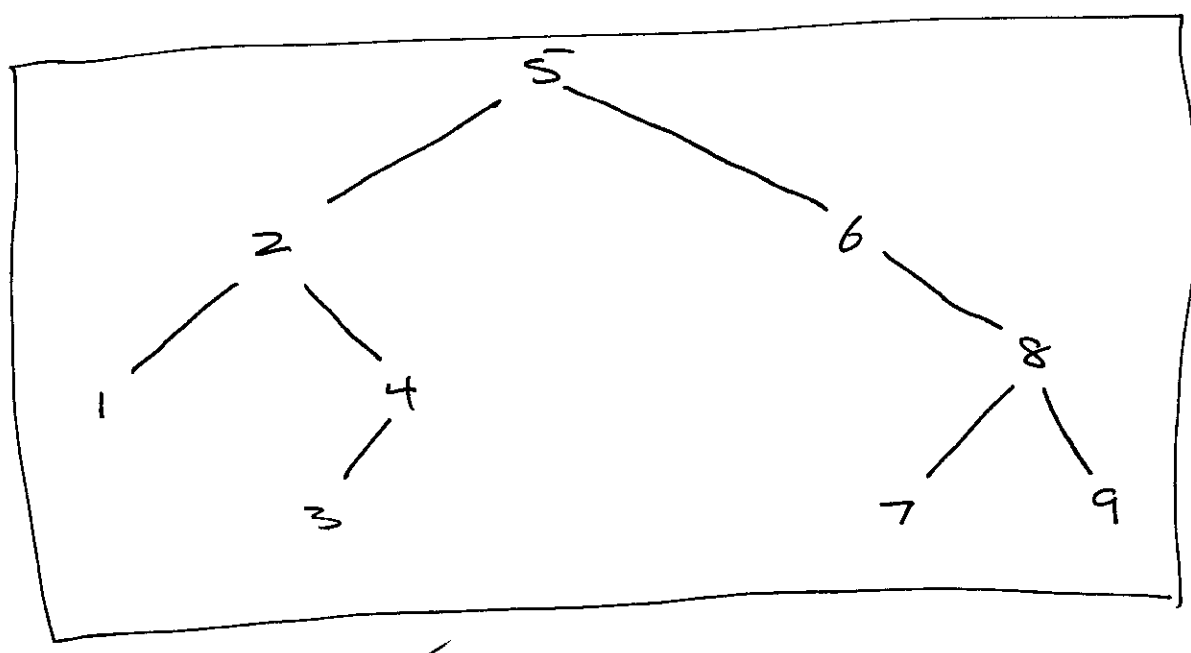


deleted 2)

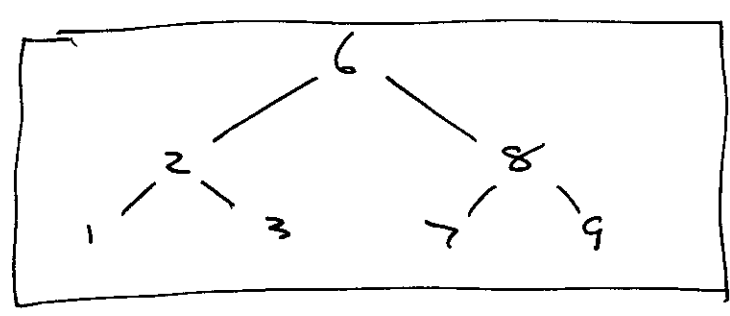
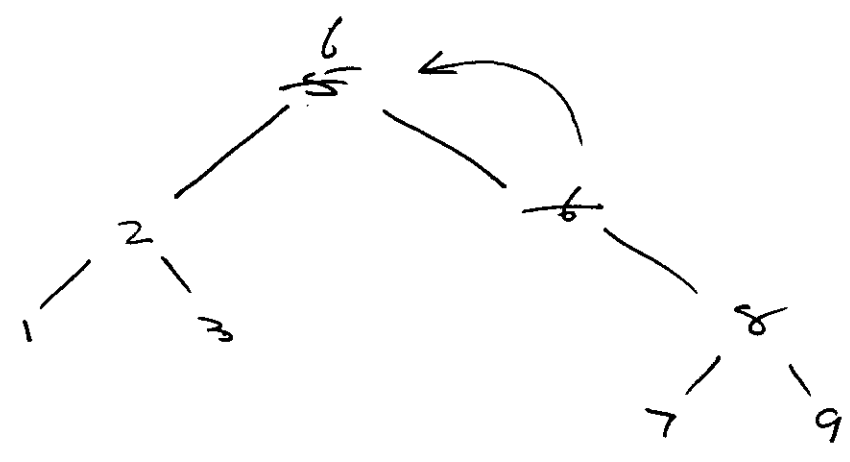


same node object with diff. data.

Ex. insert 5 2 6 4 3 8 7 1 9



Delete: 4, 5



Pseudo-code: next time.